

Projects in Wireless Communication

Lecture 1

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Outline

- ▷ Introduction to the course
- ▷ Basics of digital communications
- ▷ Discrete-time implementations
- ▷ Carrier transmission

Introduction

Lecturer and course responsible: **Fredrik Tufvesson**, E:2361A
7 scheduled lectures

Teaching assistant:
Guoda Tian, E:2367
Computer help sessions over Zoom.

Email: `firstname.lastname@eit.lth.se`

Introduction

Ultimate goals for PWC:

- 1) Two computers should communicate via speaker/microphones
- 2) Two computers should communicate using software defined radios

The projects should be performed in groups of **ONE** or **TWO** students



PWC System Simulation

In the first part of PWC we only work in software. For a passing grade you must solve three tasks:

1. A digital baseband BPSK system should be implemented in MATLAB and its performance should be measured and verified against theoretical results

$$P_e = Q \left(\sqrt{d_{\min}^2 \frac{E_b}{N_0}} \right)$$

2. Later in PWC you will encounter physical passband signals at the input of the microphone. In the first part, we will provide each group with one such signal; the bits carried by the signals correspond to the ASCII code of a secret password. If you can decode the signals and provide me with the password, you have passed task 2.

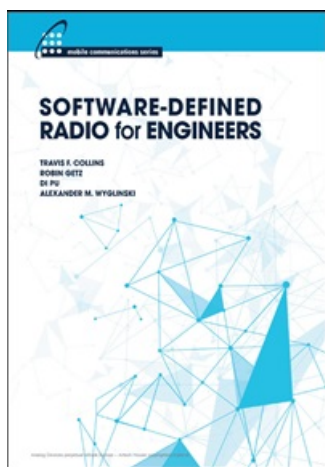
3. Same as 2 but with OFDM transmission and convolutional code.

Recommended reading

Software-Defined Radio for Engineers,

by Travis F. Collins, Robin Getz, Di Pu, and Alexander M. Wyglinski,
2018, ISBN-13: 978-1-63081-457-1.

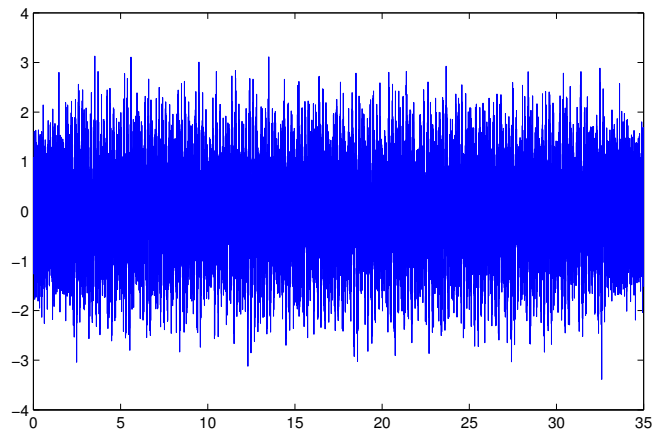
We will use some chapters in the second half of the course, and it covers many of the aspects in the first half as well.



There is a free pdf of the book available, see
<http://www.analog.com/en/education/education-library/software-defined-radio-for-engineers.html>

Example

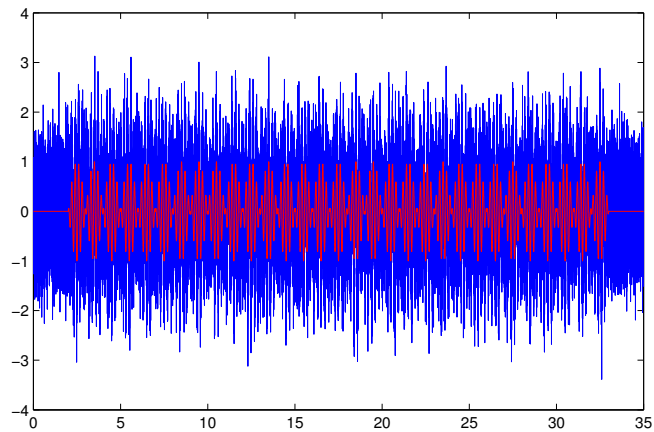
Assume that you receive the following noisy signal



You must remove the noise...

Example

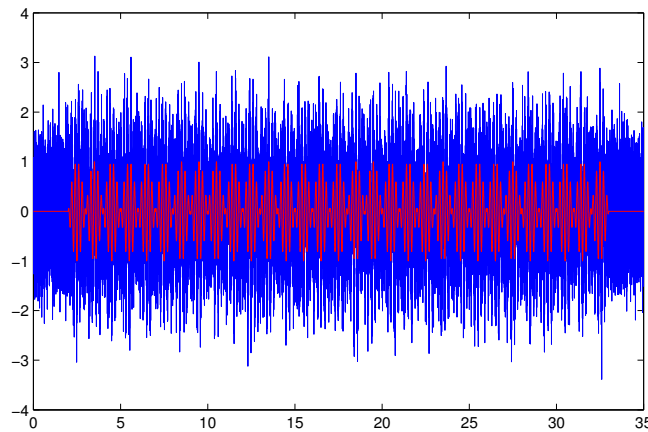
Assume that you receive the following noisy signal



You must remove the noise...Done!

Example

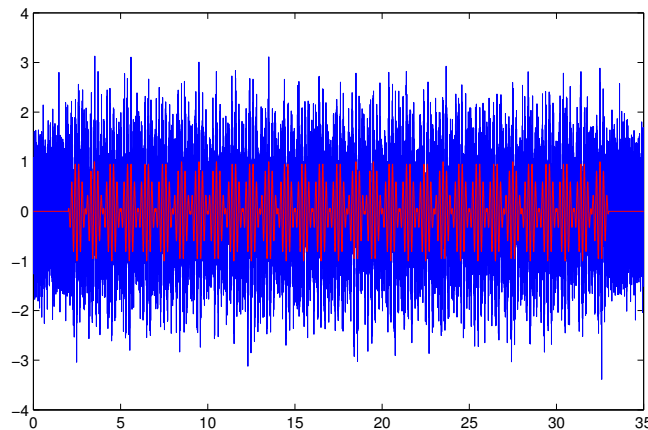
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You must remove the noise...Done!
Decode the bits:

Example

Assume that you receive the following noisy signal

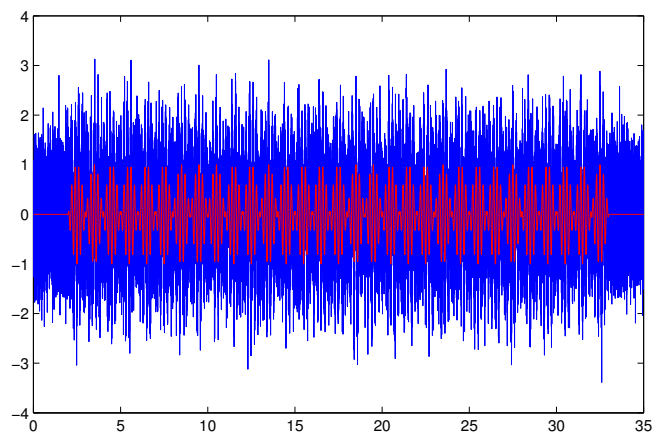


You must remove the noise...Done!

Decode the bits: 1 1 1 0 0 0 0 1 1 0 1 0 0 1 1 1 0 1.....

Example

Assume that you receive the following noisy signal



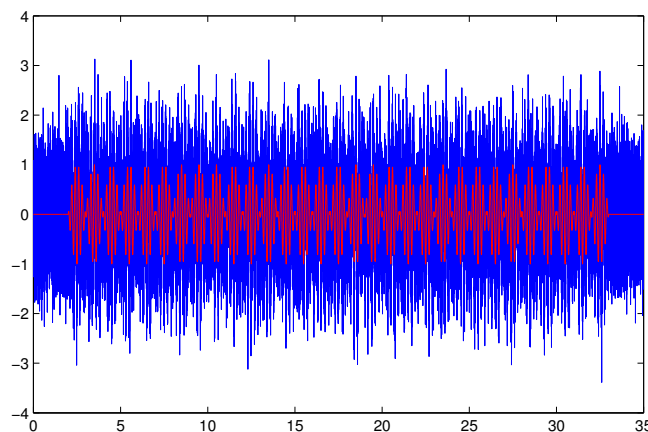
You must remove the noise...Done!

Decode the bits: 1 1 1 0 0 0 0 1 1 0 1 0 0 1 1 1 0 1.....

Convert to ASCII:

Example

Assume that you receive the following noisy signal



You must remove the noise...Done!

Decode the bits: 1 1 1 0 0 0 0 1 1 0 1 0 0 1 1 1 0 1.....

Convert to ASCII: You have passed PWC1, congratulations.....

Introduction

Formal descriptions of the tasks can be found online.

Basics of Digital Communications

This is a recall of baseband digital communications....

We need to transmit a bit sequence $\{u_k\} = 0111010\dots$

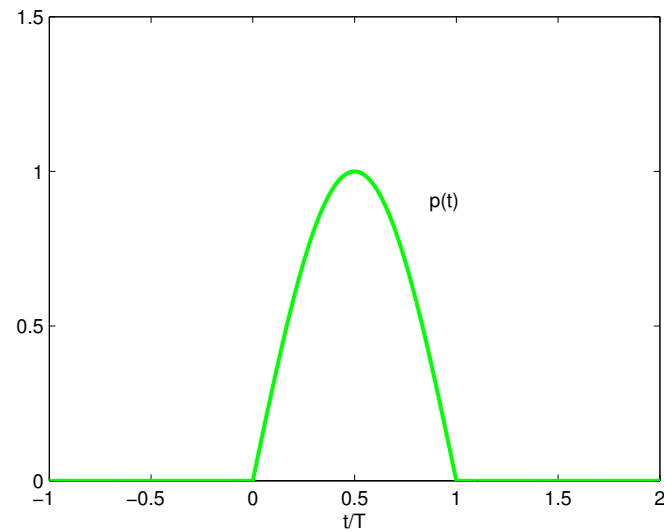
Map to symbols $\{a_k\}$

$$\text{BPSK : } a_k = \begin{cases} 1, & u_k = 0 \\ -1, & u_k = 1 \end{cases}$$

$$\text{QPSK : } a_k = \begin{cases} 1, & u_{2k}u_{2k+1} = 00 \\ i, & u_{2k}u_{2k+1} = 01 \\ -1, & u_{2k}u_{2k+1} = 10 \\ -i, & u_{2k}u_{2k+1} = 11 \end{cases}$$

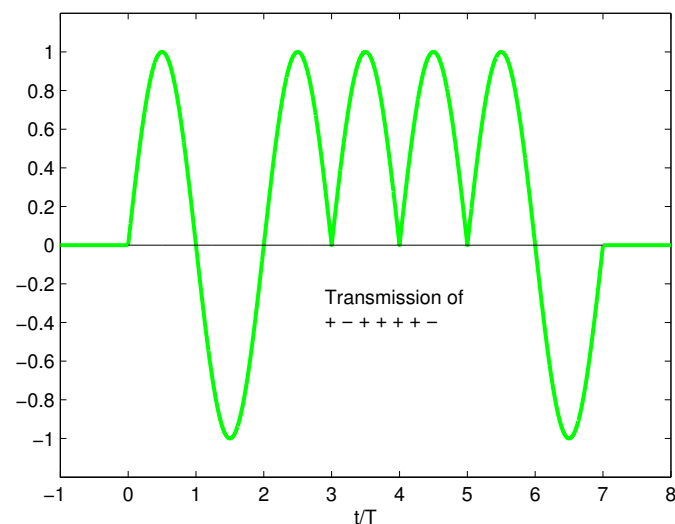
Basics of Digital Communications

Each symbol is carried by a **base pulse** $p(t)$ of length T , e.g. the **half-cycle sinus**



Basics of Digital Communications

So the transmission of bits 0 1 0 0 0 0 1 generates the pulse train $y(t)$



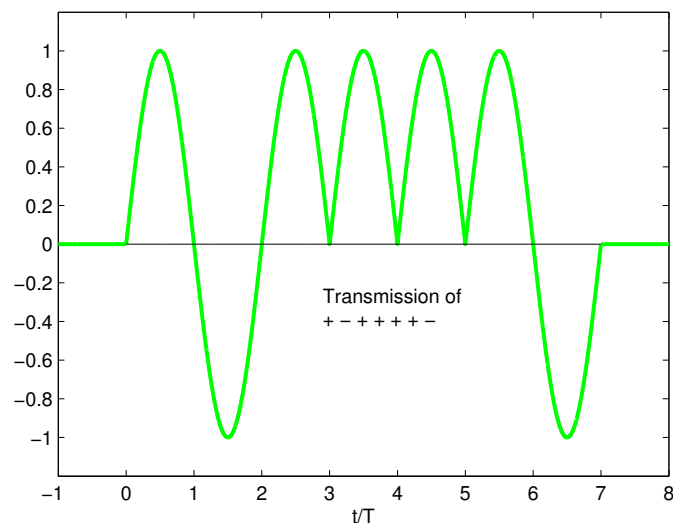
Mathematically we have

$$y(t) = \sum_k a_k p(t - kT_s)$$

Note that T_s is the symbol time while T is the duration of the pulse $p(t)$.

Basics of Digital Communications

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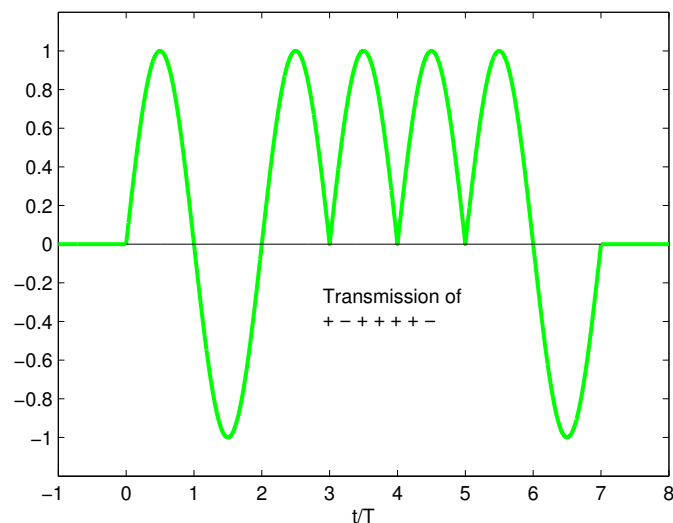
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How does T and T_s relate in this example?

Basics of Digital Communications

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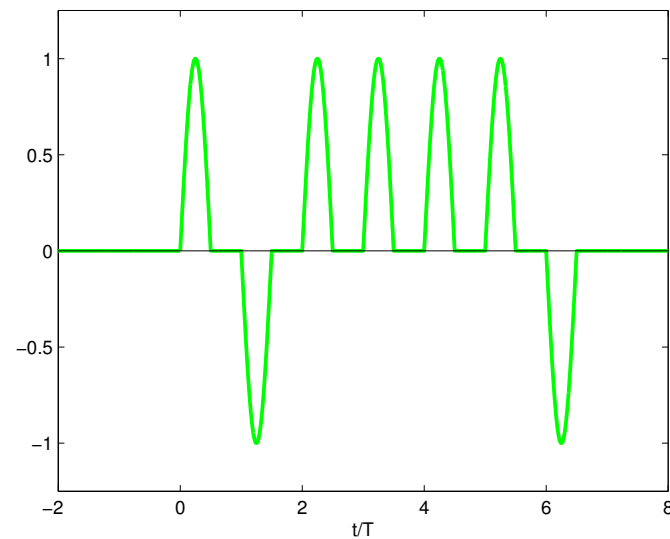
$$y(t) = \sum_k a_k p(t - kT_s)$$

Note that T_s is the symbol time while T is the duration of the pulse $p(t)$.

How does T and T_s relate in this example? $T = T_s$

Basics of Digital Communications

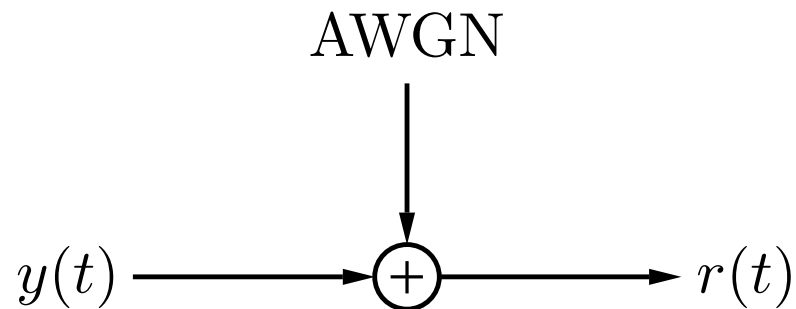
To avoid intersymbol interference one can use $T < T_s$



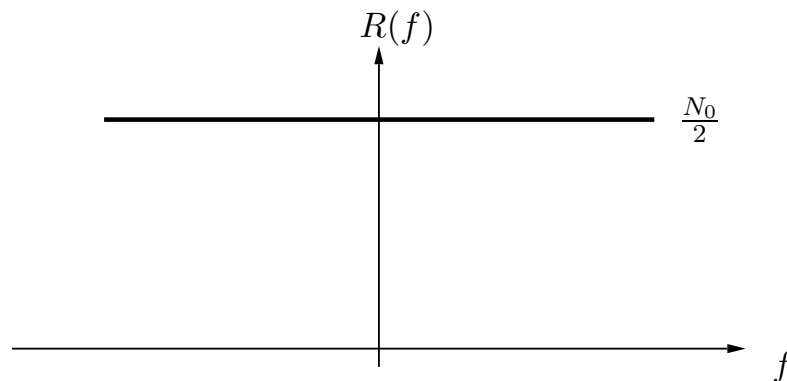
In this example we have $T = T_s/2$

Basics of Digital Communications

The channel model assumed in this review is a pure **AWGN** channel



Where the noise $n(t)$ satisfies $\mathcal{E}\{n^*(t)n(t+\tau)\} = \delta(\tau)N_0/2$; such a noise process must have power spectral density



Basics of Digital Communications

What does WGN look like?
Can we show an example?

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$$P = \int R(f)df$$

Basics of Digital Communications

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Consider the power of the process

$$P = \int R(f)df$$

$n(t)$ has infinite power!

Thus, not possible to show an example of WGN

Basics of Digital Communications

Explanation: Every signal we ever see in reality has been filtered by some low-pass filter.

Basics of Digital Communications

Mathematically, in what way should the receiver process the received signal $r(t)$.

In other words

$$\hat{\mathbf{a}} = \dots?$$

Basics of Digital Communications

Mathematically, in what way should the receiver process the received signal $r(t)$.

Maximum-likelihood detection is the answer!

$$\hat{\mathbf{a}} = \arg \max_{\mathbf{a}} \text{Prob}\{r(t)|\mathbf{a}\}$$

Basics of Digital Communications

Mathematically, in what way should the receiver process the received signal $r(t)$.

Maximum-likelihood is equivalent to minimum Euclidean distance detection

$$\hat{\mathbf{a}} = \arg \min_{\mathbf{a}} \int_{-\infty}^{\infty} |r(t) - \sum_k a_k p(t - kT_s)|^2 dt$$

Basics of Digital Communications

To decode the (**complex valued**) signal $r(t)$, we pass $r(t)$ through a **matched** filter $z(t)$

$$z(t) = p(-t)$$

For **symmetric pulses** $p(t)$, we get

$$z(t) = p(t)$$

Let

$$\begin{aligned} x(t) &= r(t) \star p(t) \\ &= \sum_k a_k g(t - kT_s) + \eta(t) \end{aligned}$$

where $\eta(t)$ is $n(t) \star p(t)$ and $g(t) = p(t) \star z(t)$. Take samples every T_s seconds: $x_k = x(kT_s)$. Then

$$x_k = E_p a_k + \eta_k$$

where η_k is a complex Gaussian random variable with variance $E_p N_0$, that is **$E_p N_0/2$ per dimension!**

Basics of Digital Communications

Energy computations and error probability:

The energy per transmitted **symbol** E_s is given by: $E_s = \underbrace{\int p^2(t)dt}_{E_p}$ while

the energy per transmitted **bit** is

$$E_b = \begin{cases} E_s, & \text{BPSK} \\ E_s/2, & \text{QPSK} \end{cases}$$

The physical minimum Euclidean distance is

$$D_{\min}^2 = \begin{cases} 4E_p, & \text{BPSK} \\ 2E_p, & \text{QPSK} \end{cases}$$

In both cases we end up with a normalized distance $d_{\min}^2 = 2$. The error probability is given by

$$P_e \approx Q\left(\sqrt{2\frac{E_b}{N_0}}\right)$$

Discrete-Time Implementations

In a computer-based package such as Matlab or C/C++, we cannot represent the signals $y(t)$ as continuous time signals. Hence we must work with sampled versions.

Let f_s be the **sample rate in samples/second** and N be the **number of samples per symbol**.

In PWC2, $f_s = 44100$ samples/second

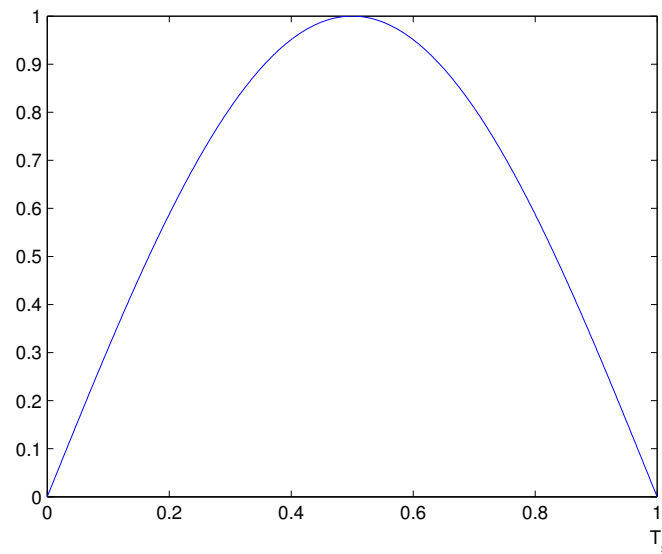
We get that $T_s = \frac{N}{f_s}$

The symbol rate becomes

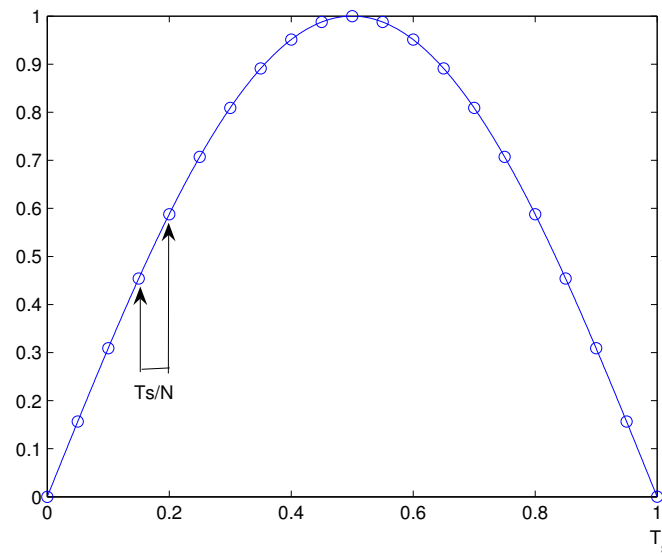
$$R_s = \frac{f_s}{N}$$

Discrete-Time Implementations

We must sample the base pulse $p(t)$.

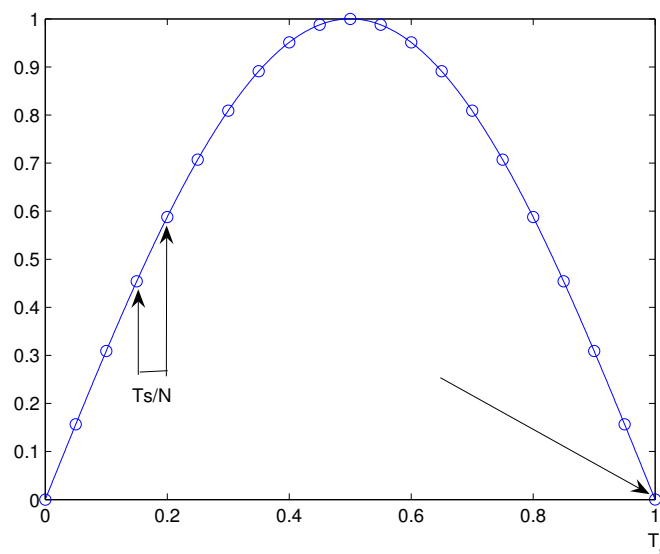


We must sample the base pulse $p(t)$. Assume a sample interval of T_s/N seconds



Discrete-Time Implementations

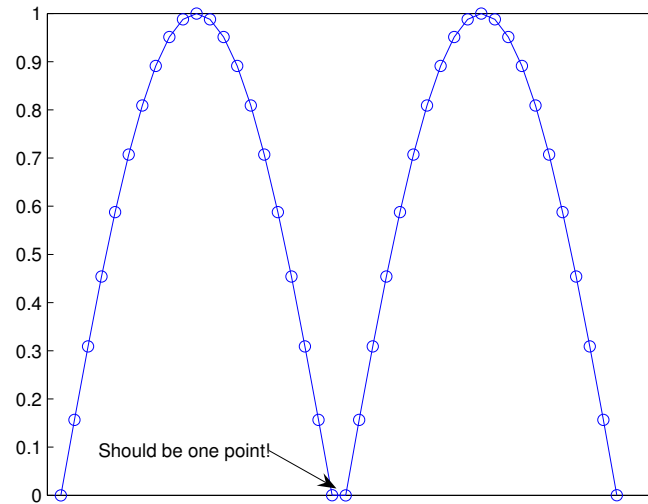
We must sample the base pulse $p(t)$. $N+1$ samples per symbol implies sample interval of T_s/N seconds



This is wrong!

Discrete-Time Implementations

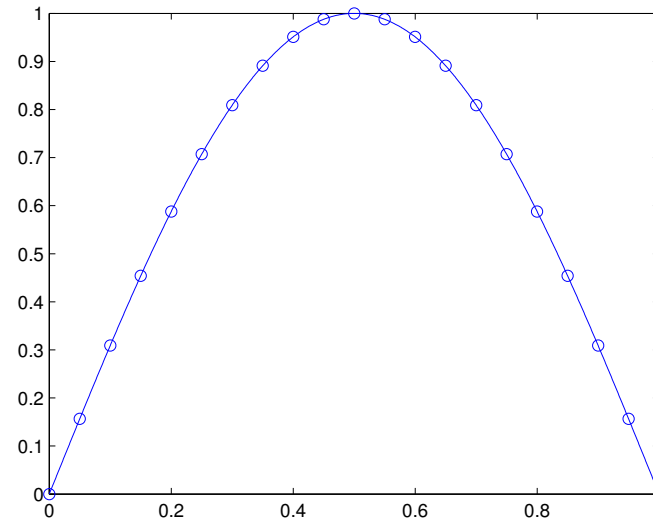
Explanation: Plot two consecutive pulses.



There should only be one point.

Discrete-Time Implementations

Correct sampling!

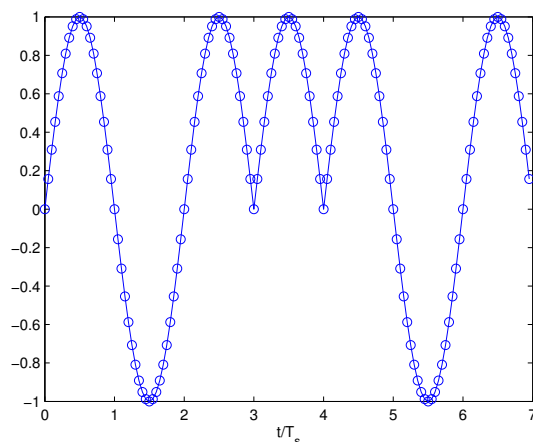


Represent the samples in a vector

$$\mathbf{p} = [0 \ 0.159 \ .309 \ \dots]$$

Discrete-Time Implementations

A sampled transmission signal of $+ - + + + - +$



Slightly harder mathematical representation. Let $\{b_k\}$ be a zero-padded version of $\{a_k\}$

$$\mathbf{b} = [a_1 \underbrace{00 \dots 0}_{N-1} \ a_2 \underbrace{00 \dots 0}_{N-1} \ a_3 \underbrace{00 \dots 0}_{N-1} \ a_4 \dots]$$

Then,

$$y_k = \sum_{\ell} b_{\ell} p_{k-\ell} \quad \text{or simply } \mathbf{y} = \mathbf{b} \star \mathbf{p}$$

Discrete-Time Implementations

Convolutions in discrete-time:

A convolution of $x(t)$ and $y(t)$ in continuous time is carried out as

$$\int x(\tau)y(t - \tau)d\tau \quad (1)$$

Let \mathbf{x} and \mathbf{y} be sampled version of $x(t)$ and $y(t)$; the sampling rate is f_s .
The discrete time version of (1) is

$$\frac{1}{f_s} \sum_{\ell} x_{\ell}y_{k-\ell}$$

The discrete time convolution must be scaled by the sampling rate!. $1/f_s$ works as $d\tau$ in (1).

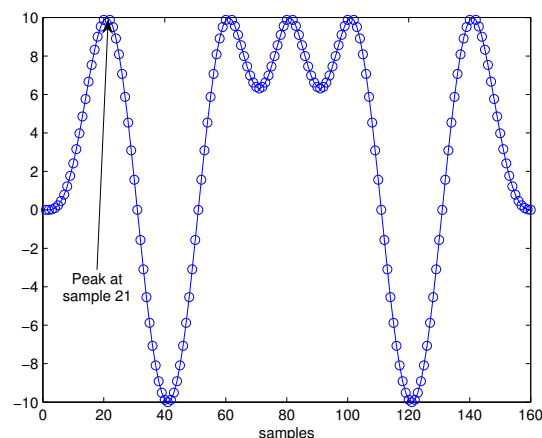
The energy of the pulse $p(t)$ must be approximated as

$$E_p = \frac{1}{f_s} \sum_k p_k^2$$

Discrete-Time Implementations

Matched filters in discrete-time:

The pulse train p should be filtered by a discrete-time matched filter. For symmetric pulses, we can take this matched filter as $z = p$ where p includes the last sample!, i.e. the length of p is $N + 1$. (This is however not crucial.) Then the output of the matched filter is ($N = 20$)



The number of samples in y is $N \times$ number of symbols and the length of the filter output is $N + N \times$ number of symbols. The peak occurs at samples $1 + kN$, $k = 1, 2, 3, \dots$

Discrete-Time Implementations

If there is a guard band ($T < T_s$), then the pulse is not symmetric and we can not take $z = p$. We must then use

$$z_k = p_{N+2-k}, \quad k = 1 \dots N + 1$$

It is still true that the peaks occur at samples $1 + kN$, $k = 1, 2, 3, \dots$

Discrete-Time Implementations

Implementation of discrete-time AWGN:

Until now we have constructed a modulation signal \mathbf{y} in discrete time. We now seek a noise vector \mathbf{n} to be added to \mathbf{y} that represents continuous time AWGN (that has inf power).

We have that both the real and the imaginary parts of the samples of

$$\eta(t) = n(t) \star z(t)$$

are zero-mean and have variance $E_p N_0/2$.

In discrete-time, a sample of the filtered noise process is given by

$$\eta_k = \frac{1}{f_s} \sum_{\ell} n_{\ell} z_{k-\ell}$$

Discrete-Time Implementations

Assume that the variance of each n_k is σ_n^2 . From probability theory it follows that η_k has variance $\sigma_n^2 \sum z_k^2 / f_s^2$.

Since

$$\sigma_n^2 \sum_k z_k^2 / f_s^2 = E_p \frac{N_0}{2}$$

we get that

$$\sigma^2 = E_p \frac{N_0}{2} \frac{f_s^2}{\sum_k z_k^2} = \frac{N_0}{2} f_s$$

Thus, The sampling rate affects the variance of the discrete time representation of continuous AWGN