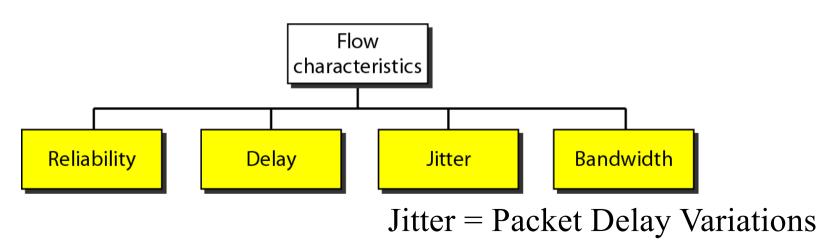
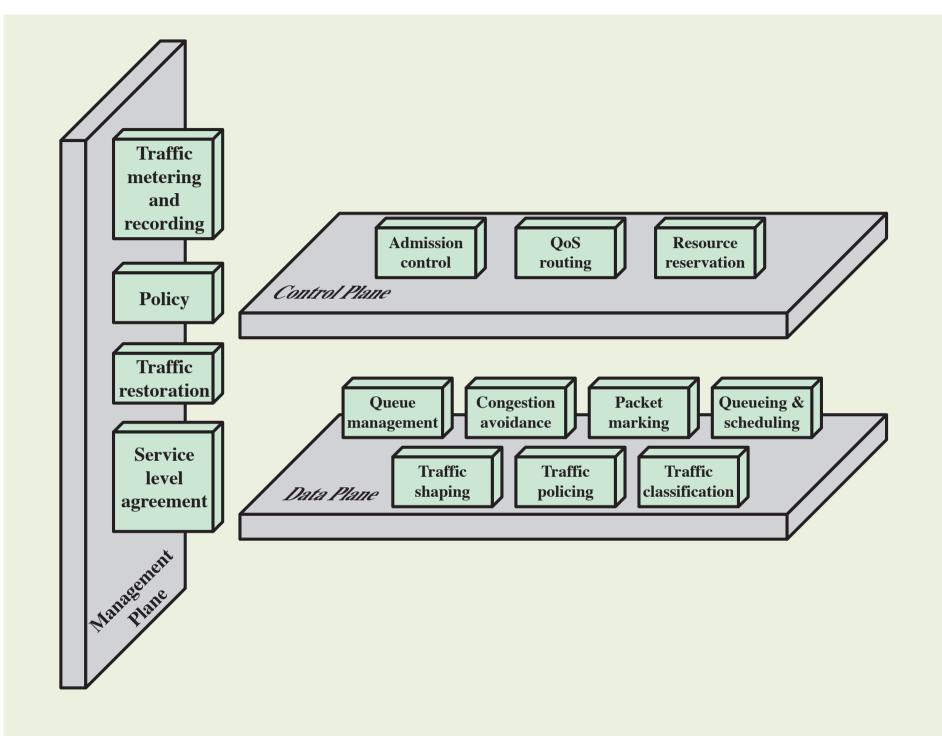


Quality of Service (QoS)

- Maintaining a functioning network
 - Meeting applications' demands
 - User's demands = QoE (Quality of Experience)
 - Dealing with flow characteristics





Data Plane

- Includes those mechanisms that operate directly on flows of data
 - Queue management algorithms
 - Tail drop vs RED (Random Early Detection)
 - Queueing and scheduling
 - Congestion avoidance
 - Packet marking
 - Traffic classification
 - Traffic policing
 - Traffic shaping

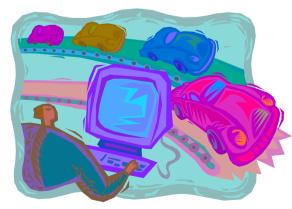
Control Plane

- Concerned with creating and managing the pathways through which user data flows
 - Admission control
 - QoS routing
 - Resource reservation



Management Plane

- Contains mechanisms that affect both control plane and data plane mechanisms
 - Service level agreement (SLA)
 - Traffic metering and recording
 - Traffic restoration
 - Policy



Network performance

- Data rate (Bandwidth)
 - Bits per second
- Throughput
 - Efficiency, always less than capacity (<1)
 - Alternatively: available data rate
- Latency (Delay)
 - Transmission, propagation, processing, queueing
 - One way or RTT (Round Trip time)
- PDV = Packet Delay Variation (Jitter)
 - Real-time applications!

Other parameters

- Bit Error Rate
 - L1 parameter that heavily impacts on L3
 - Frame/Packet Loss on higher layers
- Inter Packet Gap variations
 - "Jitter"
 - Could be non-zero already at sender
- Ratio of packets out of order
 - Impact on delay in TCP

Packet loss

- Due to
 - Bit error in packet
 - Routers discard erronous packet
 - Link or Physical Layer?
 - Queue overflow
 - Discard packets
 - Node problems
- In real-time multimedia late packets considered lost
- Packet loss ratio (%)
- Note TCP's sensitivity to packet loss

Circuit Switching	Datagram Packet Switching	Virtual Circuit Packet Switching
Dedicated transmission path	No dedicated path	No dedicated path
Continuous transmission of data	Transmission of packets	Transmission of packets
Fast enough for interactive	Fast enough for interactive	Fast enough for interactive
Messages are not stored	Packets may be stored until delivered	Packets stored until delivered
The path is established for entire conversation	Route established for each packet	Route established for entire conversation
Call setup delay; negligible transmission delay	Packet transmission delay	Call setup delay; packet transmission delay
Busy signal if called party busy	Sender may be notified if packet not delivered	Sender notified of connection denial
Overload may block call setup; no delay for established calls	Overload increases packet delay	Overload may block call setup; increases packet delay
Electromechanical or computerized switching nodes	Small switching nodes	Small switching nodes
User responsible for message loss protection	Network may be responsible for individual packets	Network may be responsible for packet sequences
Usually no speed or code conversion	Speed and code conversion	Speed and code conversion
Fixed bandwidth	Dynamic use of bandwidth	Dynamic use of bandwidth
No overhead bits after call setup	Overhead bitsein cach packetnte	nQverheadsbits in each packet

Table 9.1

Comparison
of Communication
Switching
Techniques

Figure 9.15 Event Timing for Circuit Switching and Packet Switching

Virtual Circuits vs. Datagram

- Virtual circuits
 - Network can provide sequencing and error control
 - Packets are forwarded more quickly
 - Less reliable (compare Circuit Switching)
- Datagram (Best Effort)
 - No call setup phase
 - Individual packet handling
 - More flexible
 - More reliable



IP Performance Metrics (IPPM wg)

- Chartered by IETF to develop standard metrics that relate to the quality, performance, and reliability of Internet data delivery
- Measurement techniques
 - Active: Transmit packets over network for measurement purposes
 - Passive: Use existing traffic for measurements

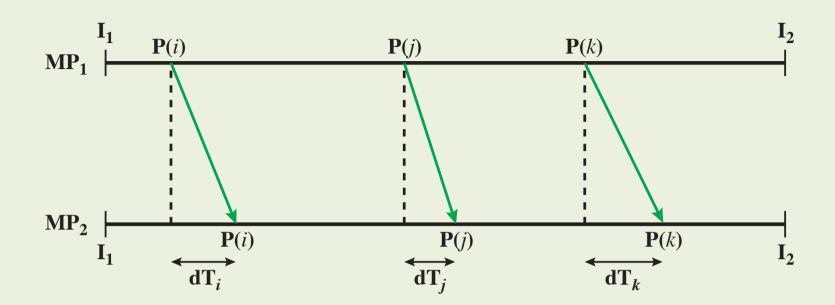
Table 22.3 IP Performance Metrics

Metric Name	Singleton Definition	Statistical Definitions
One-Way Delay	Delay = dT, where Src transmits first bit of packet at T and Dst received last bit of packet at T + dT	Percentile, median, minimum, inverse percentile
Round-Trip Delay	Delay = dT, where Src transmits first bit of packet at T and Src received last bit of packet immediately returned by Dst at T + dT	Percentile, median, minimum, inverse percentile
One-Way Loss	Packet loss = 0 (signifying successful transmission and reception of packet); = 1 (signifying packet loss)	Average
One-Way Loss Pattern	Loss distance: Pattern showing the distance between successive packet losses in terms of the sequence of packets Loss period: Pattern showing the number of bursty losses (losses involving consecutive packets)	Number or rate of loss distances below a defined threshold, number of loss periods, pattern of period lengths, pattern of inter-loss period lengths.
Packet Delay Variation	Packet delay variation (pdv) for a pair of packets with a stream of packets = difference between the one-way-delay of the selected packets Percentile, inverse percentile, jitter, peak-to-peak pdv	

Table 22.3 IP Performance Metrics

Metric Name	General Definition	Metrics
Connectivity	Ability to deliver a packet over a transport connection.	One-way instantaneous connectivity, Two-way instantaneous connectivity, one-way interval connectivity, two-way interval connectivity, two-way temporal connectivity
Bulk Transfer Capacity	Long-term average data rate (bps) over a single congestion-aware transport connection.	BTC = (data sent)/(elapsed time)

(b) Other metrics



 I_1 , I_2 = times that mark that beginning and ending of the interval in which the packet stream from which the singleton measurement is taken occurs.

 MP_1 , MP_2 = source and destination measurement points

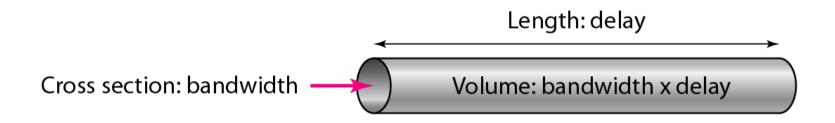
P(i) = ith measured packet in a stream of packets

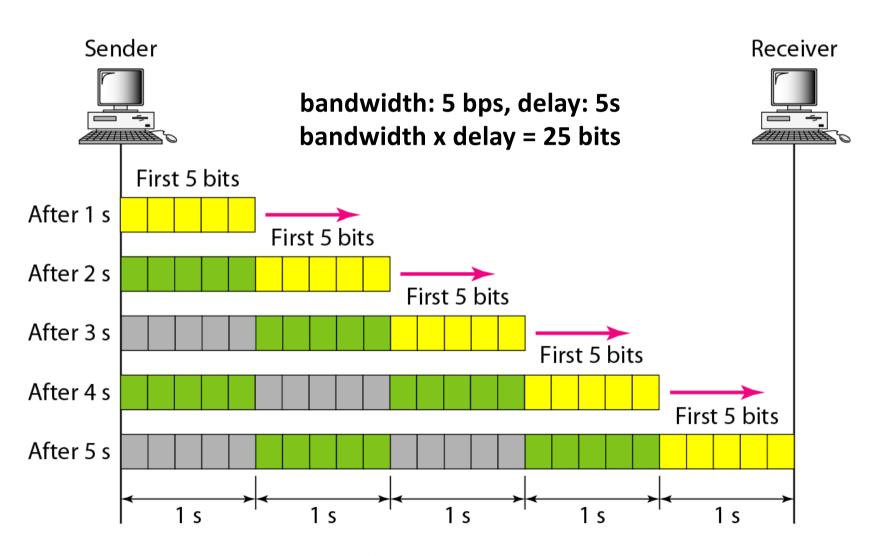
$$dT_i = \underbrace{\text{one-way delay for } P(i)}$$
 Time sync!

PDV_i=dT_i-dT_{i-1} Alt. STD(dT)

Figure 22.12 Model for Defining Packet Delay Variation

- How much data fills the link
- One Way Delay
- Two Way Delay = Round Trip Time (RTT)
 Time for data + time for ACK





- Important for congestion avoidance
 - Don't overfill the link
- Important for efficiency
 - Keep the link filled at all times
 - For max efficiencyData chunks > 2 * bandwidth * delay

- Important for tuning (TCP)
- Long Fat Network (LFN, "elephant")
 BDP >> 10⁵ bits

- Very long (high delay) links:
 - -> Bandwidth = BDP/delay
 - But it takes long time before ACK arrives ...

Performance vs ARQ

- Method
 - Stop-&-Wait
 - Go-Back-N
 - Selective-Repeat
- Utilisation = f(window size)

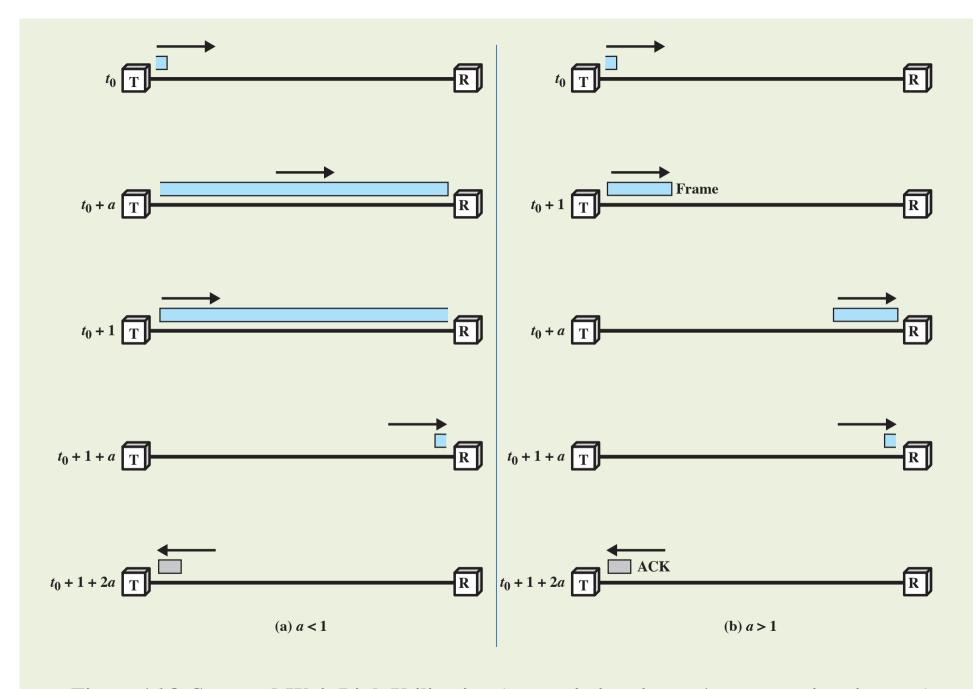
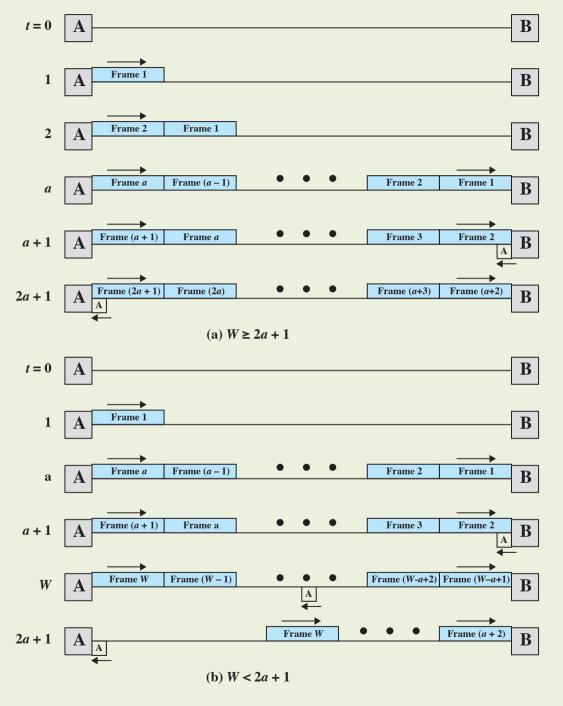


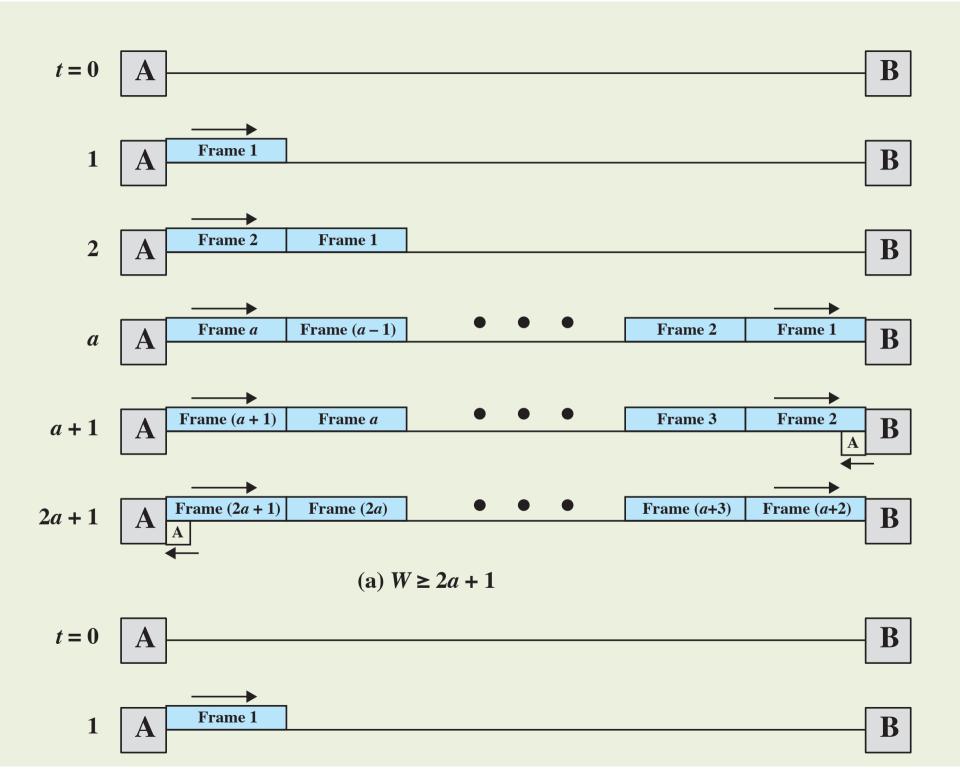
Figure 16.8 Stop-and-Wait Link Utilization (transmission time = 1; propagation time = a)

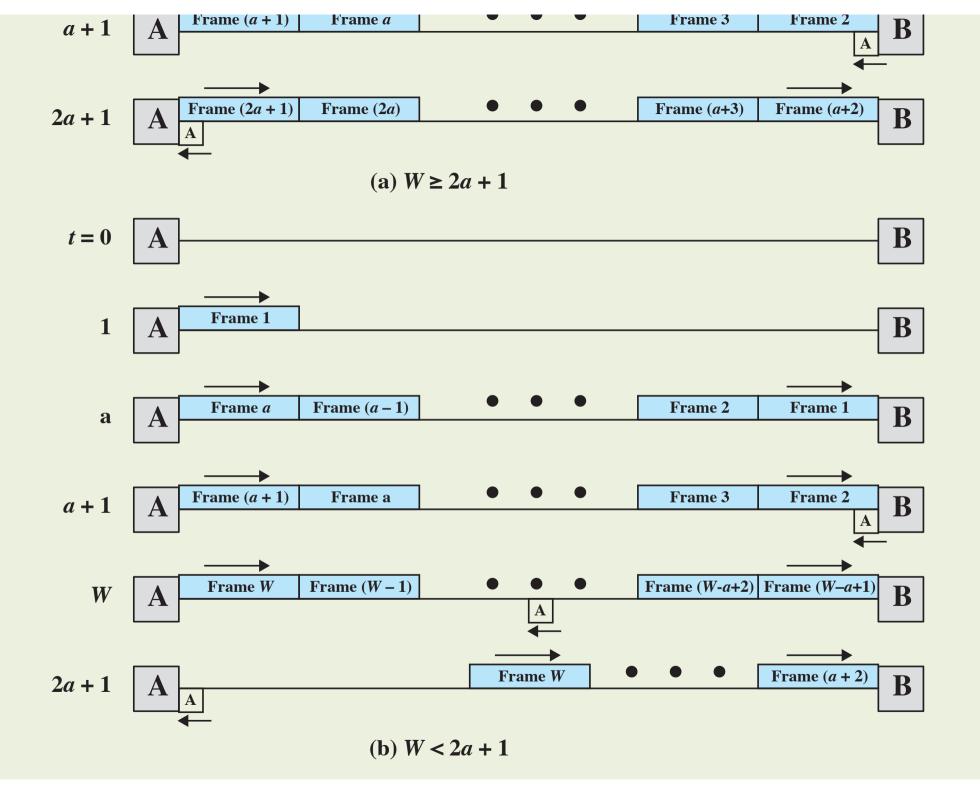


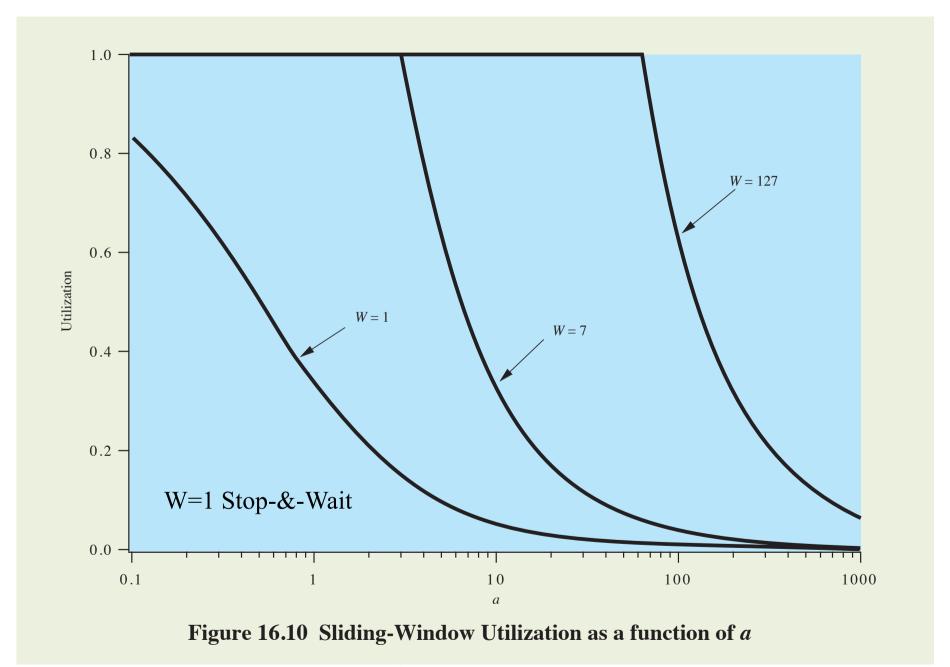
Sliding Windows based

- *a* = propagation time
- W = window size

 Compare with Bandwitdh-Delay Product







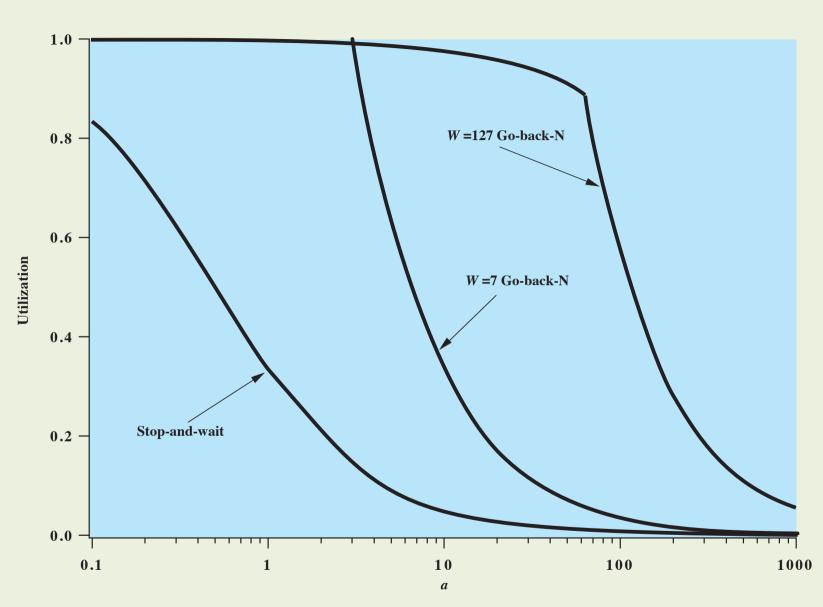
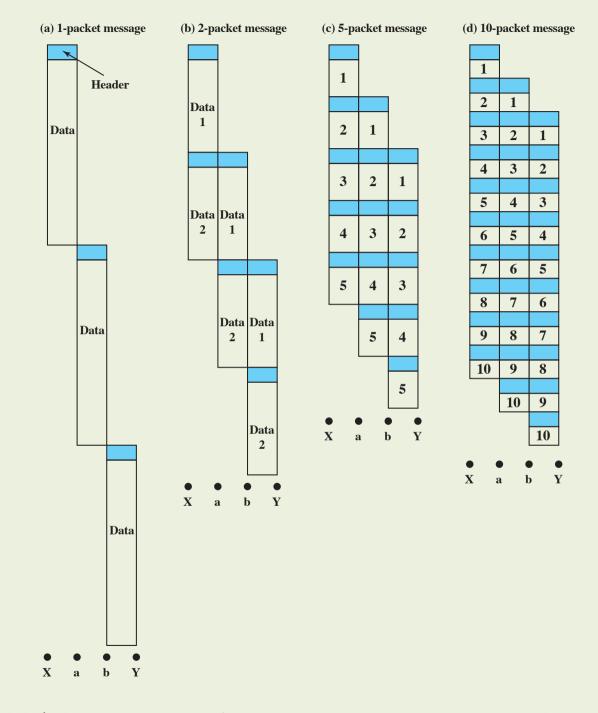


Figure 16.11 ARQ Utilization as a Function of a ($P = 10^{-3}$)

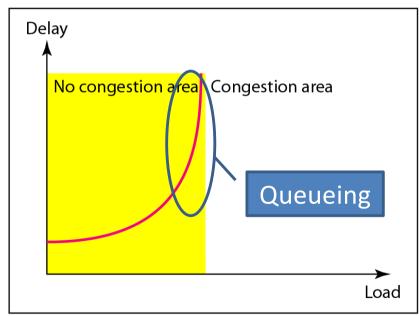
ETSF05/ETSF10 - Internet Protocols P = Propability of frame error

Packet Size vs Transmission Time

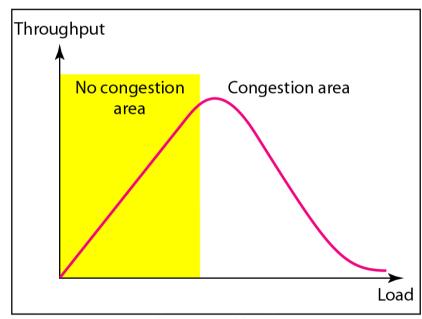


Delay and throughput

- Related to network load
 - Normal operation vs. congestion



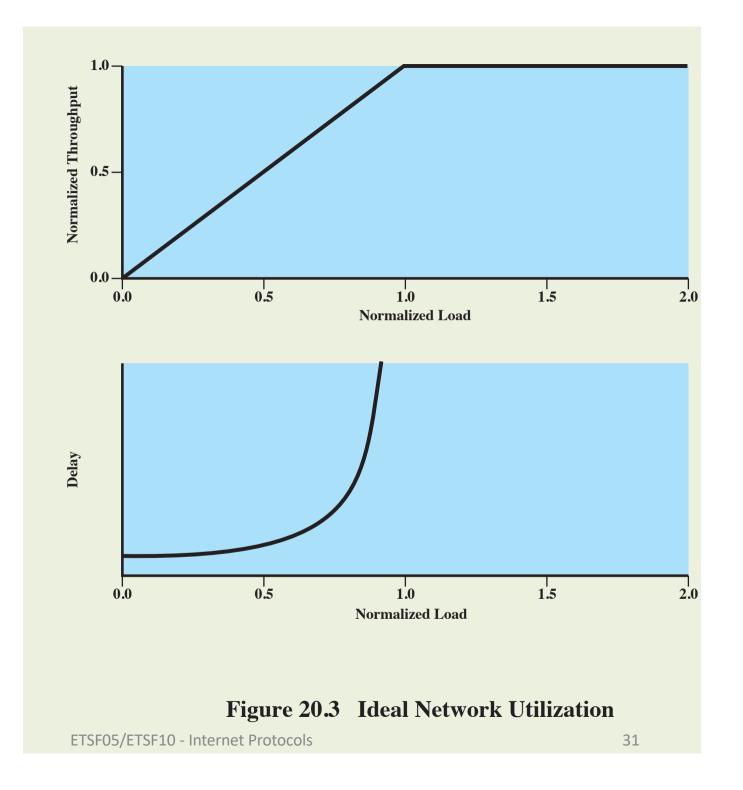




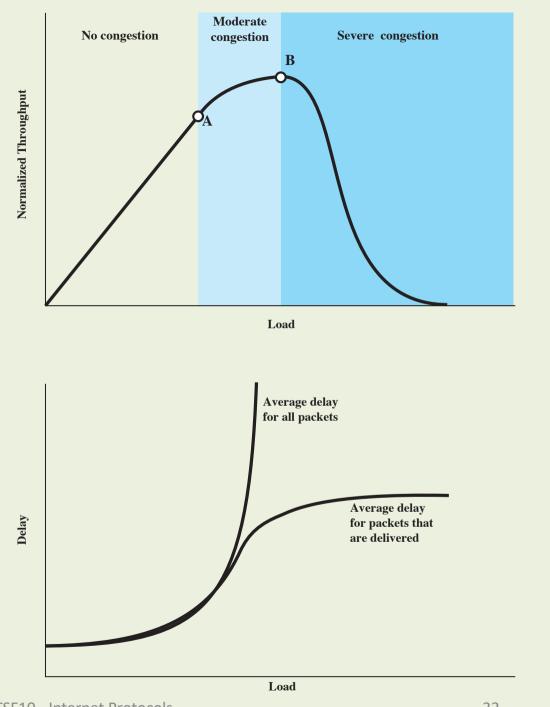
b. Throughput as a function of load

M/M/1 FIFO $\rho/(1-\rho)$

Delay and throughput: Ideal case



Delay and throughput: Finite buffers No congestion control



Nomenclature

- Throughput
 - All data successfully delivered end to end
- Goodput
 - Data successfully delivered to end application (without retransmissions)
- Goodput <= Throughput

Congestion Control in Packet-Switching Networks

Send control packet to some or all source nodes

Requires

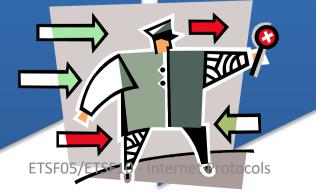
 additional
 traffic during
 congestion

Rely on routing information

May cause overreaction

End to end probe packets

Adds to overhead



Add congestion information to packets in transit

Either backwards or forwards

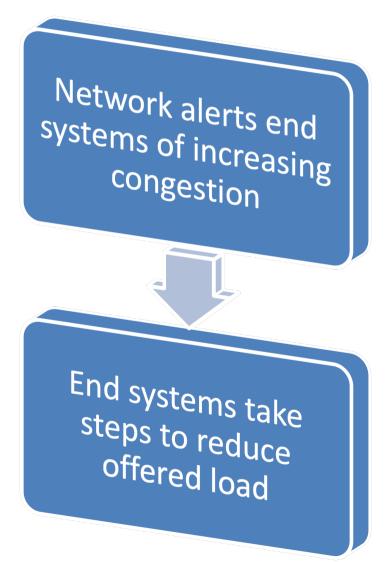
Explicit Congestion Signaling

Backward

 Congestion avoidance notification in opposite direction to packet required

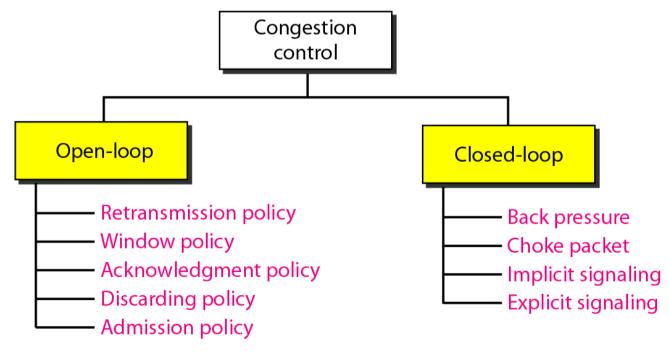
Forward

 Congestion avoidance notification in same direction as packet required



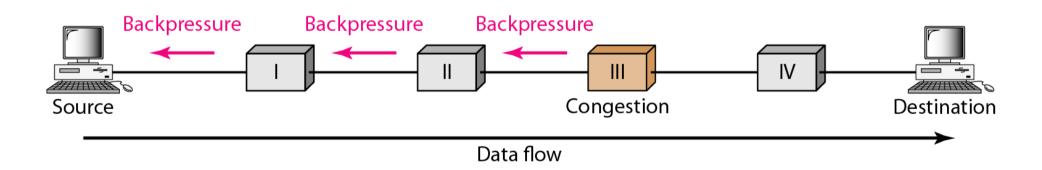
Congestion control

- Avoiding and eliminating congestion
 - Open-loop = proactive, prevent congestion
 - Closed-loop = reactive, control congestion



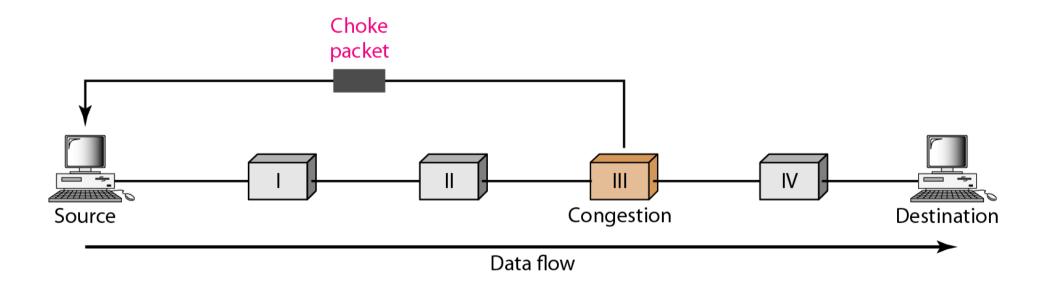
Closed-loop congestion control (1)

Backpressure



Closed-loop congestion control (2)

Choke packet



Implicit Congestion Signaling

- With network congestion:
 - Transmission delay increases
 - Packets may be discarded (Packet loss)
- Source can detect congestion and reduce flow
- Responsibility of end systems
- Effective on connectionless (datagram) networks
- Also used in connection-oriented networks

Explicit Signaling Categories

- Binary
 - A bit set in a packet indicates congestion
- Credit based
 - Indicates how many packets source may send
 - Common for end-to-end flow control
- Rate based
 - Supply explicit data rate limit
 - Nodes along path may request rate reduction

How to improve QoS?

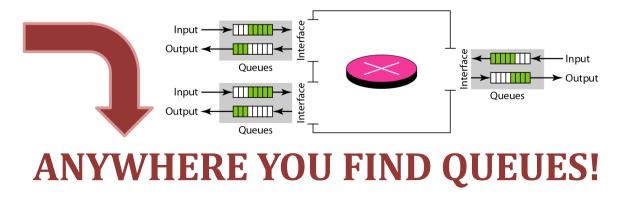
- Admission control
- Resource reservation
- Scheduling
- Traffic shaping

Routing?

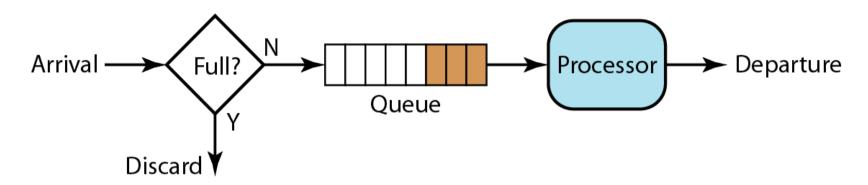
Where to improve QoS?

- Admission control
 - DIFFSERV: Service classes
 - INTSERV: Reservation architectures
- Resource reservation
 - RSVP (Resource Resevation Protocol)

- Scheduling
- Traffic shaping



Scheduling: FIFO queuing



Work conserving – always serve

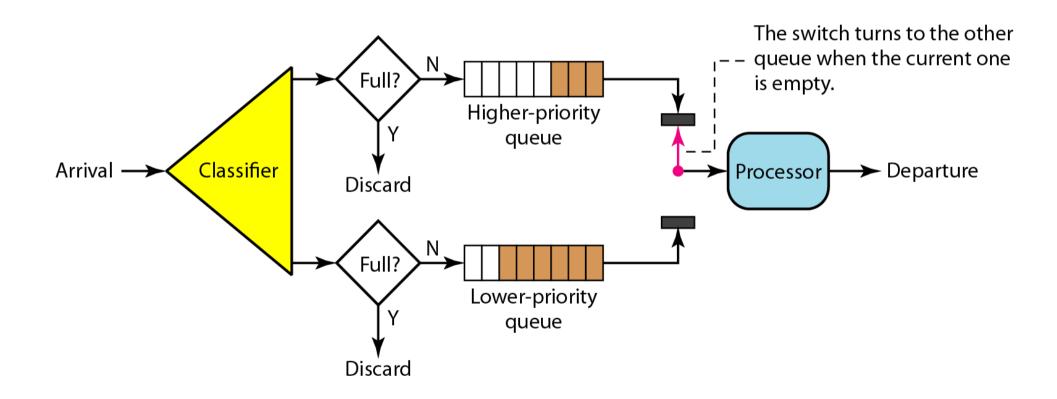
Kleinrock conservation law, ρ = utilisation

$$q = delay$$

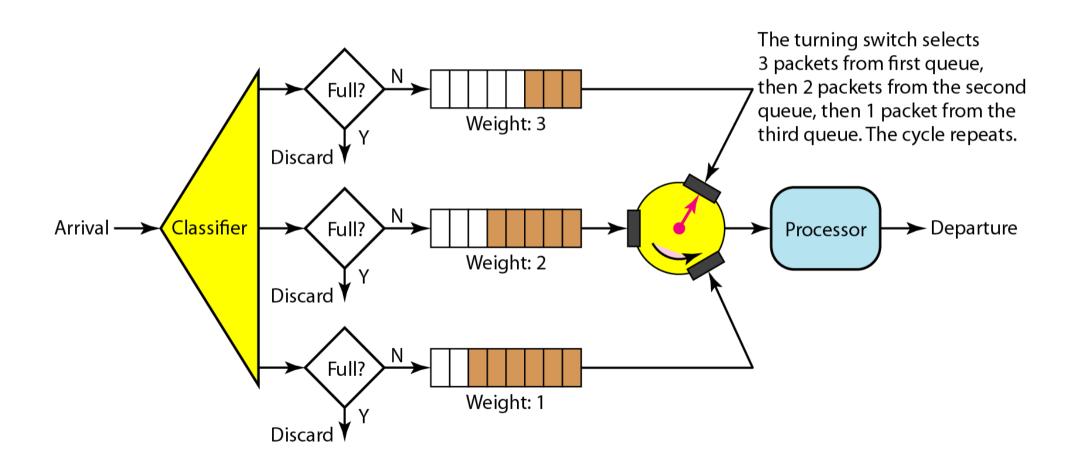
$$\sum_{n=1}^{N} \rho_n q_n = C$$

If delay for a flow is lowered, delay for some other flow has to increase

Scheduling: Priority queuing



Scheduling: Weighted fair queuing



Traffic Shaping/Traffic Policing

- Two important tools in network management:
 - Traffic shaping (Non-work conserving)
 - Concerned with traffic leaving the switch
 - Reduces packet clumping
 - Produces an output packet stream that is less bursty and with a more regular flow of packets

Traffic policing

- Concerned with traffic entering the switch
- Packets that don't conform may be treated in one of the following ways:



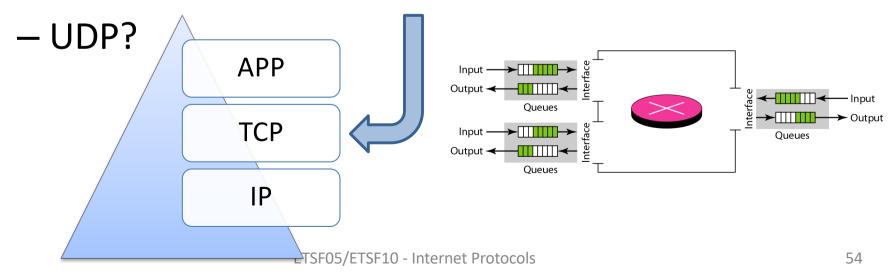
- Give the packet lower priority compared to packets in other output queues
- Label the packet as nonconforming by setting the appropriate bits in a header
- Discard the packet

Traffic Management

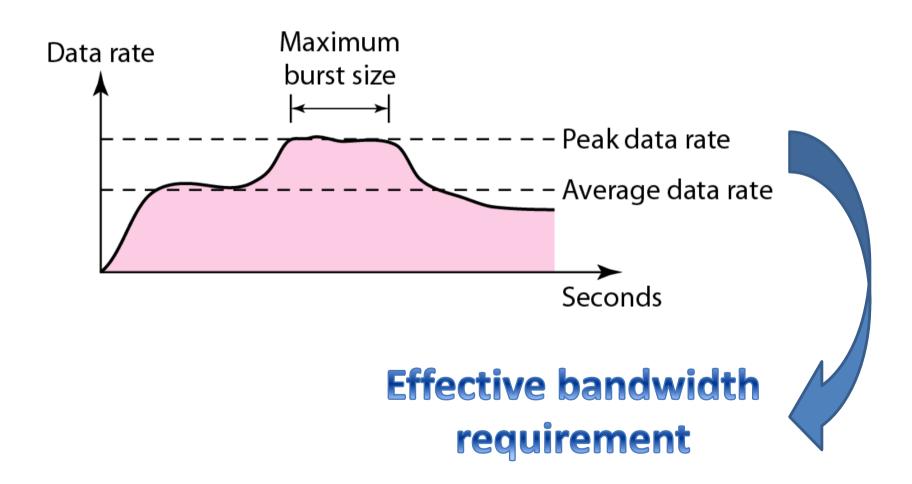
- Fairness
 - Provide equal treatment of various flows
- Quality of service
 - Different treatment for different flows
- Reservations
 - Traffic contract between user and network
 - Excess traffic discarded or handled on a best-effort basis

Layer 3 congestion avoidance

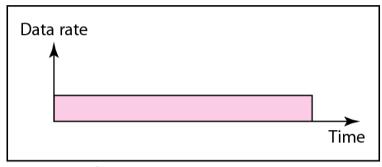
- Congestion = data load > network capacity
 - Arrival rate > processing rate
 - Processing rate > departure rate
- A simple method
 - Random early discard (RED)

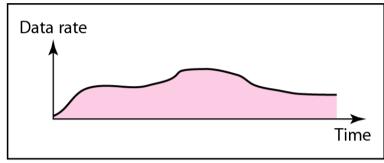


Traffic descriptors



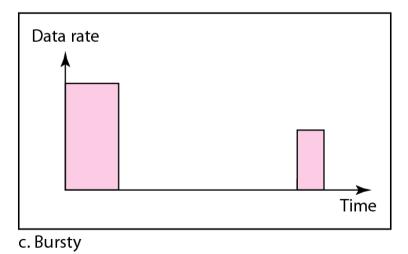
Traffic profiles





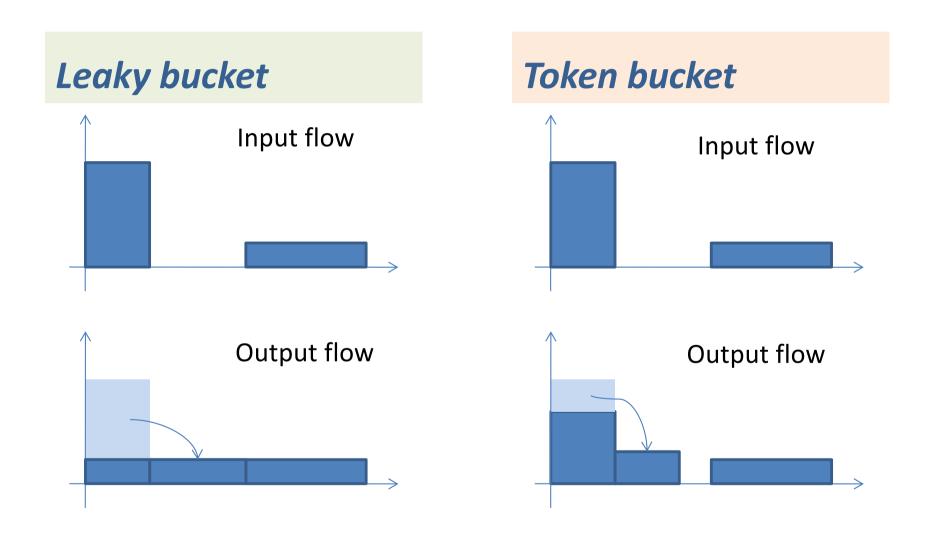
a. Constant bit rate

b. Variable bit rate

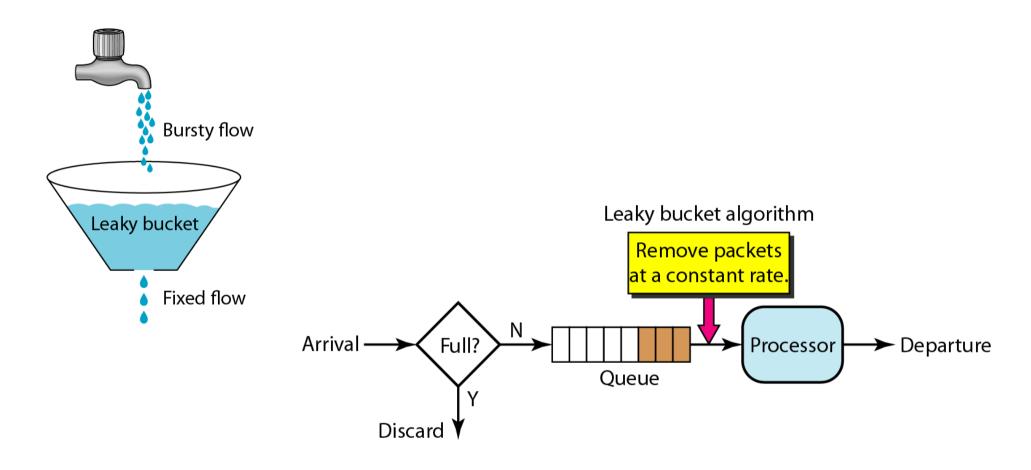


ETSF05/ETSF10 - Internet Protocols

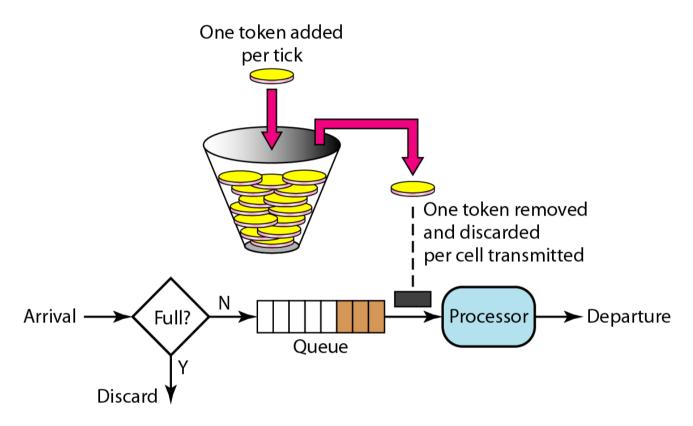
Traffic shaping: Two approaches



Traffic shaping: Leaky bucket



Traffic shaping: Token bucket



Token Bucket

- Widely used traffic management tool
- Advantages:
 - Many traffic sources can be defined easily and accurately
 - Provides a concise description of the load to be imposed by a flow, enabling the service to determine easily the resource requirement
 - Provides the input parameters to a policing function

Bonus: QoE, Quality of Experience

- The user's subjective perception of the presentation of the content
 - Example: Mean Opinion Score (MOS)
- What QoS give what QoE??