



LUND
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EITF35: Introduction to Structured VLSI Design

Part 3.1.1: FSMD

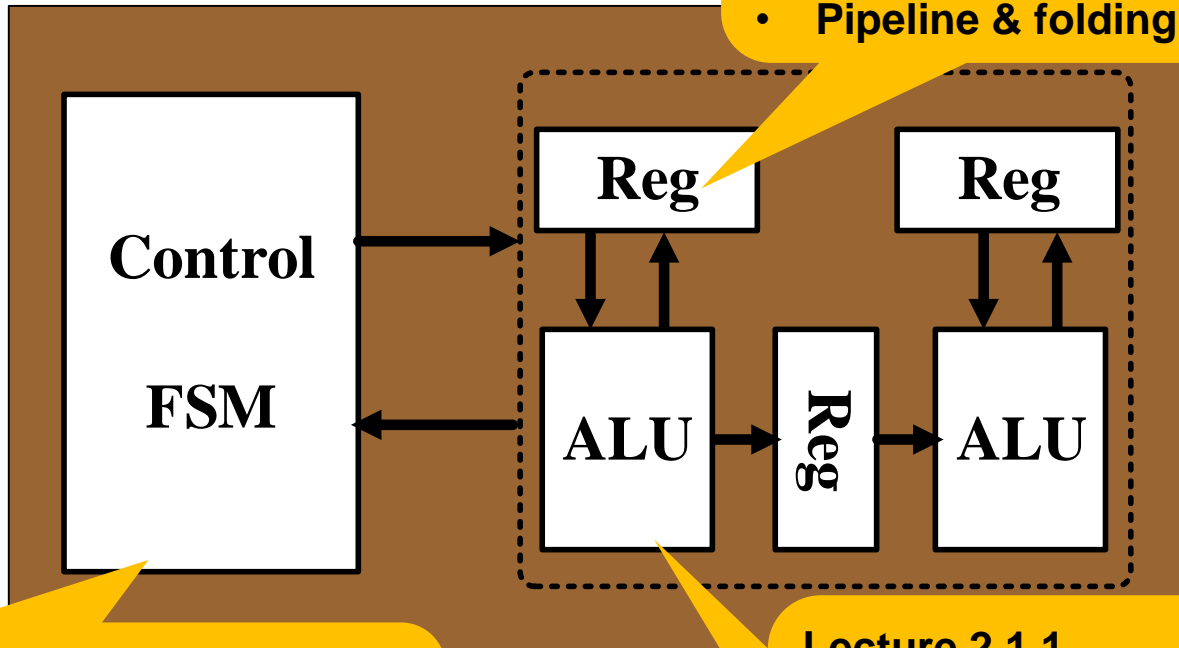
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Digital VLSI

Lecture 2.2.1

- Data storage
- Setup & hold timing
- No Latch
- Pipeline & folding



Lecture 1.2.1

- Current state -> next state
- Output logic

Lecture 2.1.1

- Order-irrelevant logic
- Complexity
- Propagation delay
- Resource sharing



Outline

- **FSMD Overview**
- **Algorithmic state machine with data-path (ASMD)**
- **FSMD design of a repetitive-addition multiplier**



Why FSMD? Start with algorithm

- **Task:** sums four elements of an array, divides the sum by 8 and rounds the result to the closest integer

```
size = 4
sum = 0;
for i in (0 to size-1) do {
    sum = sum + a(i);}
q = sum / 8;
r = sum rem 8;
if (r > 3) {
    q = q+1;}
outp = q;
```

*Algorithm: a sequence
steps of actions*

- **Two characteristics of an algorithm:**
 - Use of **variables**
e.g., sum, or $q = q + 1$
 - **Sequential execution**
e.g., sum must be finished before division



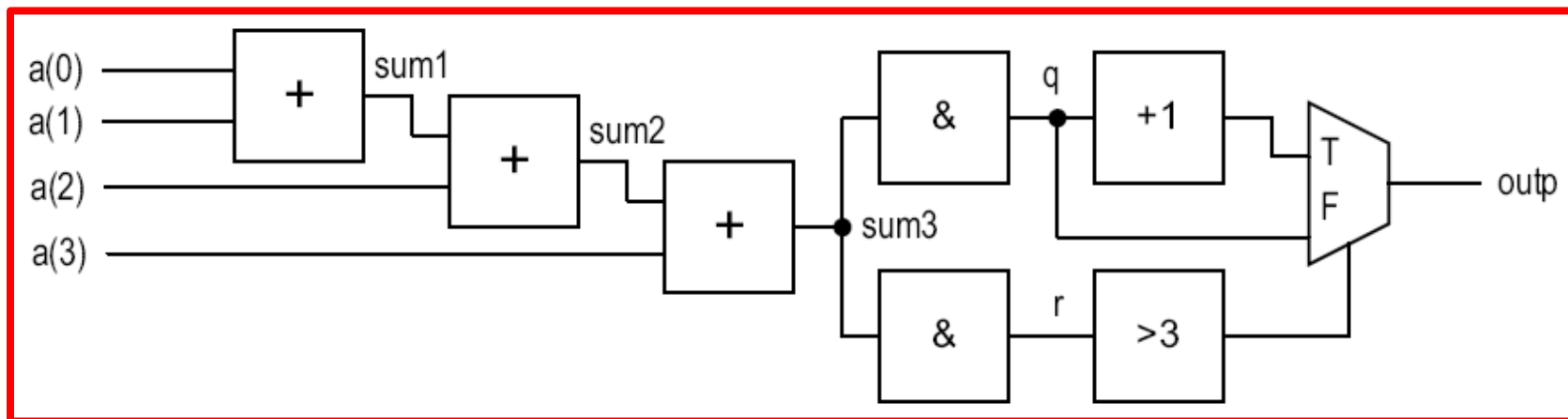
Converting algorithm to hardware

□ “Dataflow” implementation in VHDL

- Convert the algorithm in to **combinational circuit**

```
sum <= 0;  
sum0 <= a(0);  
sum1 <= sum0 + a(1);  
sum2 <= sum1 + a(2);  
sum3 <= sum2 + a(3);  
q <= "000" & sum3(8 downto 3);  
r <= "00000" & sum3(2 downto 0);  
outp <= q + 1 when (r > 3) else q;
```

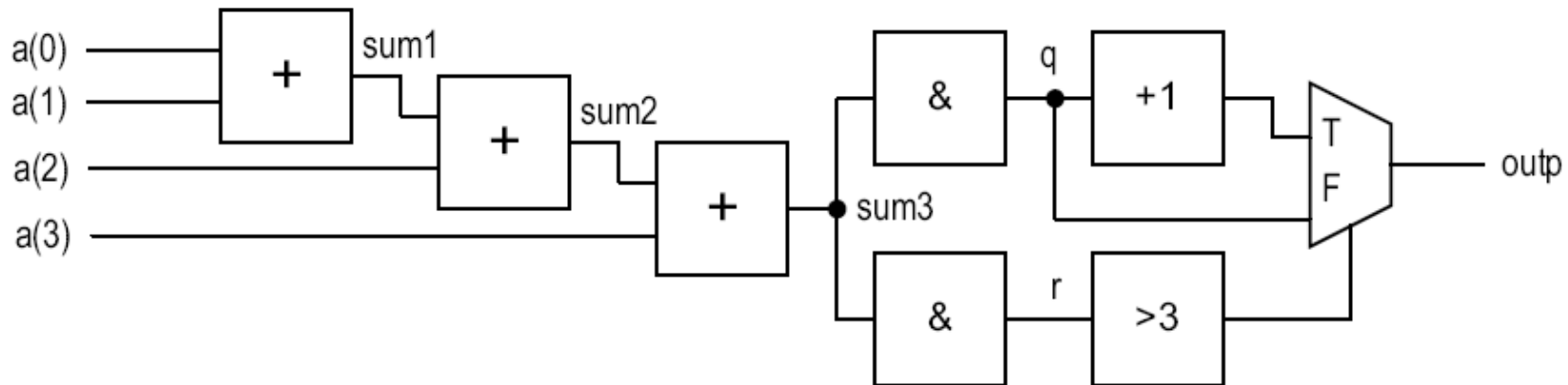
The "sequential" operations are represented by the data flow from left to right



Dataflow Implementation: Drawbacks

□ Problems with dataflow implementation:

- Can only be applied to simple trivial algorithm
- **Not flexible**
 - *What if size=10, 100, 1000 ...*
 - *or size = n, i.e., size is determined by an external input*
 - *or changing operation depending on instructions*



Alternatively?

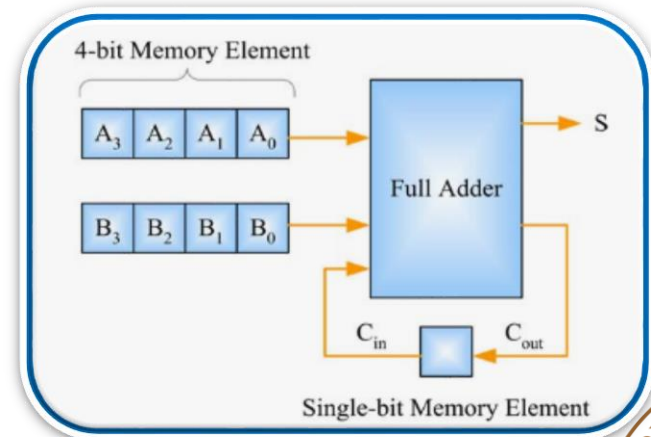
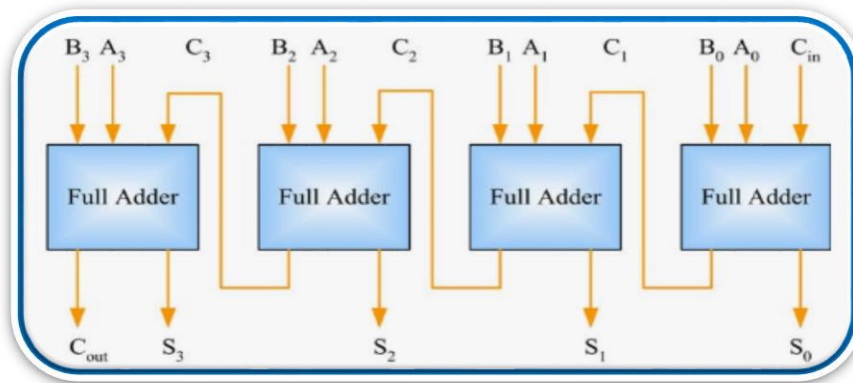
□ Hardware resembles the *variable* and *sequential* execution model

- Use **register** to store *intermediate data* and imitate *variable*
e.g. $\text{sum} = \text{sum} + a \Rightarrow \text{sum_reg} + a_reg \rightarrow \text{sum_reg}$

- Basic format of **RT operation**

$$r_{\text{dest}} \leftarrow f(r_{\text{src1}}, \dots, r_{\text{srcn}})$$

- Sequence of data manipulation and transfer among registers (**RTL**)

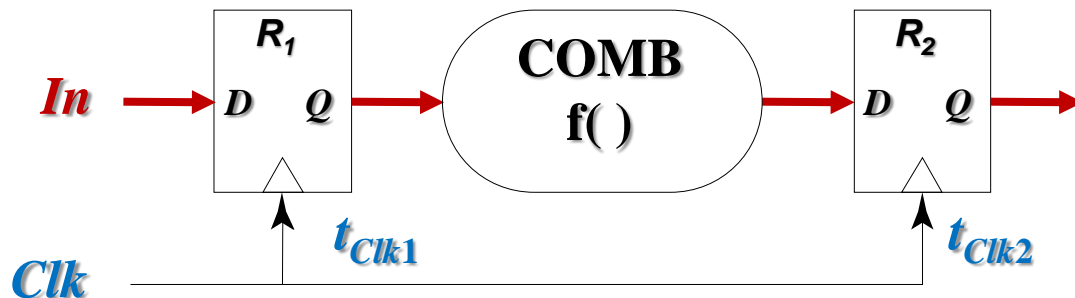


RT Operation: Timing

$$r_{\text{dest}} \leftarrow f(r_{\text{src1}}, r_{\text{src2}}, \dots, r_{\text{srcn}})$$

□ Timing:

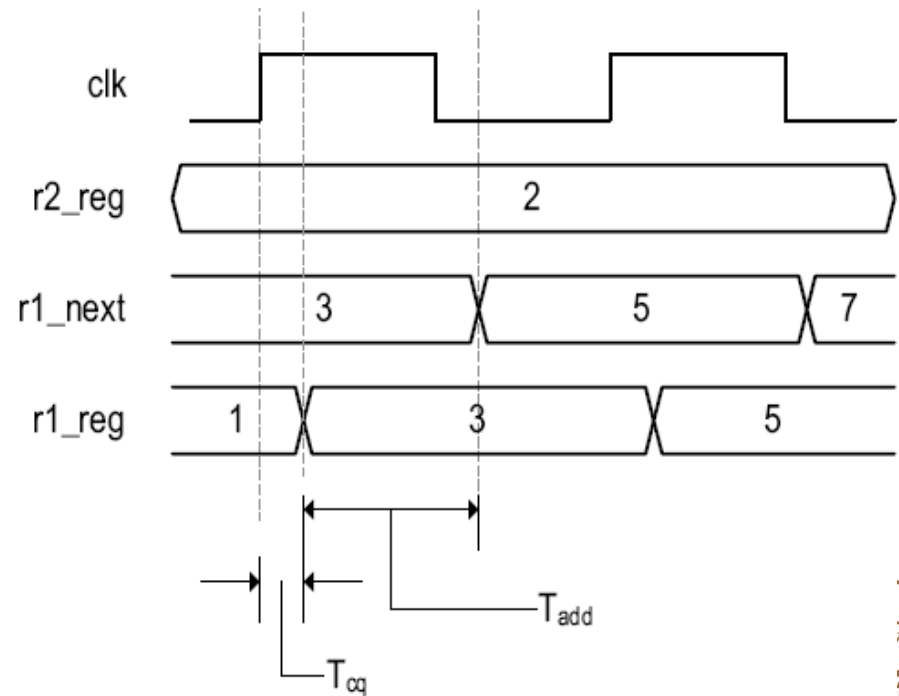
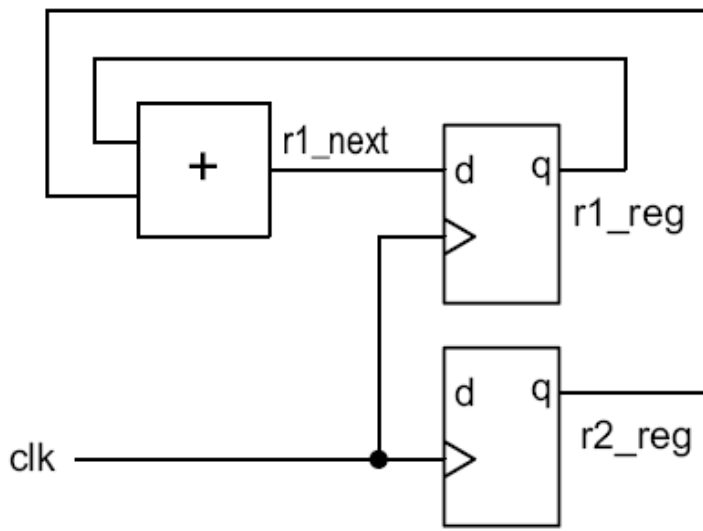
- **Hardware!** an explicit **clock** is embedded in an RT operation
- Rising edge of clk: outputs of source reg r_{src1} r_{src2} etc. are available
- The output are passed to a combinational circuit that performs $f()$
- At the **NEXT rising edge** of the clock, the result is stored into r_{dest}



Hardware Mapping of RT: Example 1

□ E.g. $r1 \leftarrow r1+r2$

- C1: $r1_next \leftarrow r1_reg + r2_reg$
- C2: $r1_reg \leftarrow r1_next$

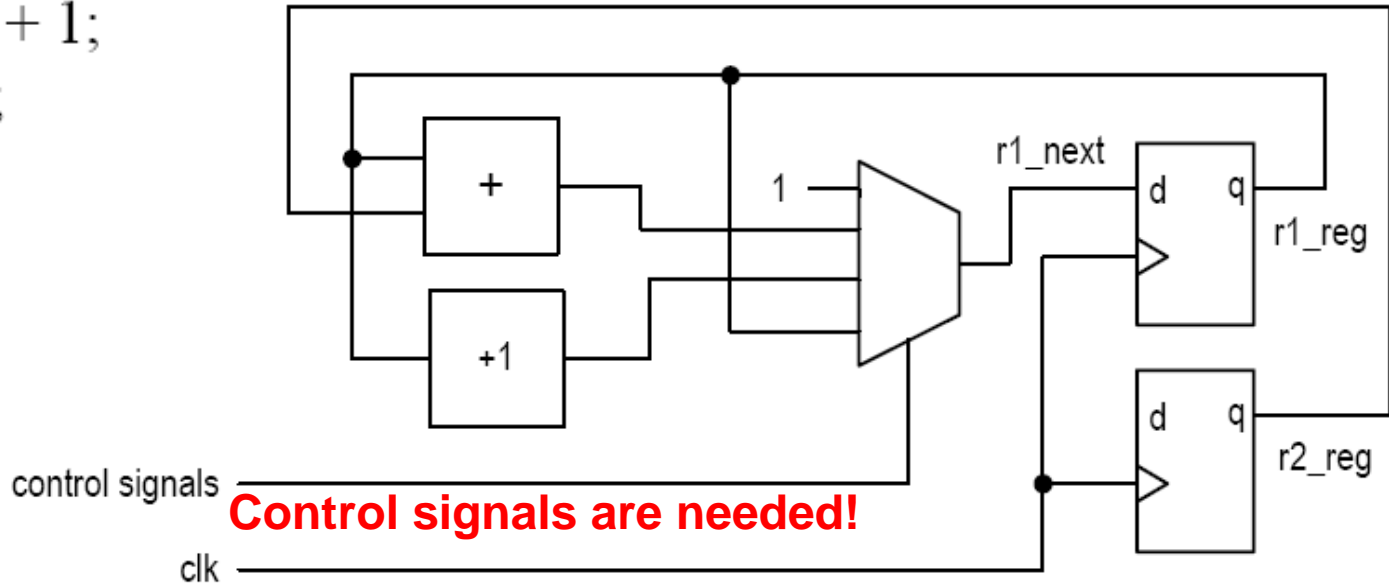


Hardware Mapping of RT: Example 2

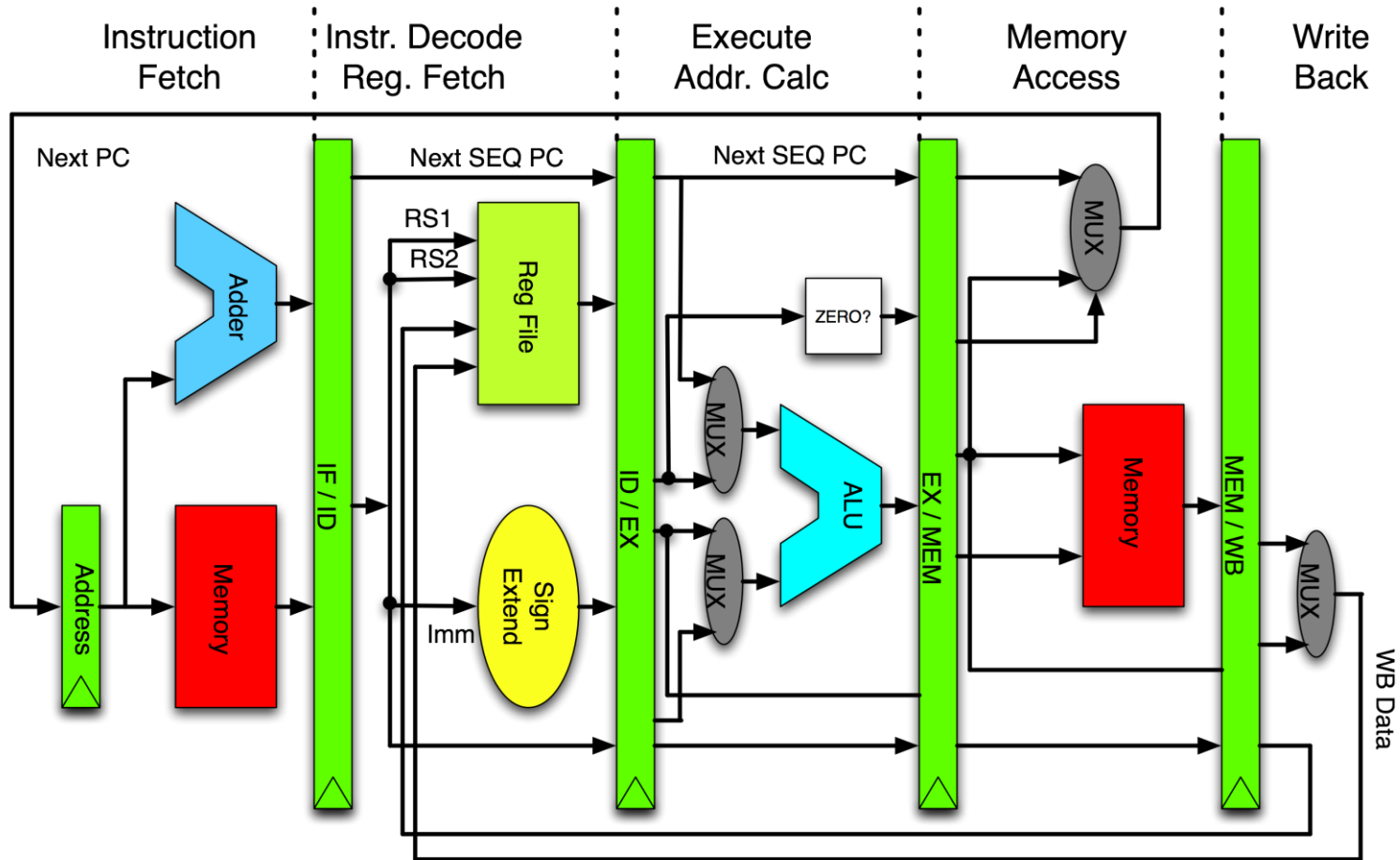
Multiple RT operations

How can we organize multiple operations on one register (in a time-multiplexing way)?

```
r1 ← 1;  
r1 ← r1 + r2;  
r1 ← r1 + 1;  
r1 ← r1;
```



MIPS pipeline



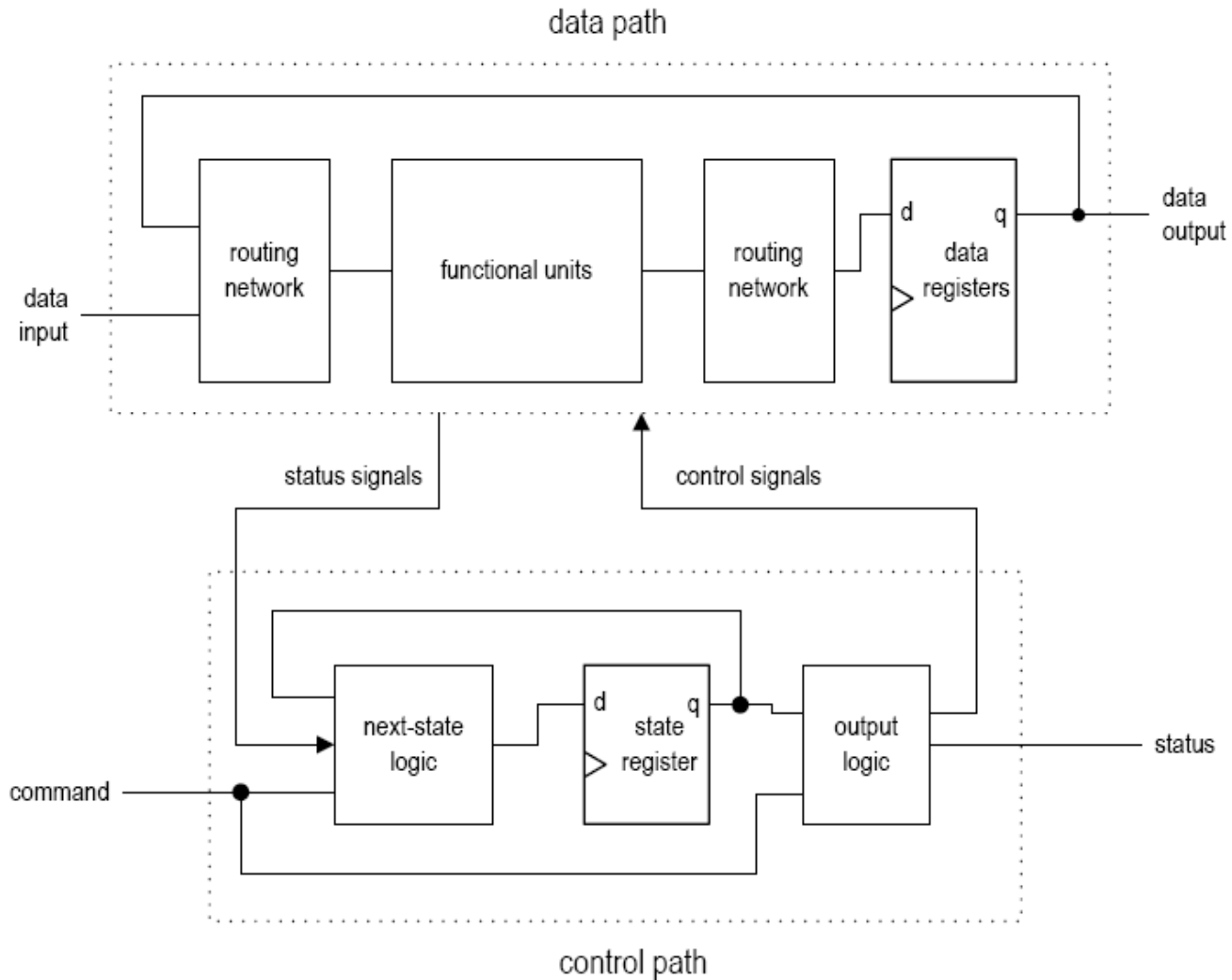
FSM as Control Path

□ FSMD: FSM with data path

- Use a **data path** to realize all the required RT *operations*
- Use a **control path (FSM)** to specify the *order* of RT operation



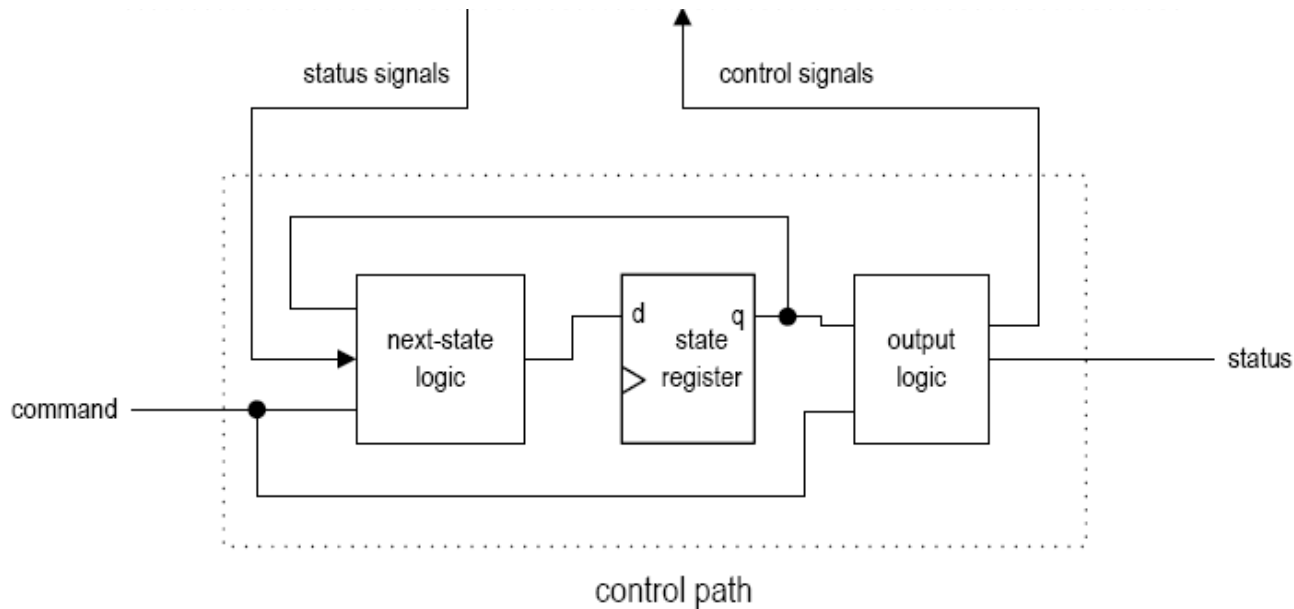
FSMD (FSM with Date Path)



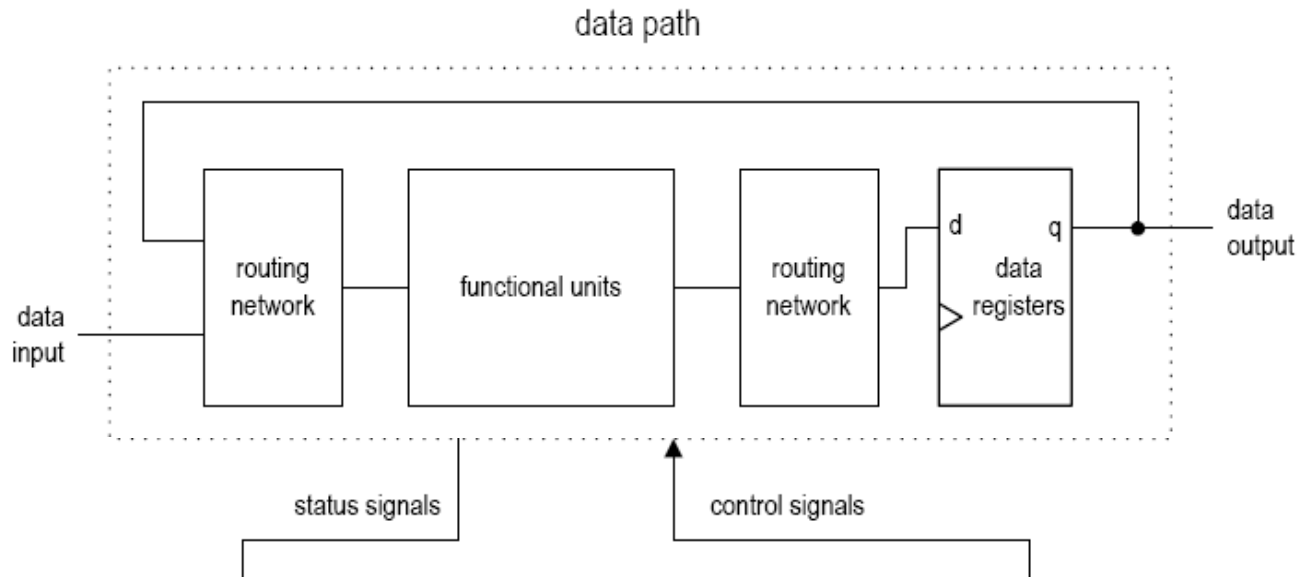
FSMD (FSM with Data Path)

□ Control Path: FSM

- **Command**: the external command signal to the FSMD
- **Internal status**: signal from the data path.
- **Control signal**: output, used to control data path operation.
- **External status**: output, used to indicate the status of the FSMD



FSMD (FSM with Date Path)



□ **Data Path:** perform all the required RT operations

- **Data registers:** store the intermediate results.
- **Functional units:** perform RT operations
- **Routing circuit:** connection, selection (multiplexers)



Outline

- Overview of FSMD
- **Algorithmic state machine with data-path (ASMD)**
- FSMD design of a repetitive-addition multiplier
- Timing analysis of FSMD



ASM (algorithmic state machine)

□ ASM (algorithmic state machine) chart

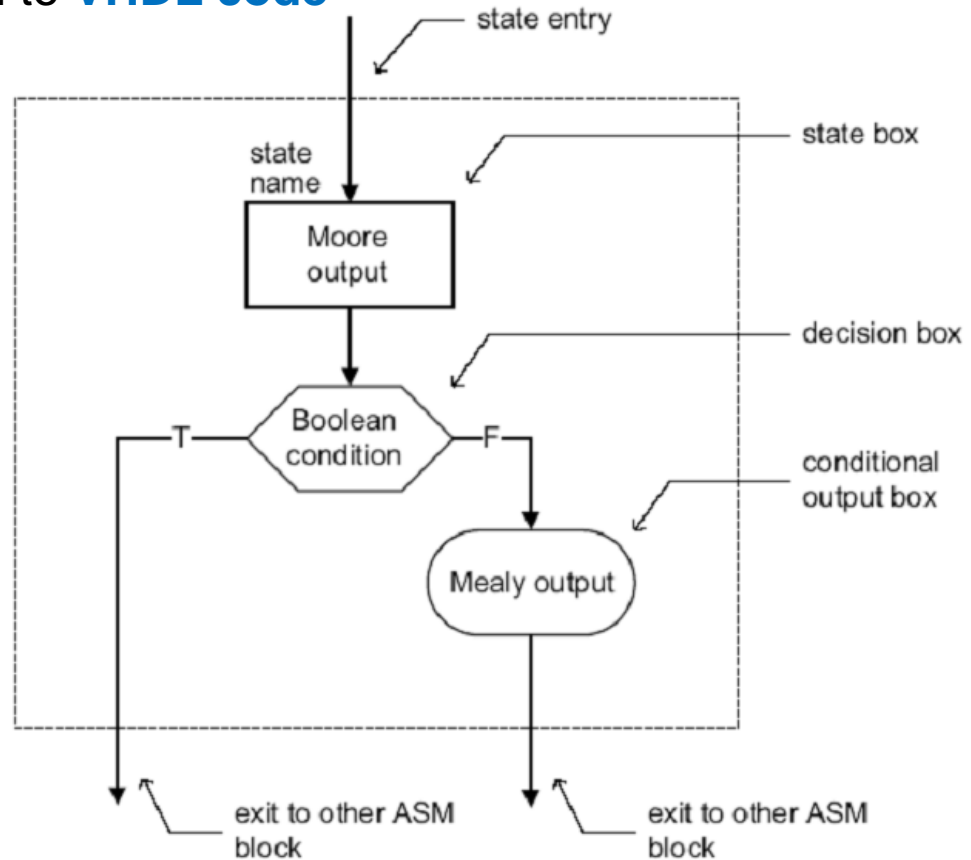
- **Flowchart-like** diagram, provide the same information as an FSM
- More **descriptive**, better for complex algorithm (both condition and uncondition operations)
- Can easily be transformed to **VHDL code**

An ASM chart is a network of **ASM blocks**

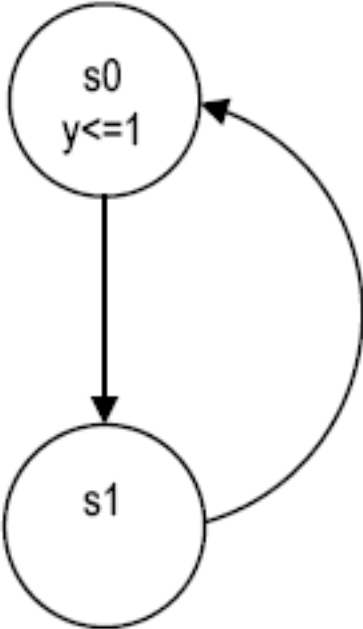
□ **One state box: FSM state**

□ **Decision boxes: with *T* or *F* exit path: next state logic**

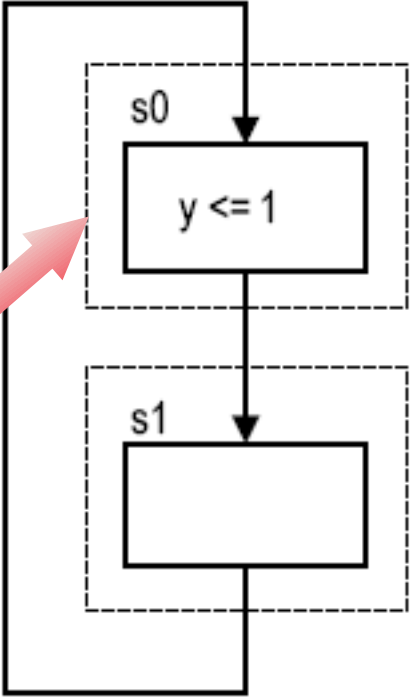
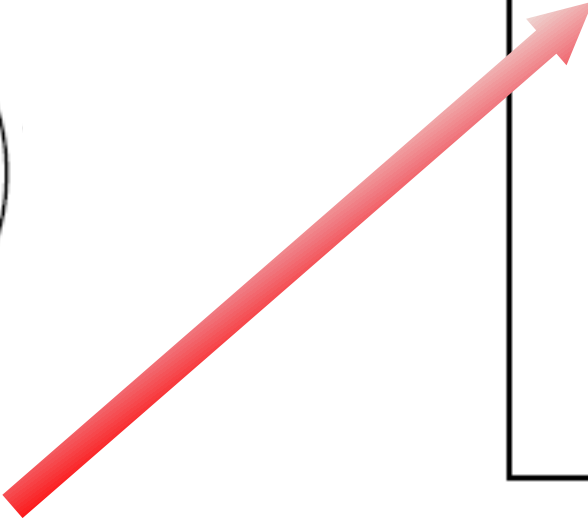
□ **Conditional output (operation) boxes: for Mealy output**



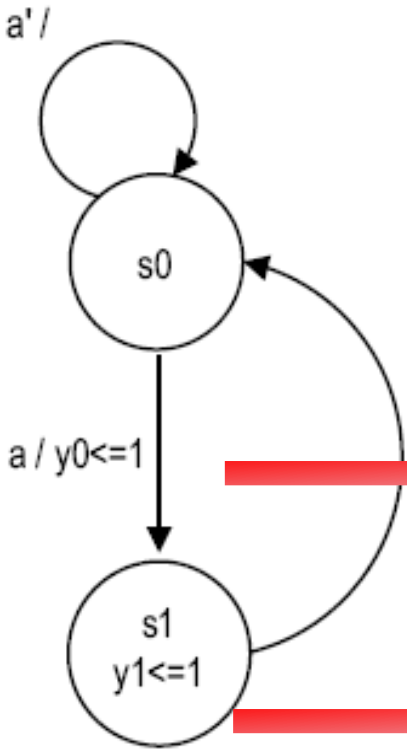
State Diagram and ASM Chart: Example 1



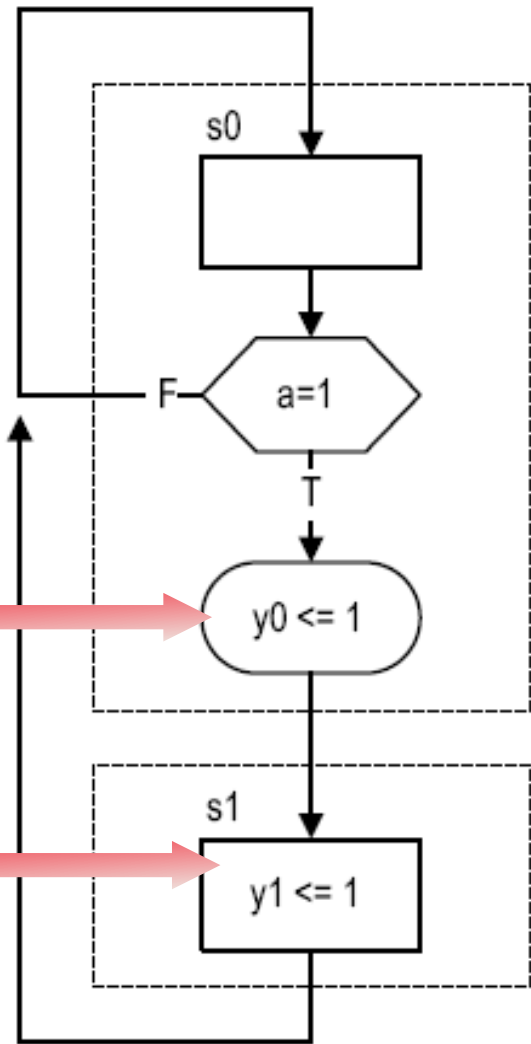
Moore FSM



State Diagram and ASM Chart: Example 2



Moore and Mealy



ASMD

ASMD:

- Extend ASM chart to incorporate **RT operations**
- RT operations are treated as another type of activity and be placed where the output signals are used

S0:

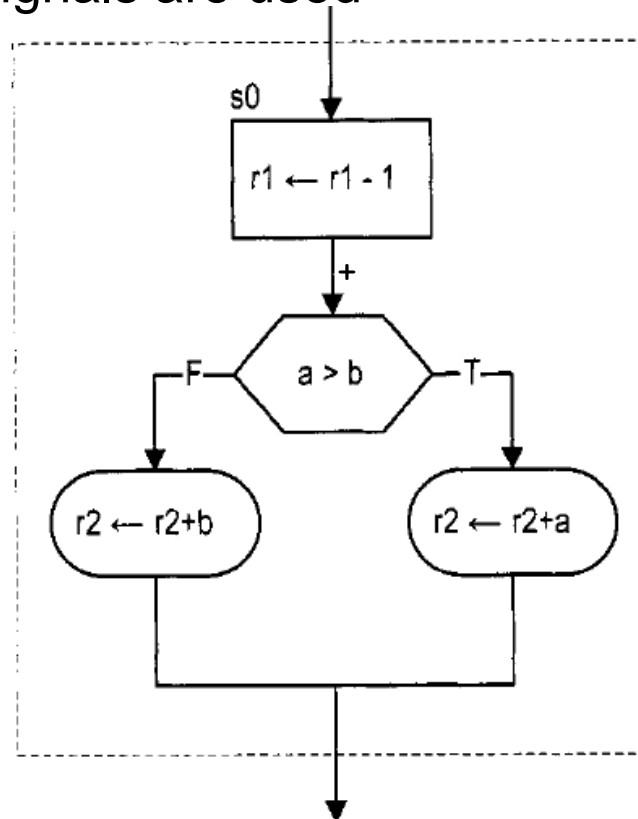
$r1 \leftarrow r1 + 1$

if $a > b$

$r2 \leftarrow r2 + a$

else

$r2 \leftarrow r2 + b$



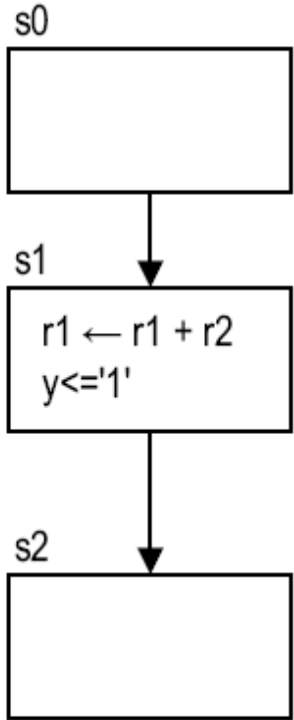
\leftarrow :RT operation, **infer a register**

= or \leftarrow :infer combinational logic

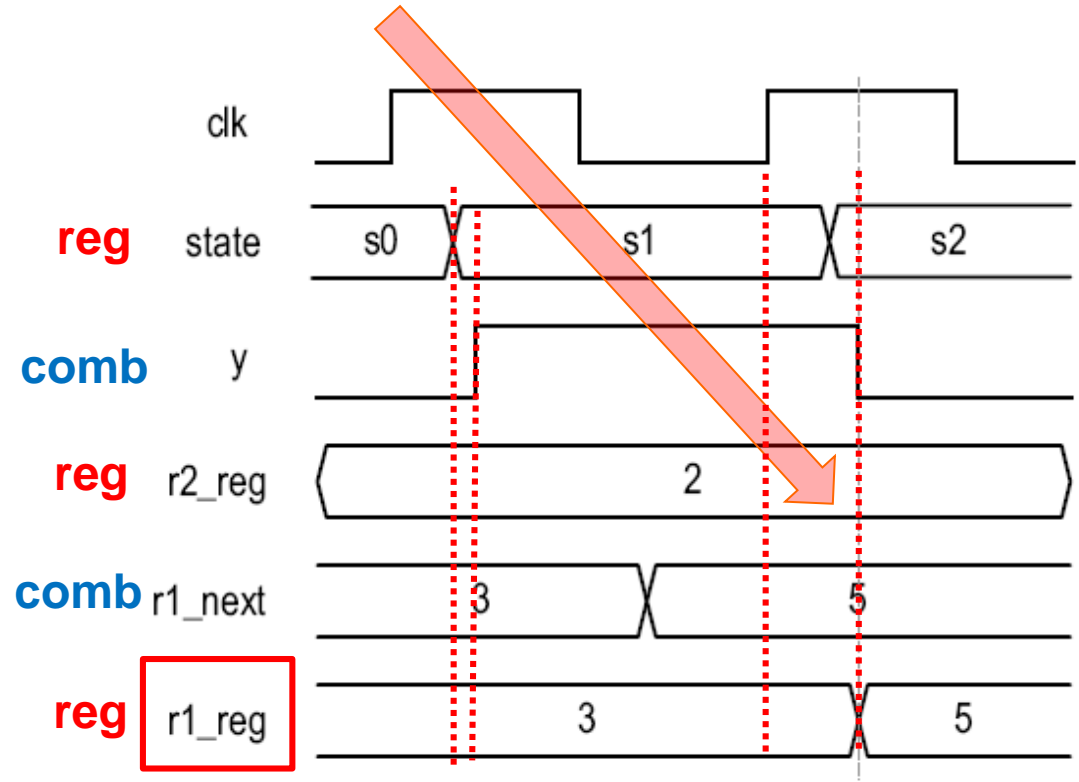


ASMD: Timing

R1_reg is updated at the **NEXT** clock tick



(a)



(c) Timing diagram

Suggestion: use meaningful names for your signals (direction_function_type)



Outline

- Overview of FSMD
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- **FSMD design of a repetitive-addition multiplier**
- Timing analysis of FSMD



Map Algorithm to FSMD

Example: Repetitive addition multiplier

□ Basic algorithm: $7*5 = 7+7+7+7+7$

```
if (a_in=0 or b_in=0) then {
    r = 0;}
else{
    a = a_in;
    n = b_in;
    r = 0;
    while (n != 0 ){
        r = r + a;
        n = n-1;}
}
return (r)
```

Pseudo code

```
if (a_in=0 or b_in=0) then {
    r = 0;}
else{
    a = a_in;
    n = b_in;
    r = 0;
op:   r = r + a;
      n = n-1;
      if (n = 0) then{
          goto stop;}
      else{
          goto op;}
}
stop: return (r);
```

ASMD-friendly code



ASMD Chart

Input:

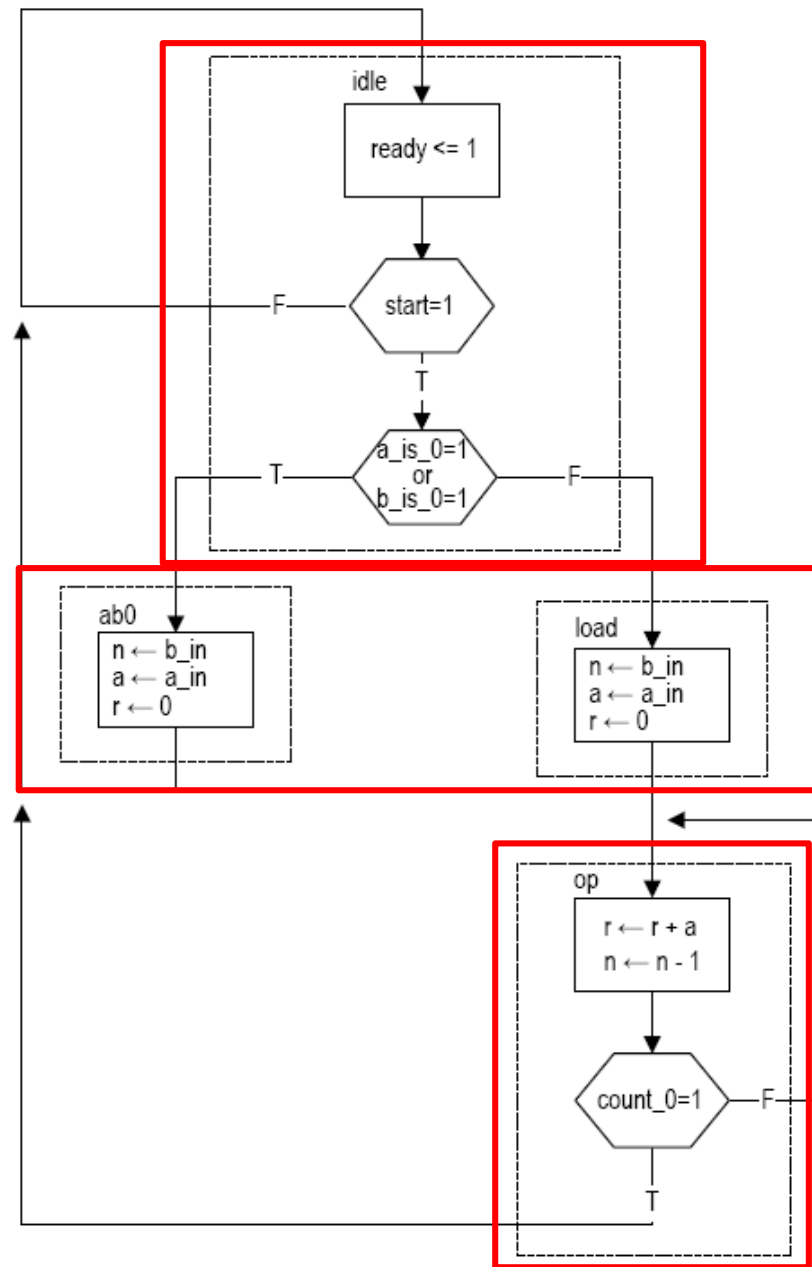
- a_in, b_in: 8-bit unsigned
- **clk, reset**
- start: command

Output:

- r: 16-bit unsigned
- ready: ready for new input

ASMD chart

- 3 registers (n,a,r)
- 4 states
- Data-path: RT operations
- FSM: state transition



Translate ASMD to Hardware



Construction of FSMD

□ Construction of the **data path**

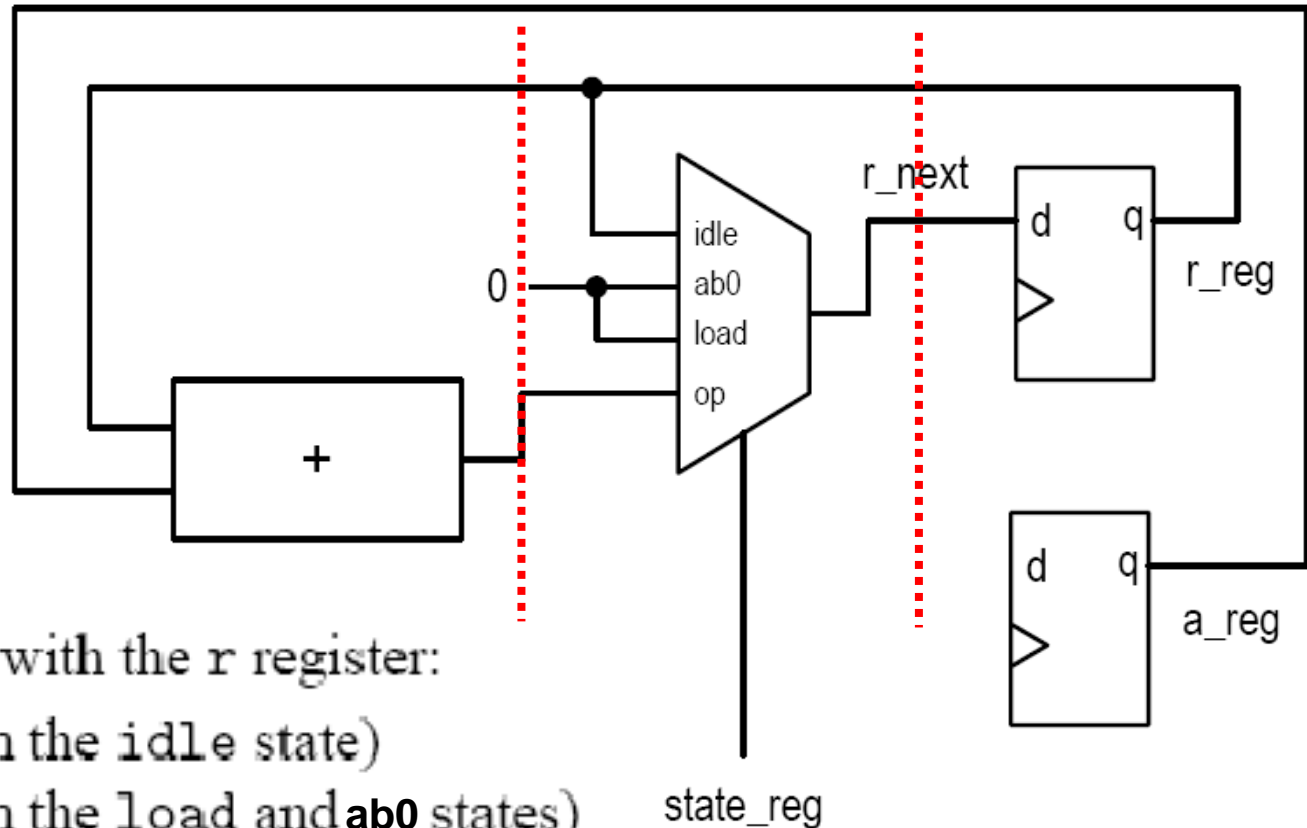
- List all possible RT operations
 - Group RT operation according to the **destination register**
 - Add combinational circuit/mux
- RT operations with the **r** register:
 - $r \leftarrow r$ (in the **idle** state)
 - $r \leftarrow 0$ (in the **load** and **ab0** states)
 - $r \leftarrow r + \mathbf{a}$ (in the **op** state)
 - RT operations with the **n** register:
 - $n \leftarrow n$ (in the **idle** state)
 - $n \leftarrow \mathbf{b_in}$ (in the **load** and **ab0** states)
 - $n \leftarrow n - 1$ (in the **op** state)
 - RT operations with the **a** register:
 - $\mathbf{a} \leftarrow \mathbf{a}$ (in the **idle** and **op** states)
 - $\mathbf{a} \leftarrow \mathbf{a_in}$ (in the **load** and **ab0** states)

Grouping RT Operations



Construction of the Date Path

Circuit associated with **r register**

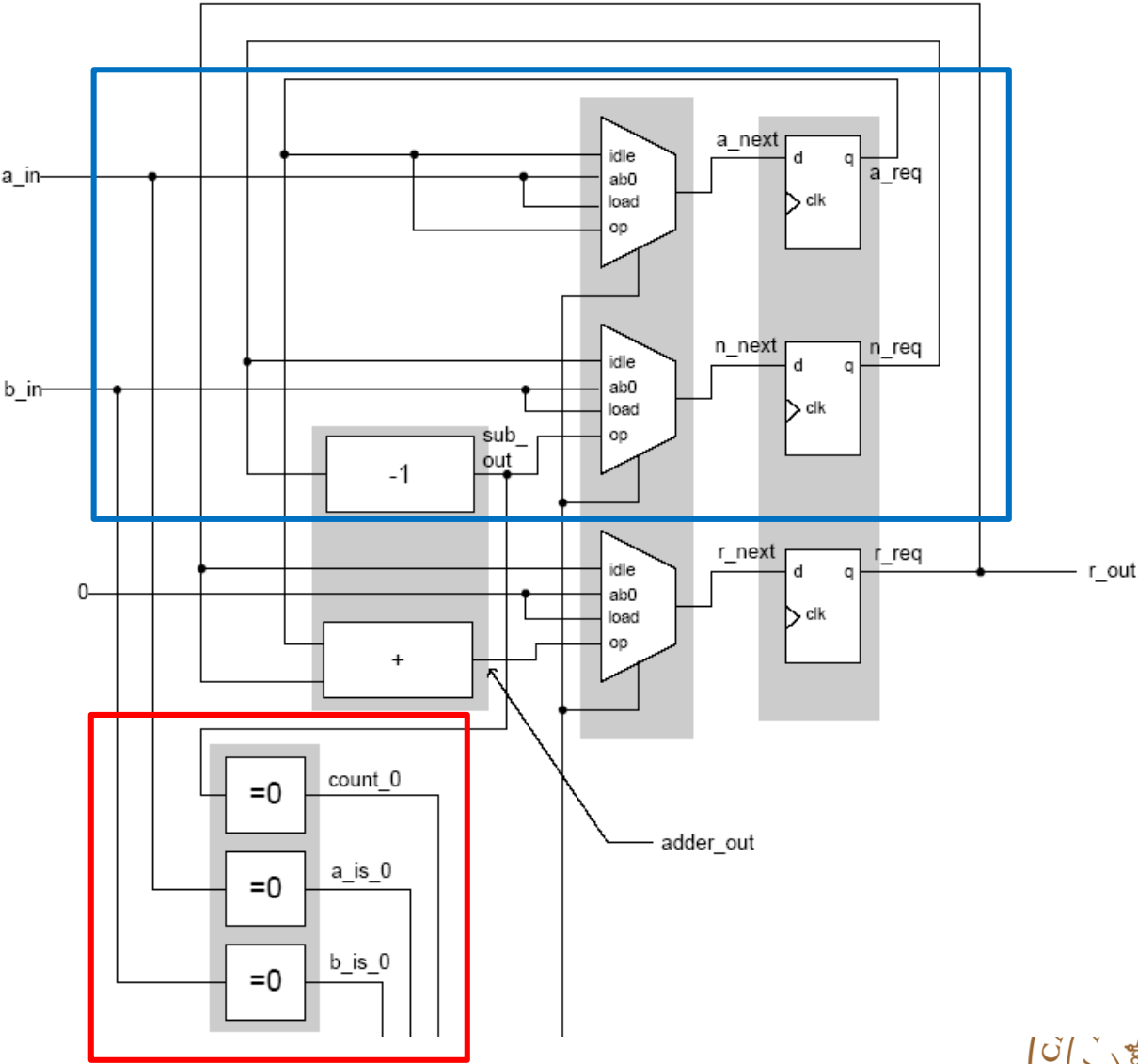


- RT operations with the **r** register:
 - $r \leftarrow r$ (in the **idle** state)
 - $r \leftarrow 0$ (in the **load** and **ab0** states)
 - $r \leftarrow r + a$ (in the **op** state)



Construction of the Date Path

Continue with
n-register
a-register
Add status circuits

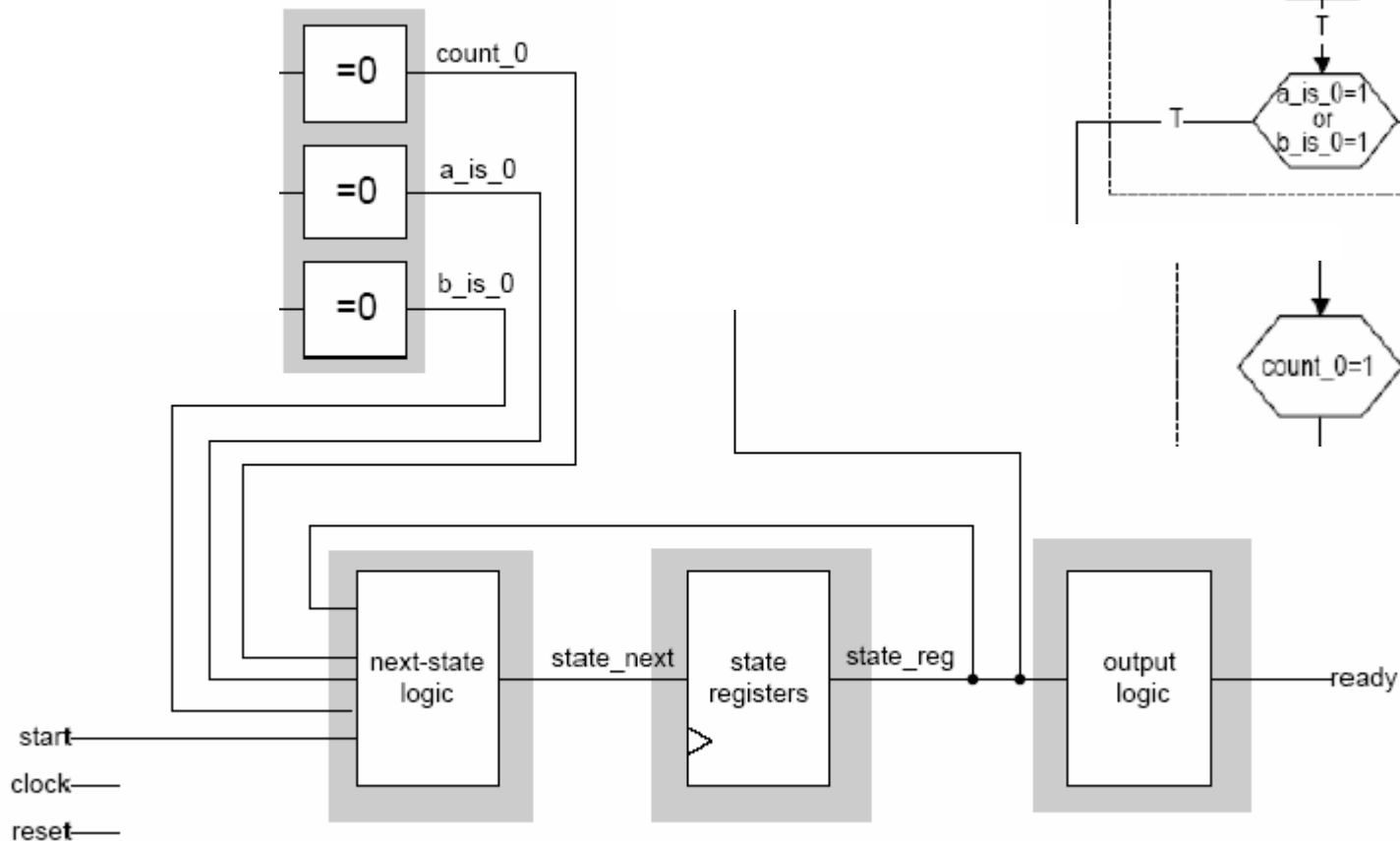


Construction of the Control Path

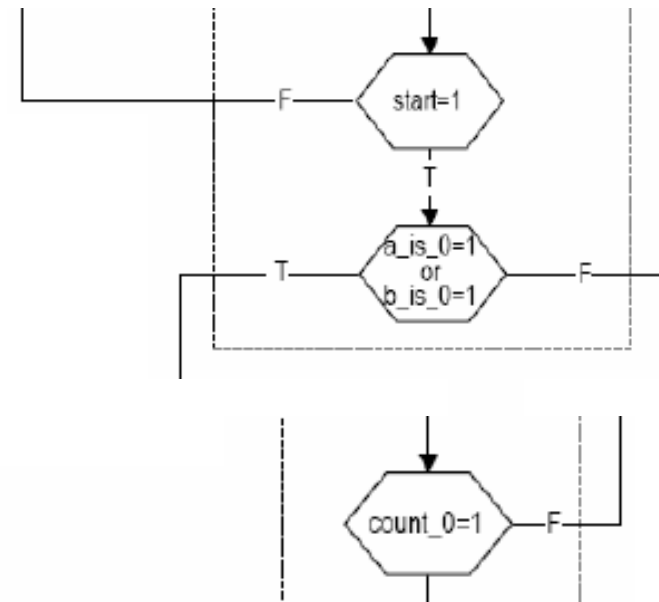
Input of FSM

- External: start, clock, reset
- Internal: decision box in ASMD

Output of FSM



Decision Box



VHDL Follow the Block Diagram

□ Entity

```
library ieee;
use ieee.std_logic_1164.all;
use ieee.numeric_std.all;

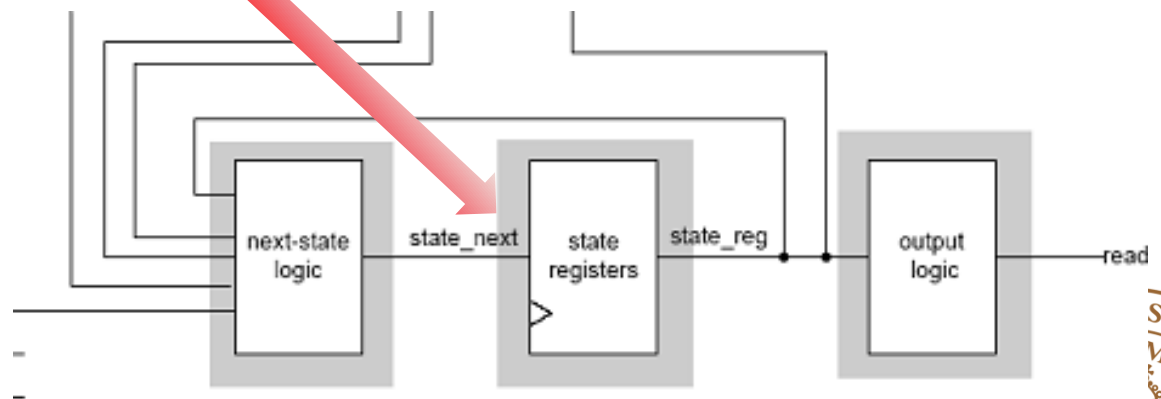
entity seq_mult is
  port (
    clk, reset: in std_logic;
    start: in std_logic;
    a_in, b_in: in std_logic_vector(7 downto 0);
    ready: out std_logic;
    r: out std_logic_vector(15 downto 0)
  );
end seq_mult;
```



FSM (state registers)

— *control path: state register*

```
process (clk, reset)
begin
  if reset='1' then
    state_reg <= idle;
  elsif (clk'event and clk='1') then
    state_reg <= state_next;
  end if;
end process;
```



FSM (next-state/output logic)

```
process(state_reg,start,a_is_0,b_is_0,count_0)
begin
```

```
  case state_reg is
```

```
    when idle =>
      if start='1' then
        if (a_is_0='1' or b_is_0='1') then
          state_next <= ab0;
        else
          state_next <= load;
        end if;
      else
        state_next <= idle;
      end if;
```

```
    when ab0 =>
      state_next <= idle;
```

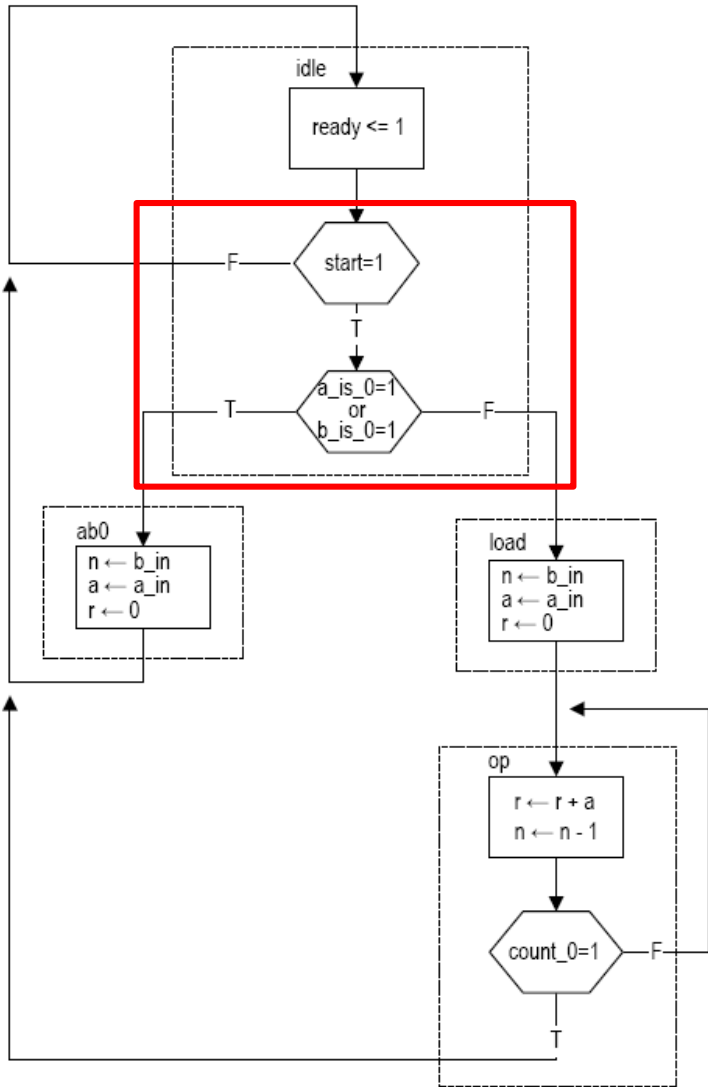
```
    when load =>
      state_next <= op;
```

```
    when op =>
      if count_0='1' then
        state_next <= idle;
      else
        state_next <= op;
      end if;
```

```
  end case;
```

```
end process;
```

```
ready <= '1' when state_reg=idle else '0';
```

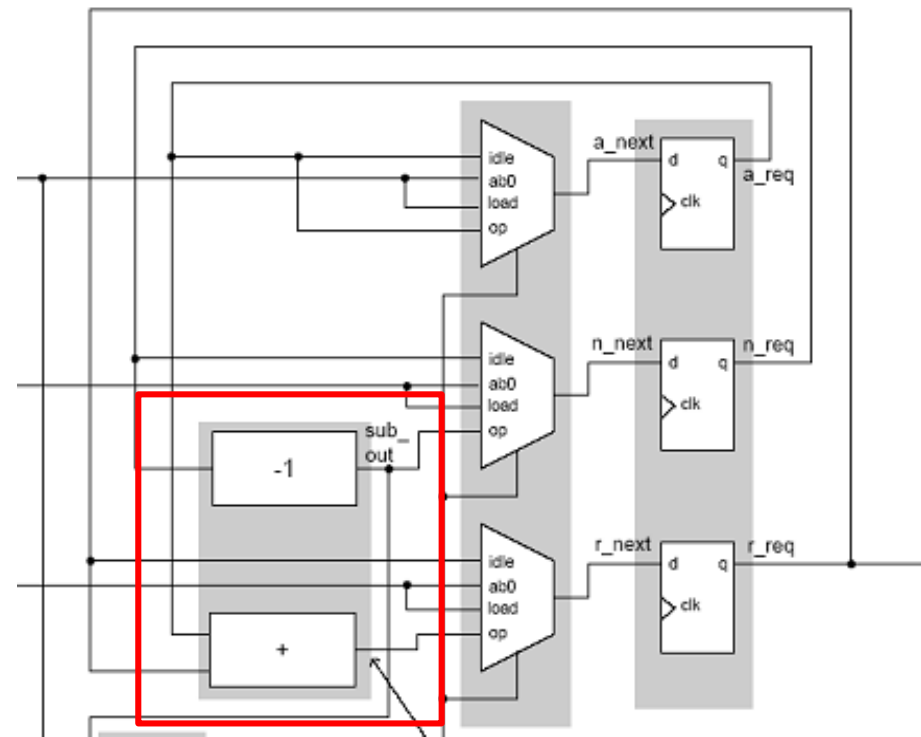
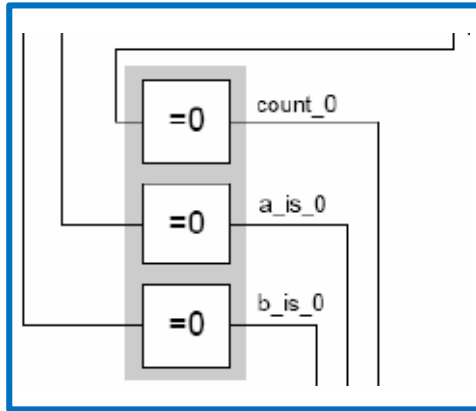


Data Path (Data Registers)

```
— data path: data register
process (clk, reset)
begin
    if reset='1' then
        a_reg <= (others=>'0');
        n_reg <= (others=>'0');
        r_reg <= (others=>'0');
    elsif (clk'event and clk='1') then
        a_reg <= a_next;
        n_reg <= n_next;
        r_reg <= r_next;
    end if;
end process;
```



Data Path (Function Unit)



```
-- data path: functional units  
adder_out <= ("00000000" & a_reg) + r_reg;  
sub_out <= n_reg - 1;
```

```
-- data path: status  
a_is_0 <= '1' when a_in="00000000" else '0';  
b_is_0 <= '1' when b_in="00000000" else '0';  
count_0 <= '1' when n_next="00000000" else '0';
```



Data Path (Multiplexer Routing)

```
process (state_reg, a_reg, n_reg, r_reg,  
        a_in, b_in, adder_out, sub_out)
```

```
begin
```

```
  case state_reg is
```

```
    when idle =>
```

```
      a_next <= a_reg;
```

```
      n_next <= n_reg;
```

```
      r_next <= r_reg;
```

```
    when ab0 =>
```

```
      a_next <= unsigned(a_in);
```

```
      n_next <= unsigned(b_in);
```

```
      r_next <= (others => '0');
```

```
    when load =>
```

```
      a_next <= unsigned(a_in);
```

```
      n_next <= unsigned(b_in);
```

```
      r_next <= (others => '0');
```

```
    when op =>
```

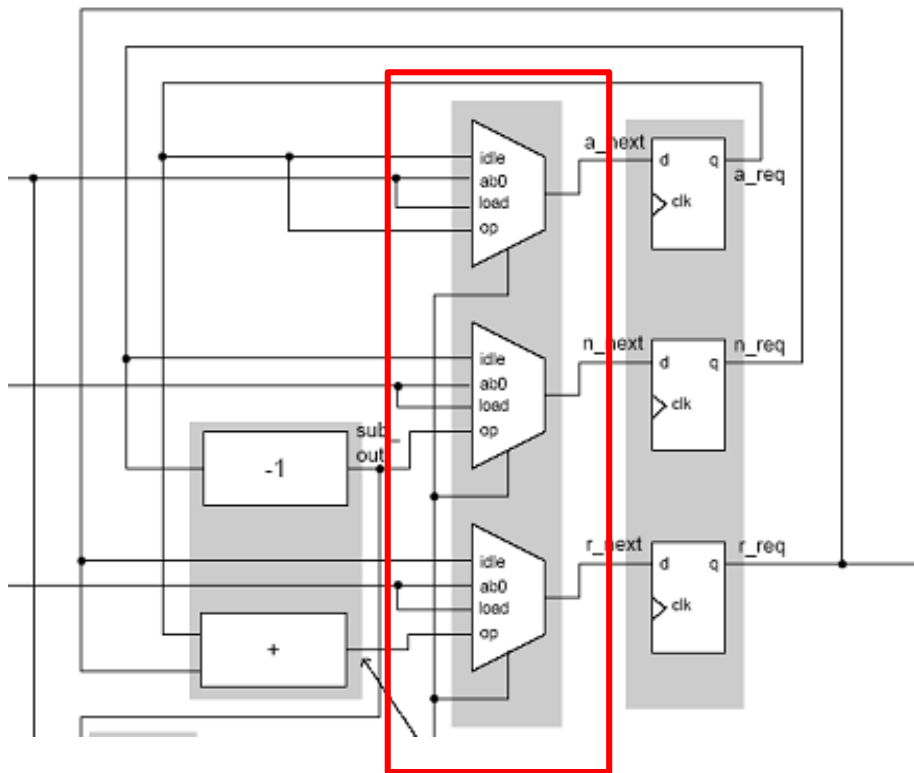
```
      a_next <= a_reg;
```

```
      n_next <= sub_out;
```

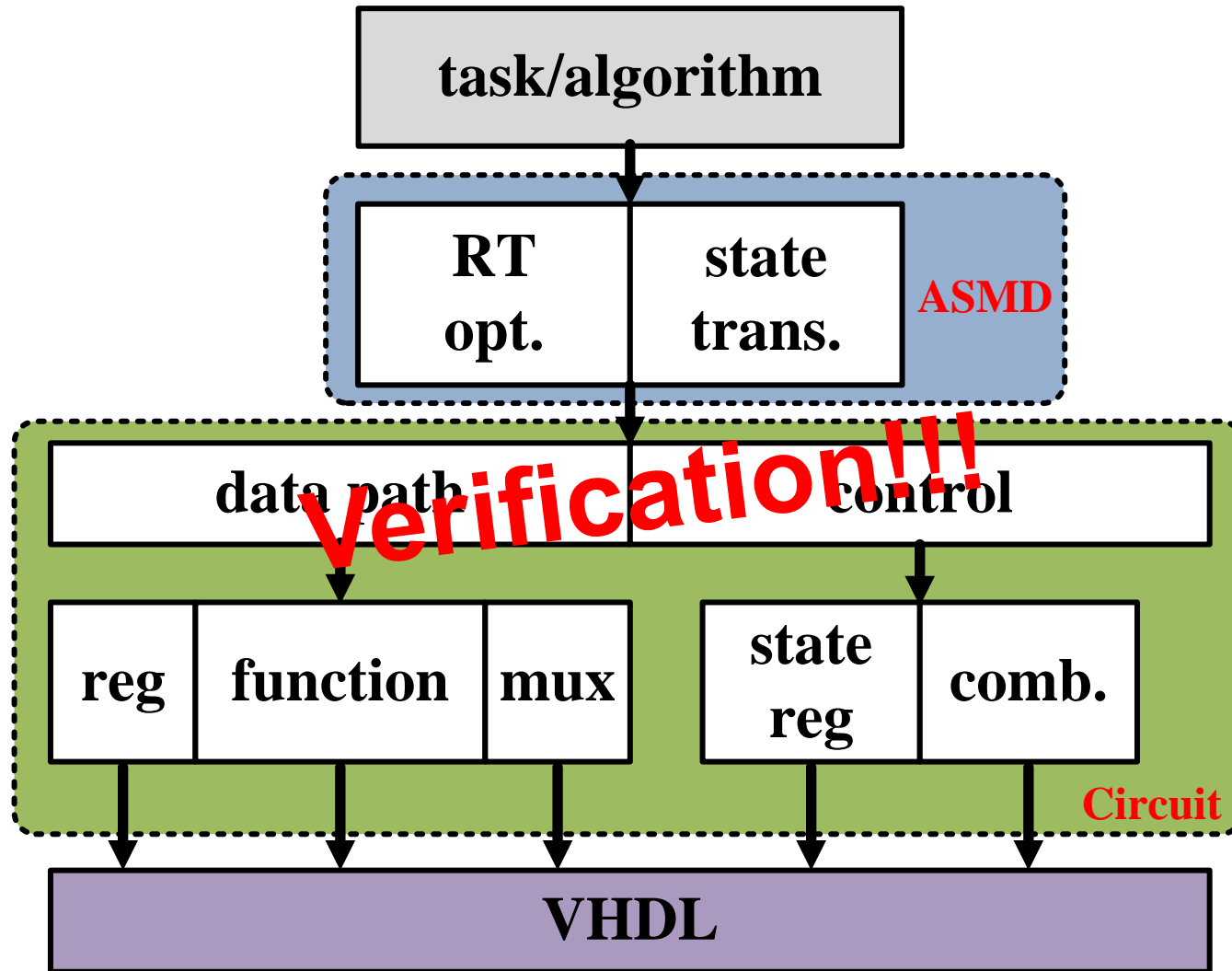
```
      r_next <= adder_out;
```

```
  end case;
```

```
end process;
```



Design Flow



Thanks!

