

## Point-to-Point Protocol Routing Algorithms Internet Addressing

ETSF05  
Internet Protocols  
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### Previously on ETS05

- Introduction
- Network topologies §1.1-2
- Network models §2.1-5
- Frames and data link control §11.1-5
- Local area networks and Ethernet §13.1-3

### Student Representative Election Day!

2011-09-26

2

### Today's lecture

- **Point-to-point protocol, PPP** §11.7
- Circuit-switched and datagram networks §8.1-2
- **Routing algorithms** §22.3
- IPv4 and IPv6 addresses, NAT §19.1-2
- Internetworking §20.1
- Address mapping, ARP §21.1

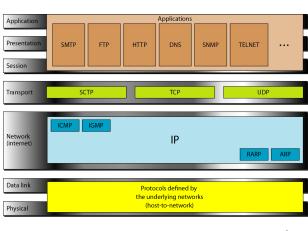
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3

## Point-to-point protocol (PPP)

- Direct connection between two nodes
    - Internet access

- Home user to ISP
  - Telephone line



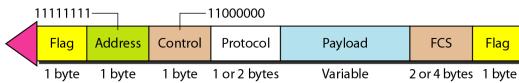
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4



## PPP frame format

- Support for several (sub)protocols
  - Address & control not used
  - Maximum payload 1500 bytes



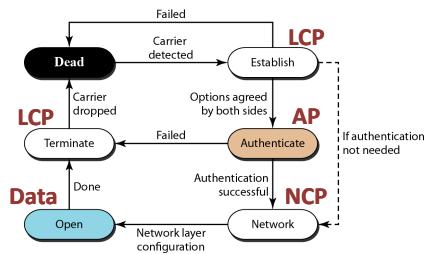
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5



## State transitions in PPP

- That's why we need (sub)protocols

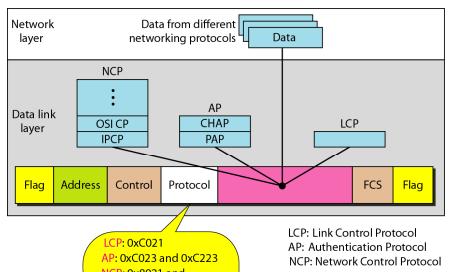


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6



## (Sub)protocols in PPP



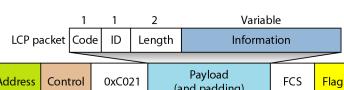
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7

## Link control protocol (LCP)

- Establish
- Configure
- Terminate

Code	Packet Type	Description
0x01	Configure-request	Contains the list of proposed options and their values
0x02	Configure-ack	Accepts all options proposed
0x03	Configure-nak	Announces that some options are not acceptable
0x04	Configure-reject	Announces that some options are not recognized
0x05	Terminate-request	Request to shut down the line
0x06	Terminate-ack	Accept the shutdown request
0x07	Code-reject	Announces an unknown code
0x08	Protocol-reject	Announces an unknown protocol
0x09	Echo-request	A type of hello message to check if the other end is alive
0x0A	Echo-reply	The response to the echo-request message
0x0B	Discard-request	A request to discard the packet

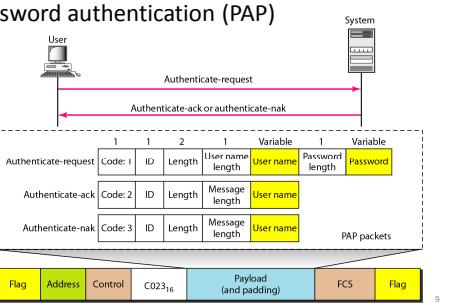


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8

## Authentication protocols (AP)

- Password authentication (PAP)

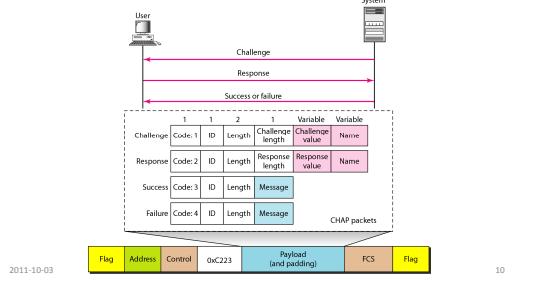


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9

## Authentication protocols (AP)

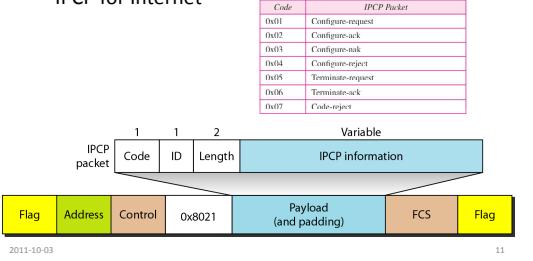
- Challenge handshake authentication (CHAP)



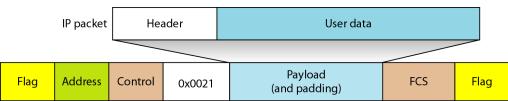
## Network control protocols (NCP)

- Preparations for the network layer

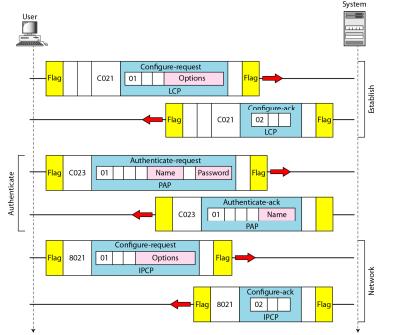
– IPCP for Internet



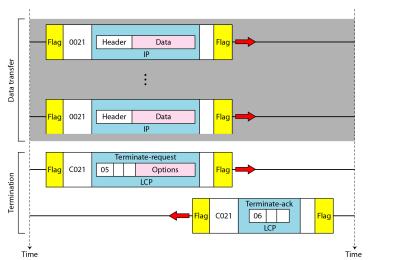
## IP datagram encapsulation in PPP



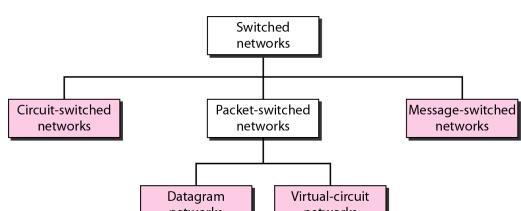
### PPP session example



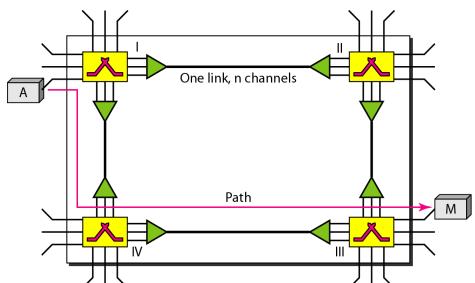
### PPP session example (cont.)



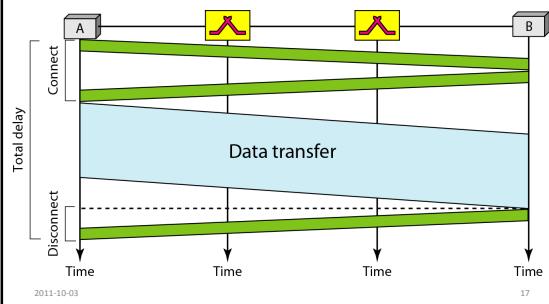
### Switching



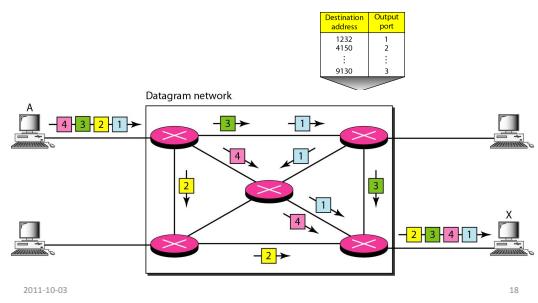
### Circuit-switched networks



### Phases of circuit switching

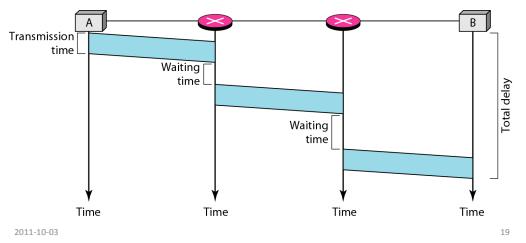


### Datagram networks



## Datagram transmission

- Higher delay
  - Higher throughput

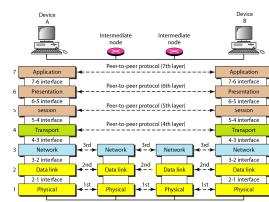


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19

## Routing

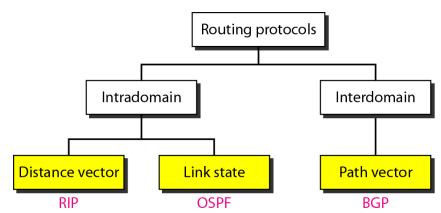
- Choosing the optimal path
    - Using a cost metric
  - Algorithms
    - Rules and procedures
    - Updates
  - Sharing information
    - Central
    - Distributed



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20

## Routing protocol classification



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23

## Distance vector routing

- Best path info shared locally
    - Periodically
    - Upon any change
  - Routing tables updated for
    - New entries
    - Cost changes

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22

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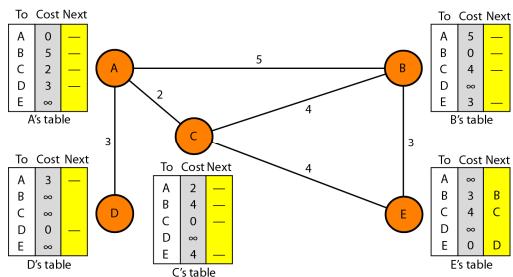
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## Initialisation of routing tables



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23

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## Updating algorithm

```

if (advertised destination not in table)
{
    add new entry // rule #1
}
else if (adv. next hop = next hop in table)
{
    update cost // rule #2
}
else if (adv. hop count < hop count in table)
{
    replace old entry // rule #3
}

```

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24

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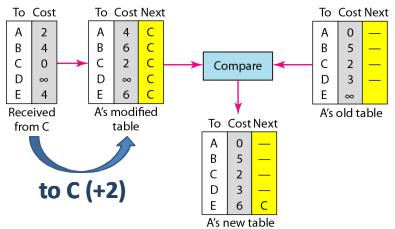
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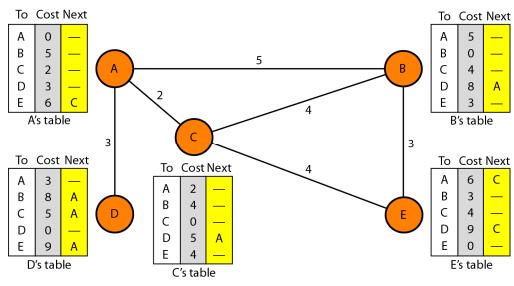
## Updating a routing table



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25

## Distance vector concept



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26

See you in 15' :)



- Now:  
**Student Representative Election**
- After the break
  - IPv4/IPv6 addresses
  - Internetworking
  - Address mapping

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27

## Link state routing

- Local topology info **flooded** globally
    - Periodically
    - Upon any change
  - Routing tables updated for
    - New entries
    - Cost changes

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28

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29

## Updating algorithm (Dijkstra)

```

put yourself to tentative list
while tentative list not empty
{
    pick node with least cumulative cost
    put it to permanent list
    add its least cost link to tree
    put its neighbours to tentative list
        (if not already there)
}

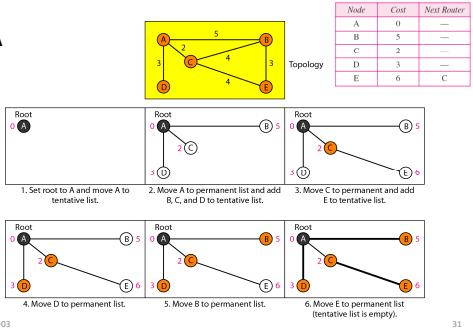
```

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30

## Shortest path tree generation

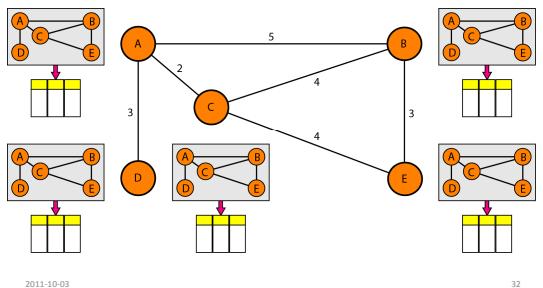
- A



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31

## Link state concept

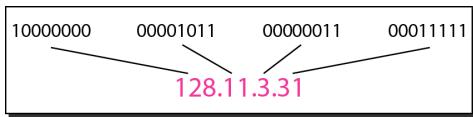


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32

## IPv4 addresses

- 32 bits = 4 bytes
- $2^{32} = (2^8)^4 = 256^4 = 4\ 294\ 967\ 296$
- Notations



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33

## Classful addressing

- 3 main classes for all organisation types
  - Problem: not flexible
- Netid, hostid

	First byte	Second byte	Third byte	Fourth byte
Class A	0			
Class B	10			
Class C	110			
Class D	1110			
Class E	1111			

a. Binary notation

	First byte	Second byte	Third byte	Fourth byte
Class A	0-127			
Class B	128-191			
Class C	192-223			
Class D	224-239			
Class E	240-255			

b. Dotted-decimal notation

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34

## Classless addressing

- Addresses in blocks
  - Count a power of 2
  - First address divisible by address count

Block	
First →	205.16.37.32
	205.16.37.33
	⋮
Last →	205.16.37.47

a. Decimal

Block	
11001101	00010000
11001101	00010000
⋮	⋮
11001101	00010000
00100101	00101111

b. Binary

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35

## Classless addressing example

- 205.16.37.39/28

Address:	11001101 00010000 00100101 00100111
Mask:	<b>11111111 11111111 11111111 11110000</b>
First address:	11001101 00010000 00100101 00100000
Address:	11001101 00010000 00100101 00100111
Mask complement:	<b>00000000 00000000 00000000 00001111</b>
Last address:	11001101 00010000 00100101 00101111

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36

## Network address translation (NAT)

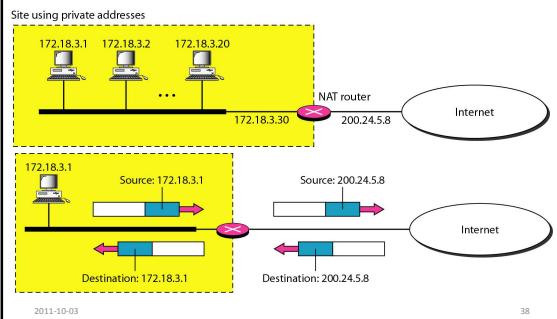
- Private Internet usage increasing
  - Address needs becoming permanent
- Not enough addresses!
  - Solution: Separate internal from external
  - Let private networks do address translation

Range	Total
10.0.0.0 to 10.255.255.255	$2^{24}$
172.16.0.0 to 172.31.255.255	$2^{20}$
192.168.0.0 to 192.168.255.255	$2^{16}$

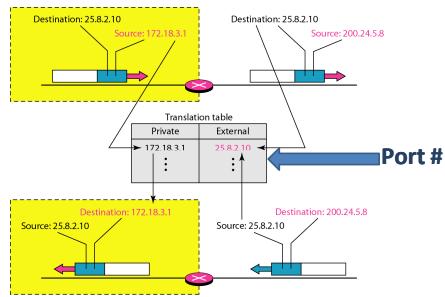
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37

## NAT implementation

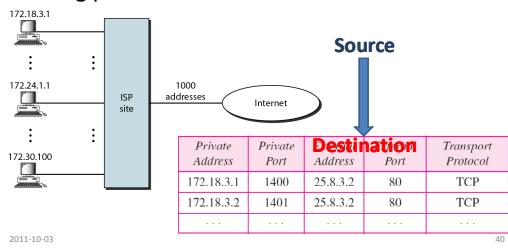


## NAT address translation



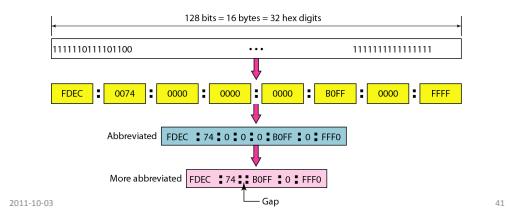
## NAT improvements

- Using multiple external addresses
- Using port numbers



## IPv6 addresses

- 128 bits = 16 bytes
- $2^{128} = (2^{16})^8 = 65\,536^8 > 3 \cdot 10^{38}$
- Notations

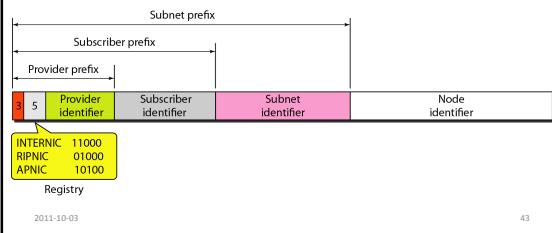


## IPv6 address categories

Type Prefix	Type	Fraction
0000 0000	Reserved	1/256
0000 0001	Unassigned	1/256
0000 001	ISO network addresses	1/128
0000 010	IPX (Novell) network addresses	1/128
0000 011	Unassigned	1/128
0000 1	Unassigned	1/32
0001	Reserved	1/16
001	Reserved	1/8
<b>010</b>	<b>Provider-based unicast addresses</b>	<b>1/8</b>

## IPv6 unicast addresses

- Individual computer
  - Address type, area, ISP, organisation, subnet, user



## A few special IPv6 addresses

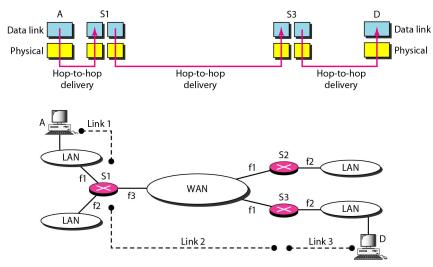
8 bits	120 bits	a. Unspecified
00000000	All 0s	a. Unspecified
8 bits	120 bits	b. Loopback
00000000	0000000000000000.....0000000001	b. Loopback
8 bits	88 bits	c. Compatible
00000000	All 0s	IPv4 address
8 bits	72 bits	d. Mapped
00000000	All 0s	All 1s
10 bits	70 bits	IPv4 address
111111010	All 0s	Node address
10 bits	38 bits	a. Link local
111111011	All 0s	Subnet address
		Node address
		b. Site local

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44

## Internetworking

- L2 is host-to-host

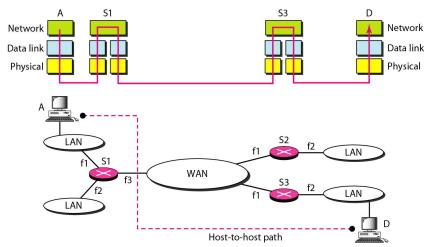


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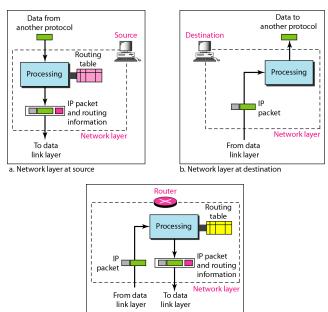
45

## Network layer

- L3 is end-to-end

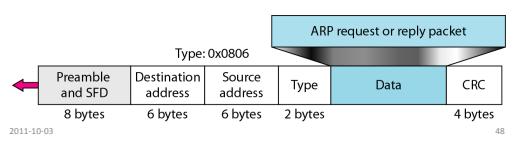


## Routing at network layer

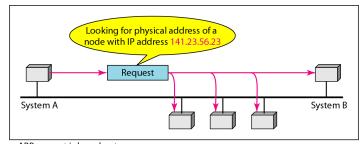


## Address mapping

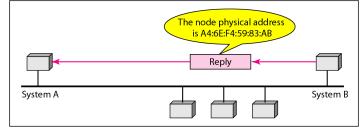
- Internetworking**
  - Network of networks (Internet)
  - Connected by routers
- Routers need information**
  - Logical (IP)  $\leftrightarrow$  physical (MAC)



## Address resolution protocol (ARP)

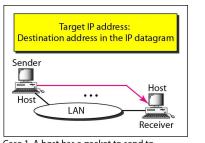


a. ARP request is broadcast

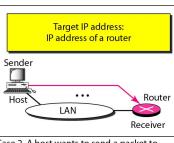


b. ARP reply is unicast

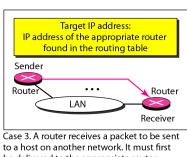
## Four use cases for ARP



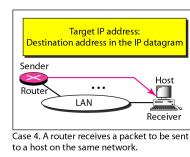
Case 1. A host has a packet to send to another host on the same network.



Case 2. A host wants to send a packet to another host on another network. It must first be delivered to a router.



Case 3. A router receives a packet to be sent to a host on another network. It must first be delivered to the appropriate router.



Case 4. A router receives a packet to be sent to a host on the same network.

## Coming up next week

- Transmission modes §4.3
- Multiplexing §6.1
- SONET/SDH §17
- Asynchronous Transfer Mode, ATM §18.2-3