Unix (and Linux) Security

- Identification and Authentication
- Access Control
- Other security related things:
  - Devices, mounting filesystems
  - Search path
  - Race conditions
- NOTE: filenames may differ between OS/distributions

Users

- Principals have unique UID
  - System cares about ID, not name
  - Several users can have different names but same ID. Then they are treated as the same.
- Superuser (root) has UID = 0
  - There is only one superuser
- Stored in /etc/passwd

- Processes are subjects.

UIDs for Processes

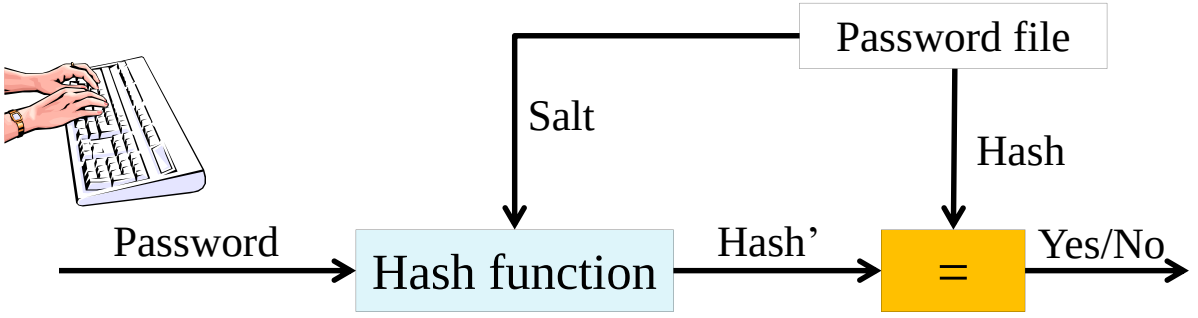
- Real user ID – The ID of the logged in principal
  - Can only be changed by root (effective user ID = 0) → this is how login works
- Effective user ID – The ID used for access control
  - Can be changed by root (effective user ID = 0) to anything
    - Used by processes with effective user ID = 0 when they temporarily access files as a less privileged user
  - Can be changed by anyone (any effective user ID) to real user ID
    - This process has to be able to get back to effective user ID = 0
- Same rules apply to group ID

Groups

- Can not associate multiple user IDs with one file
  - We have to put users in groups if we want several users to have access to the file
- Every user belongs to a primary group.
- Older Unix: Can only be in one group at a time
- Newer Unix and Linux: Can be in several groups at the same time
  - New files are associated with current group ID of user
  - Process group ID is the current group ID of user running the process
- Change group (newgrp)
- Primary group given in /etc/passwd
- Other groups in /etc/group
  - A group can not belong to a group

```
users:x:100:
Students:x:1000:alice,bob
```

Authentication



- Salt is always used
- Hash function and salt will depend on OS
- We look at three variants

Traditional crypt (Password Hashing)

- Design dates back to 1976
- Based on DES
- Password up to 8 characters, salt 12 bits
  - Take least significant 7 bits → 56 bit key
  - Encrypt zero string 25 times with DES
  - If bit i = 1 in salt, swap bits i and i + 24 in E-box output
  - Output 12 + 64 = 76 bits. Encode to 13 characters.

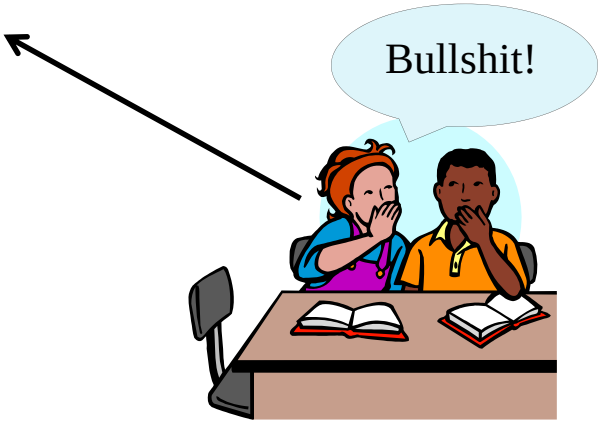
- Problems: Short passwords, short salts, constant cost (and fast function)

Other Alternatives – MD5 crypt

- MD5 crypt
  - Developed for FreeBSD to avoid export restrictions and allow longer passwords (up to 2^{64} bits)
  - Algorithm uses 1000 iterations → slow
  - Salt 12-48 bits
  - Output: 1 'salt' $ 128 bit hash output
- Problem: Constant cost

Other Alternatives – bcrypt

- Based on block cipher blowfish
- Password up to 72 characters, 128 bit random salt
- Internal loop with variable cost
- Output $2a$cost$salt + 192 bit hash output
- Default in OpenBSD
- All problems solved



Current defaults on different systems

- Arch Linux SHA-512-crypt with 5000 rounds
- Ubuntu 18.04 SHA-512-crypt with 5000 rounds
- FreeBSD 12 SHA-512-crypt with 5000 rounds
- OpenBSD 6.4 bcrypt with cost 10
- macOS 10.8+ PBKDF2-SHA512 with ~40000 rounds

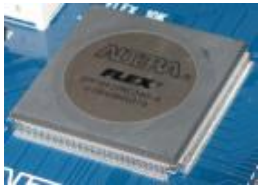
Comparison

	DES crypt	MD5 crypt	bcrypt	sha512crypt
Password length	max 8 chars	virtually any	max 72 chars	virtually any
Salt length	12 bits	12-48 bits	128 bits	typically 48 bits
Variable cost	No	No	Yes	Yes
Hashes/sec	9 800 000	120 000	170	2 500

- Hashes/sec based on 3.4 GHz processor with 4 cores, approximate values given
- The given performance for bcrypt with cost 10, and for sha512crypt with 5000 rounds

Final words on our password discussion

- ”All problems solved” is kind of bullshit
- Some devices can be really fast to a low cost
  - With enough money they are really really really fast
  - Several instances can be implemented in parallel
- Can no longer compare
  - CPU – ”needed” when verifying password
  - GPU, FPGA, ASIC – used by attackers



FPGA/ASIC

GPU

- Make this more fair by making hashing more difficult (costly) for GPUs, FPGAs and ASICs
- **Example**: Argon2 – variable cost in both time and memory

The File /etc/passwd

- Store user (principal) information

Format:

Username:password:UID:GID:ID string:home directory:login shell

- File is world readable
- Example:

```
alice:x:1004:100::/home/alice:/bin/bash
bob:x:1005:100::/home/bob:/bin/bash
```

The File /etc/shadow

- Save passwords in a non-world readable file
  - Username
  - (hashed) password
  - Date of last change (days since Jan 1, 1970)
  - Minimum days between password changes (0 means anytime)
  - Maximum days of validity
  - Days in advance to warn user about change
  - Days account is active after password expired
  - Date of account disabling (days since Jan 1, 1970)
  - Last entry is reserved

```
alice:$6$Gar7uDv0$Ihuwd...wKGlNnWavx:17912:0:99999:7:::
bob:$6$q1/LoHbE$7Md2k...hAtXiw4hW.:17912:0:99999:7:::
```

Access Control

- Discretionary access control – owner of file can change permissions
- Three categories: User (owner), Group, Other (world)
- Three access rights: Read, Write, Execute

```
alice@eita25:/data/1$ ls -l
total 8
drwxr-xr-x 1 alice Students 26 Jan 16 10:05 directory
-rw-r--r-- 1 alice Students 31 Jan 16 10:04 file1
-rw-r--r-- 1 alice Students 10 Jan 16 10:04 file2
```

Other info from ls -l

Link counter, owner, group, size, date of last change, name

Order of Checking

1. Owner
2. Group
3. Other

Consequence:

if owner = r and other = rw then owner has no write permission

```
alice@eita25:/data/2$ ls -l
total 0
-r--rw-rw- 1 alice Students 0 Jan 16 10:06 file
alice@eita25:/data/2$ echo hello > file
bash: file: Permission denied
```

```
bob@eita25:/data/2$ ls -l
total 0
-r--rw-rw- 1 alice Students 0 Jan 16 10:06 file
bob@eita25:/data/2$ echo hello > file
bob@eita25:/data/2$
```

Permissions For Directories

- Read = list the directory
- Write = Delete, rename and insert files in directory
- Execute = access directory and access files in directory

```
alice@eita25:/data/3$ ls -la
total 0
dr-xr-xr-x 1 alice Students  8 Jan 16 10:09 .
drwxr-xr-x 1 root  root     36 Jan 16 10:33 ..
-rw-rw-rw- 1 alice Students  0 Jan 16 10:09 file
alice@eita25:/data/3$ rm -f file
rm: cannot remove 'file': Permission denied
```

```
alice@eita25:/data/4$ ls -la
total 0
drwxr-xr-x 1 alice Students  8 Jan 16 10:10 .
drwxr-xr-x 1 root  root     36 Jan 16 10:33 ..
-rw-r--r-- 1 root  root      0 Jan 16 10:10 file
alice@eita25:/data/4$ rm -f file
alice@eita25:/data/4$
```

Change Permissions – chmod

- Used to change permissions on files
- Mnemonics can be used: user, group, other, all, read write execute.
- Examples:

chmod u+rw file
chmod u=r file
chmod a+rwx file
chmod u-w,g+r,o+r file
chmod a-rwx,u+r file1 file2

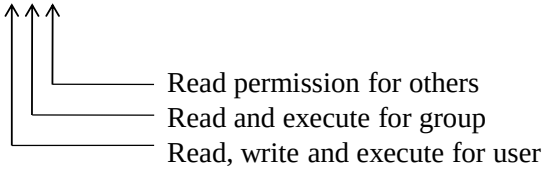
Change Permissions – chmod

- Alternatively, numbers can be used.
- See each group of permissions as one number.
  - Read = 4
  - Write = 2
  - Execute = 1

Sum gives permission

- Example:

chmod 754 file




```
alice@eita25:/data/5$ chmod 754 file; ls -l file
-rwxr-xr-- 1 alice Students 0 Jan 16 10:12 file
```

Setuid and Setgid (programs)

- Controlled invocation
- Effective ID of process is ID of program owner (usually root)
  - Here is the situation when RUID ≠ EUID
- Used to temporarily change access rights
- *x* is replaced by *s*

```
alice@eita25$ ls -l
total 40
-rwxr-sr-x 1 root root 16568 Jan 16 10:17 prog_setgid
-rwsr-xr-x 1 root root 16568 Jan 16 10:17 prog_setuid
alice@eita25$ ./prog_setuid &
[1] 249
alice@eita25$ ./prog_setgid &
[2] 250
alice@eita25$ ps -C prog_setgid,prog_setuid -o pid,ruser,euser,rgroup,egroup,args
  PID RUSER    EUSER    RGROUP   EGROUP   COMMAND
  249 alice    root     users    users    ./prog_setuid
  250 alice    alice    users    root     ./prog_setgid
```

Setuid and Setgid (Directories)

- Setuid on directory usually ignored
- Setgid on directory causes new files to get the same group as directory

```
alice@eita25:/data/7$ ls -l
total 0
drwxr-s--- 1 alice root 0 Jan 16 10:19 directory
alice@eita25:/data/7$ cd directory; touch file; ls -l
total 0
-rw-r--r-- 1 alice root 0 Jan 16 12:55 file
```

Without setgid, file would get the group which is current group ID for user (set by newgrp and defaults to primary group).

Allows users to share files more easily

Important SUID Programs

- **/usr/bin/passwd** change password
- **/bin/su** change UID program

```
alice@eita25:/data$ ls -l /usr/bin/passwd /bin/su
-rwsr-xr-x 1 root root 44664 Jan 25  2018 /bin/su
-rwsr-xr-x 1 root root 59640 Jan 25  2018 /usr/bin/passwd
```

Setuid and setgid:
chmod u+s file or chmod 4XXX file
chmod g+s file or chmod 2XXX file

The dangers of the setuid bit

- Extremely important to write correct code, since it will run as root
- Anything you code *will* be used against you
- Example of vulnerability: CVE-2018-14665 from October 2018
  - Affected most major Linux and BSD distributions
  - Xorg-server, binary `Xorg` has setuid bit set, so regular users can launch X
  - Launching X with `Xorg -logfile /etc/shadow :1` overwrites /etc/shadow with garbage (the output from X)
  - `Xorg -fp "root::16431:0:99999:7:::" -logfile /etc/shadow`
  - The command above also adds an extra line to shadow, which sets a blank root password
  - Any user can now get root privileges simply by issuing `su`

Sticky Bit

- Historically used to keep program code in memory when exiting program (still the case in, e.g. HP-UX)
- Now used to only let owner delete file
  - directory owner and superuser can also delete it

```
bob@eita25:/data/8$ ls -la
total 0
drwxrwxr-t 1 alice Students  8 Jan 16 10:26 .
drwxr-xr-x 1 root  root     36 Jan 16 10:33 ..
-rw-rw-r-- 1 alice Students  0 Jan 16 10:26 file
bob@eita25:/data/8$ rm file
rm: cannot remove 'file': Operation not permitted
```

```
bob@eita25:/data/9$ ls -la
total 0
drwxrwxr-x 1 alice Students  8 Jan 16 10:26 .
drwxr-xr-x 1 root  root     36 Jan 16 10:33 ..
-rw-rw-r-- 1 alice Students  0 Jan 16 10:26 file
bob@eita25:/data/9$ rm file
bob@eita25:/data/9$
```

- Typical example: the directory /tmp has sticky bit set

Default Access Rights (umask)

- Control default permissions, stored in /etc/profile
- Override in ~/.profile or in prompt
- `umask` tells which permissions to exclude by default
- Access = full access AND NOT(umask)
  - Full access for programs and directories: 0777
  - Full access for files: 0666

```
alice@eita25:/data/a$ umask 0027; mkdir directory; touch file; ls -
l
total 0
drwxr-x--- 1 alice users 0 Jan 16 12:59 directory
-rw-r----- 1 alice users 0 Jan 16 12:59 file
```

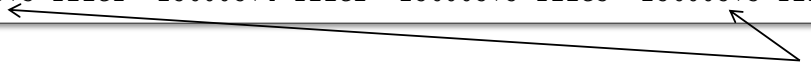
Change Owner and Group (chown and chgrp)

- chown is used to change the owner of a file (or directory)
- chgrp is used to change the group of a file (or directory)
  - chown can set group also
- Possible problem: A user creates a suid program and owner gets changed to root
- Common solution:
  - Only root can change owner and setuid and setgid bits are removed when owner is changed
  - Anyone can change group to a group they are member of, but setuid and setgid bits are removed when group is changed
- Other solutions possible
  - Let only root use chown, but preserve setuid and setgid bits
  - Let any user change owner on his/her own files, but remove setuid and setgid bits

The inode

- Stores file information
- Directory contains filename and inode number

```
alice@eita25:/data/b$ ls -i
25600573 file1  25600574 file2  25600575 file3  25600573 file4
```

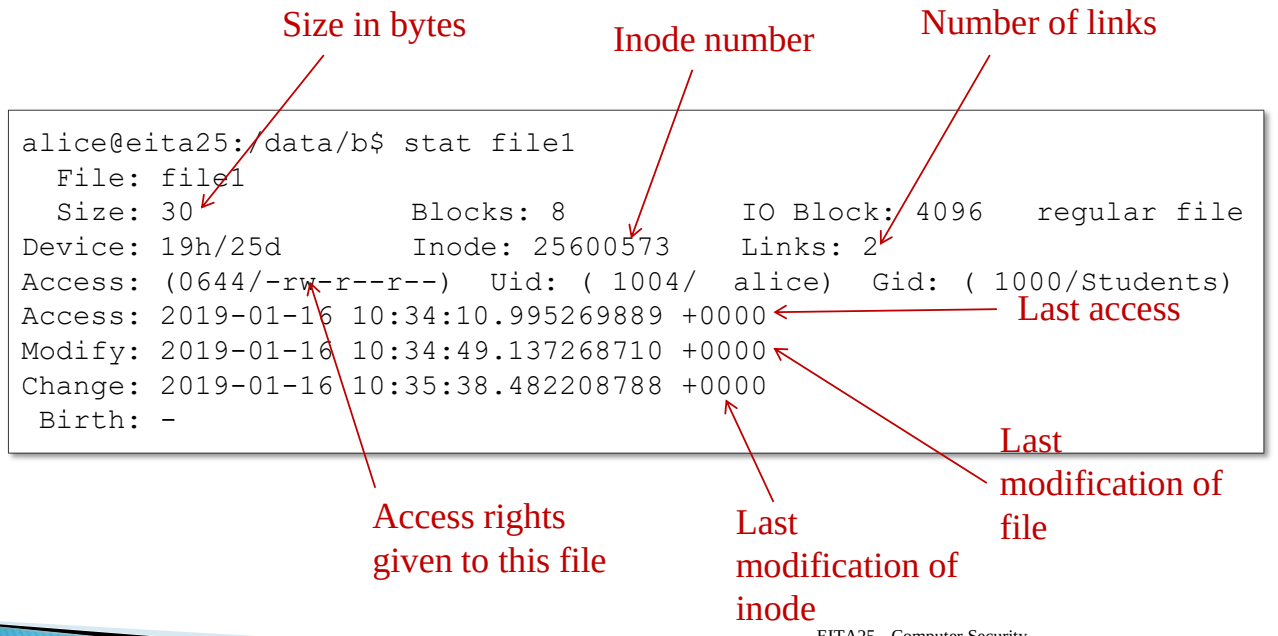


Note that file1 and file4 points to the same inode

- inode contains e.g.:
  - Access rights
  - Owner (UID)
  - Group (GID)
  - Time of latest access, modification and change
  - Size of file
  - Pointers to block of data

inode Information (stat)

- Some information about an inode can be found using `stat`



Copy files

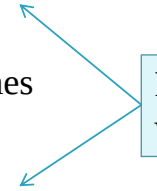
- Files can be copied in two ways
- `cp src dest`
  - Creates a new inode and new physical file owned by user running cp
- `ln target linkname`
  - Creates filename and pointer to target's inode. No new file is created.
  - When one filename is deleted the other is still there and the file is not deleted
  - rm subtracts the number of links in the inode by 1. If it becomes zero the corresponding data block is freed
- `ln –s target linkname`
  - Creates a symbolic link, not a real link
  - When opening symbolic link for reading or writing link is automatically dereferenced
  - If target is deleted, the symbolic link remains, pointing to nothing

Race conditions

- Assume process ”proc” with effective user ID = 0 writes to files in /tmp directory
  - Process creates e.g., /tmp/file and writes temporary data to this file
- What if malicious user creates /tmp/file as symbolic link to /etc/passwd?
  - The file /etc/passwd will be overwritten since ”proc” has write access to this file
  - System is damaged
- Race condition: Who creates the file first

Solutions To This Race Condition

- Create files with unpredictable filenames in /tmp
  - Still, attacker can try thousands of filenames and will succeed with probability > 0
- Use O_CREAT and O_EXCL flag when opening file
  - Then open fails if file already exists
  - Will not follow symbolic links during creation either



Function mkstemp() will do this

Protection of devices

- Devices are treated as files
- **Example:** If you can read/write physical memory all access control is overruled!
- **/dev/mem** is the physical memory
- **/dev/sda** is the first disk drive (in Linux)

Mounting File Systems

- Different physical devices put under a single root "/"
- The mounted file system may contain unwelcome programs
  - nosuid – turn off SUID and SGID bits
  - noexec – no binaries can be executed
  - nodev – no devices can be accessed
  - ro – read-only
- UIDs and GIDs are local identifiers that need not be interpreted the same on different Unix systems
  - Need *global unique* identifiers on networks

Searchpath

- When executing programs, system needs to know where to look for it → PATH tells system where to look
- PATH=. : $HOME/bin:/usr/bin:/bin:
  - Programs can be located in current directory + 3 bin directories
  - Trojan horse
- Can be a bad idea to put your current directory in the search path (especially for programs executed by root)
- Order matters, current directory last reduces risk
- Alternatively, call program by full name

Before the labs

- Labs start next week (yay!)
- Sign up if you haven't already
  - Lots of free spots Friday afternoon :D
- There are preparatory assignments for all labs

- Start preparations early for lab 2
- Also, start preparations early for lab 2
- Note that there are two question hours especially for lab 2
  - Check course home page under Labs for exact hours

- You are not alone! You can ask me questions