

Background and motivation for yet another transform

For discrete-time systems, we have treated two transforms:

1. The z-transform

Good for solving difference equations Good for algebraic analysis of systems Limited insight from plotting

2. The discrete-time Fourier transform (DTFT) Excellent for understanding characteristics of systems Insights from plotting Relates to reality (bandwidth)



Consider the Processing module. It should process the signal according to the application at hand. (Can be more than a filter, for example an entire C/C++ program)



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These 6 numbers, are in the frequency domain represened by...





in the frequency domain represened by a continuous curve !



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The discrete Fourier Transform (DFT) in one sentence: A Fourier version of x(n) with 6 numbers

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Background and motivation for yet another transform

The discrete Fourier Transform (DFT) in one sentence: A Fourier version of x(n) with 6 numbers

Honey, I shrank the Discrete-time Fourier Transform

Easy to remember: Less words, less numbers

So, we have developed a transformation that we can only use to "look at on a piece of paper", not one that can be used by a computer.















Warning: This is not yet the DFT. It is only samples of the DTFT



If samples are representing X(f), then we should be able to get x(n) back from them

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Compact notation.



$$=\sum_{m=-\infty}^{\infty}x(n-mN)$$
























































Obvious Question Why would anybody want to compute a DFT of less Result length than that of $x(n)$?
if the length of $x(n)$ is N, then
$x_{\text{IDFT}}(n) \equiv x(n)$ and $X_{\text{DFT}}(k) \equiv X(f \mid f = k/N)$

Less obvious answer

Signals with infinite time-support requires this



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Signals with infinite time-support requires this

 $x(n) = a^n \cdot u(n)$

We have two options, a computer has one 1. Compute the N-point DFT according to

 $X_{\rm DFT}(k) = \sum_{n=0}^{N-1} x(n) e^{-j2\pi \cdot \frac{k}{N} \cdot n} \qquad \text{for } k = 0, 1, \dots, N-1$

2. Compute X(f) , and sample it to obtain X(k)

(computer cannot do 2 if it just encounters a signal. 2 requires us to know its math formula)

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Inversion of 1: we know the result $a^n \cdot u(n)$ n = 0, 1, ..., N-1

Let us invert 1&2, and compare 1. Compute the N-point DFT as

$$\frac{1}{N} \cdot \sum_{k=0}^{N-1} X(k) \mathrm{e}^{\mathrm{j} 2\pi \cdot \frac{n}{N} \cdot k}$$

 $X_{\rm DFT}(k) = \sum_{n=0}^{N-1} x(n) e^{-j2\pi \cdot \frac{k}{N} \cdot n} \qquad \text{for } k = 0, 1, \dots, N-1$

2. Compute ${\boldsymbol X}(f)$, and sample it to obtain ${\boldsymbol X}(k)$

Signals with infinite time-support requires this

$$x(n) = a^n \cdot u(n)$$

Inversion of 2: we know the result $\sum x(n-mN) \quad n=0,1,\ldots,N-1$ $\frac{1}{N} \cdot \sum_{k=0}^{N-1} X(k) \mathrm{e}^{\mathrm{j} 2\pi \cdot \frac{n}{N} \cdot k}$ $m \equiv -\infty$ Let us invert 1&2, and compare 1. Compute the N-point DFT as $X_{\rm DFT}(k) = \sum x(n) e^{-j2\pi \cdot \frac{k}{N} \cdot n}$ for k = 0, 1, ..., N - 12. Compute X(f), and sample it to obtain X(k)

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Inversion of 2: we know the result

$$\sum_{m=-\infty} x(n-mN) \quad n=0,1,\ldots,N-1$$



Signals with infinite time-support requires this

$$x(n) = a^n \cdot u(n)$$



Computational complexity

DFT defined as

$$X_{\text{DFT}}(k) = \sum_{n=0}^{N-1} x(n) e^{-j2\pi \cdot \frac{k}{N} \cdot n}$$

for
$$k = 0, 1, ..., N - 1$$

Number of operations needed:

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Each value requires N multiplications x(

$$x(n) \cdot \mathrm{e}^{-j2\pi kn/N}$$

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Number of operations needed:

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m DFT}(k)$ to be computed

Each value requires N multiplications

$$x(n) \cdot \mathrm{e}^{-j2\pi kn/N}$$

Total complexity N²








Computational complexityFFT not included in course, but good to know aboutTest in MatlabFast Fourier transform (FFT)If N=2k , then N log2(N) complexity to compute $X_{DFT}(k) = \sum_{n=0}^{N-1} x(n)e^{-j2\pi \cdot \frac{k}{N} \cdot n}$ for $k = 0, 1, \dots, N-1$

Made possible by some algebraic manipulations and tricks.

Cooley and Tukey 1965

Method known to, and used by, Gauss in 1805

Computational complexity FFT not included in course, but good to know about

Test in Matlab

Fast Fourier transform (FFT)

If $N=2^k$, then $N \log_2(N)$ complexity to compute

$$X_{\rm DFT}(k) = \sum_{n=0}^{N-1} x(n) e^{-j2\pi \cdot \frac{k}{N} \cdot n} \qquad \text{for } k = 0, 1, \dots, N-1$$

Made possible by some algebraic manipulations and tricks.

The importance of the FFT cannot be underestimated. WIFI and 4G, etc could not been implemented without the FFT

For a computer,

- 1. It can avoid the continuous DTFT
- 2. It can compute the DFT extremely fast

Properties

For DTFTs, we have

$$x(n) \star y(n) \leftrightarrow X(f)Y(f)$$

 $x(n) \leftrightarrow X(f) \qquad x(n-n_0) \leftrightarrow X(f)e^{-i2\pi f n_0}$

Still true ? I.e.

$$x(n) \star y(n) \leftrightarrow X(k)Y(k)$$

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Still true ? I.e. Assume length N sequences

$$\begin{aligned} x(n) \star y(n) &\leftrightarrow X(k)Y(k) \\ x(n) &\leftrightarrow X(f) \qquad x(n-n_0) \leftrightarrow X(k) e^{-i2\pi k n_0/N} \end{aligned}$$

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Assume length N sequences. Follows that DFTs also length N But this is length 2N-1. So its DFT must be length 2N-1

$$\begin{array}{c} -x(n) \leftrightarrow y(n) \leftrightarrow X(k)Y(k) \\ x(n) \leftrightarrow X(f) \qquad x(n-n_0) \leftrightarrow X(k) e^{-i2\pi k n_0/N} \end{array}$$

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Assume length N. Ex {1 2 3 4}

$$x(n) \leftrightarrow X(f)$$
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Also length N. Becomes {10 2+2i 2 2-2i}

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Assume length N. Ex {1 2 3 4} Length N+n_0. Ex {0 1 2 3 4} $x(n) \leftrightarrow X(f)$ $x(n - n_0) \leftrightarrow X(k) e^{-i2\pi k n_0/N}$

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Still length No

Makes no sense...

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For DFTs, we have

$$x_1(n) \otimes x_2(n) \leftrightarrow X(k)Y(k)$$
$$x(n - n_0 \mod N) \leftrightarrow X(k)e^{-i2\pi k n_0/N}$$
where $x_1(n) \otimes x_2(n) = \sum_{k=0}^{N-1} x_1(k)x_2(n - k \mod N)$

Circular convolution





In WIFI, 4G, 5G etc etc, we give up 5-10% of the possible data rate in order to turn channel convolution into a circular convolution. Then we can use the FFT and save complexity

Example

Circular convolution

 $x(n) = \{ \underline{1} \ 2 \ 3 \ 4 \}$ $h(n) = \{ \underline{2} \ 2 \ 1 \ 1 \}$ $y_C(n) = x(n) \otimes h(n)$

Example	Metl	thod 1
Circular convolution	h(k)	$1 1 2 \underline{2} \rightarrow$
$x(n) = \{\underline{1} \ 2 \ 3 \ 4\}$	x(k)	<u>2 3 4 1</u> 2 3 4 1 2 3
$h(n) = \{2 \ 2 \ 1 \ 1\}$	$y_C(k)$	<u>15</u> 13 15 17
$y_C(n) = x(n) \otimes h(n)$		Method 2
		<u>2 3 4 1</u> 2 3 4
		$\begin{array}{c ccccccccccccccccccccccccccccccccccc$
Method 3		2 4 6 8 2 4 6 8
$\begin{array}{cccccccccccccccccccccccccccccccccccc$	1 2 3	$\begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 1 \\ 2 \\ 3 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1$
$y_L(k)$ <u>2</u> 6 11 17 13 7 4 0		$\begin{bmatrix} - & - & - \\ - & - & - \\ - & - & - \\ - & - &$

Example	Method 1
Circular convolution	$h(k)$ 1 1 2 $\underline{2} \rightarrow$
Mathad 1	x(k) 2 3 4 <u>1</u> 2 3 4 1 2 3
>> x = $[1 \ 2 \ 3 \ 4];$ >> h = $[2 \ 2 \ 1 \ 1];$	$y_C(k)$ <u>15</u> 13 15 17
>> yc = ifft(fft(x).*fft(h)) yc = 15 13 15 17	Method 2
	<u>2 3 4 1</u> 2 3 4
	$\underline{2} \qquad 4 \qquad 6 \qquad 8 \qquad \underline{2} \qquad 4 \qquad 6 \qquad 8 \qquad \underline{3} \qquad $
Method 3	2 4 6 8 2 4 6 8
$\begin{array}{c ccccccccccccccccccccccccccccccccccc$	
$\frac{y_L(k)}{2} = 6 11 17 13 7 4 0$	

Example

Linear convolution computed via DFTs

- Given: Two length N sequences, x(n), y(n)
- Task: Compute their linear convolution by using DFT and its inverse IDFT

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```
This fails
```

Example

Linear convolution computed via DFTs

Given: Two length N sequences, x(n), y(n)

Task: Compute their linear convolution by using DFT and its inverse IDFT

>> $x=[1 \ 2 \ 3 \ 4];$ This is the result, But >> y=[2 2 1 1]; >> vL=conv(x, y)not computed via DFT vL = 2 6 11 17 13 7 4 >> xp=[1 2 3 4 0 0 0 0]; >> yp=[2 2 1 1 0 0 0 0]; >> yL=ifft(fft(xp).*fft(yp)) yL = 2.0000 6.0000 11.0000 17.0000 13.0000 7.0000 4.0000 -0.0000

Still a circular convolution carried out, but due to zero-padding, it behaves linear.

More examples

```
x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \}
```

Compute 16-point DFT (N=16) $X(k), k = 0, 1, \dots, 15$

More examples

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Compute 16-point DFT (N=16) $X(k), k = 0, 1, \dots, 15$

Compute inverse transform

$$x_{\text{IDFT}}(n) = \frac{1}{N} \sum_{k=0}^{N-1} X^2(k) e^{i2\pi kn/N}$$

More examples

$x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \}$

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More examples

Repeat for $x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ \}$

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Compute 16-point DFT (N=16) $X(k), k = 0, 1, \cdots, 15$

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More examples

Repeat for $x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 1 \ \}$

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Compute 16-point DFT (N=16) $X(k), \ k=0, \ 1, \cdots, 15$

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```
x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ \}
```

Compute DFT (N=6)



More examples

$x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0\}$

Compute DFT (N=8)



More examples

$x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0\}$

Compute DFT (N=8)



More examples

 $x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0 \cdots \}$

Compute DFT (N=16)







 $x(n) = \{1 \ 1 \ 1 \ 1 \ 1 \ 1 \ 0 \ 0 \cdots \}$

Compute DFT (N=16)



What is this line?

```
DFT size larger-or-equal to the length of x(n)
```



As a matrix computation

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-i2\pi kn/N}$$

We can write a DFT as a matrix multiplication



 $\mathbf{X} = \mathbf{W}\mathbf{x}$
EITF75 Systems and Signals

As a matrix computation

$$X(k) = \sum_{n=0}^{N-1} x(n) e^{-i2\pi kn/N}$$

We can write a DFT as a matrix multiplication

$$\begin{bmatrix} X(0) \\ X(1) \\ \vdots \\ X(N-1) \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & e^{-i2\pi/N} & e^{-i2\pi 2/N} & \cdots & e^{-i2\pi(N-1)/N} \\ 1 & e^{-i2\pi 2/N} & e^{-i2\pi 2 \cdot 2/N} & \cdots & e^{-i2\pi 2 \cdot (N-1)/N} \\ \vdots & \vdots & \vdots & \vdots & \vdots \\ 1 & e^{-i2\pi(N-1)/N} & e^{-i2\pi(N-1) \cdot 2/N} & \cdots & e^{-i2\pi(N-1) \cdot (N-1)/N} \end{bmatrix} \begin{bmatrix} x(0) \\ x(1) \\ \vdots \\ x(N-1) \end{bmatrix}$$

 $\mathbf{X} = \mathbf{W}\mathbf{x}$

Therefore

$$\mathbf{x} = \mathbf{W}^{-1}\mathbf{X}$$

Easy to show that W is unitary (orthogonal, but for complex matrices) ${f W}^H {f W} = {f I}$