

Problem 1

a)

$$mK = (\alpha, \alpha^2) \begin{pmatrix} \alpha^2 & \alpha + 1 \\ \alpha^2 + 1 & 1 \end{pmatrix} = (1, \alpha).$$

b) If K is not invertible, unique decryption is not possible.

c)

$$|\mathcal{K}| = (q^2 - 1)(q^2 - q),$$

where $q = 2^3$. It gives $H(K) = 11.78$, and $N_0 = \lceil 11.78/1 \rceil = 12$.

Problem 2

This is exactly the last exercise of Laboratory 1!

We construct two sequences, one over \mathbb{Z}_2 and one over \mathbb{Z}_5 , which are combined through the Chinese remainder theorem to give a sequence over \mathbb{Z}_{10} .

We will receive a de Bruijn sequence if both sequences are de Bruijn. Let's start with the one over \mathbb{Z}_2 . A de Bruijn sequence is obtained if we take a maximal length sequence (primitive feedback polynomial) and add the all zero state into the cycle. Take the primitive polynomial $x^4 + x + 1$ over \mathbb{Z}_2 as feedback polynomial. Then we have a state transition $1000 \rightarrow 0001$ which should be changed to the state transition $1000 \rightarrow 0000 \rightarrow 0001$. This is done by adding a nonlinear part f to the feedback. Let the LFSR state be denoted by s_3, s_2, s_1, s_0 . Then $f(s_3, s_2, s_1, s_0) = s_2' s_1' s_0'$, where s_i' means inversion of the s_i variable.

Then we do the same for the sequence over \mathbb{Z}_5 . Finally, we combine them. Let \mathbf{u} be the de Bruijn sequence over \mathbb{Z}_2 , and let \mathbf{v} be the de Bruijn sequence over \mathbb{Z}_5 . The resulting sequence w over \mathbb{Z}_{10} is given by

$$w_i = 5 \cdot u + v.$$

DRAW A PICTURE, choose two starting states for the LFSRs, and produce 10 symbols of the sequence.

Problem 3

a)

C	$M = (0, 0)$	$M = (0, 1)$	$M = (1, 0)$	$M = (1, 1)$
$K = (0, 0)$	(0, 0, 0)	(0, 1, 0)	(1, 0, 0)	(1, 1, 0)
$K = (0, 1)$	(0, 0, 0)	(0, 1, 1)	(1, 0, 1)	(1, 1, 0)
$K = (1, 0)$	(0, 0, 1)	(0, 1, 1)	(1, 0, 1)	(1, 1, 1)
$K = (1, 1)$	(0, 0, 1)	(0, 1, 0)	(1, 0, 0)	(1, 1, 1)

$$P_I = \frac{1}{2},$$

For all C we can find a ciphertext C' such that $P(C' \text{ valid} | C) = 1$, e.g., $C = (0, 0, 0)$, $C' = (1, 1, 0)$. Thus, we get $P_S = 1$.

b) Since $k = 3$ the secret polynomial is of the form $a(x) = a_0 + a_1x + a_2x^2$. The three shares

give the following three equations,

$$\begin{aligned}a_0 + a_1 + a_2 &= 6, \\a_0 + 3a_1 + 9a_2 &= 8, \\a_0 + 5a_1 + 12a_2 &= 12.\end{aligned}$$

Solving this system of equations gives $a(x) = 9 + 10x^2$ and $K = a_0 = 9$.

Problem 4

- a) Correct.
 - b) Correct, $p_2(z)$ primitive.
 - c) Wrong, $\mathbf{s}_1 = [2]^\infty$.
 - d) Wrong, $L(\mathbf{s}_1 + \mathbf{s}_2) \leq L(\mathbf{s}_1) + L(\mathbf{s}_2) = 1 + 3 = 4 < 6$.
 - e) Wrong, since $p_1(z)$ is not irreducible.
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Problem 5

a) Two possible primes are $p = 4007$, $q = 4013$ giving $n = 16080091$. Since $\gcd(3, \phi(n)) = 1$, we can choose $e = 3$.

public parameters: $(n, e) = (16080091, 3)$

b) $\phi(n) = (p - 1)(q - 1) = 16072072$

If we choose $e = 3$, then by using Euklides algorithm and Bezouts identity we get $d = 10714715$.

trapdoor parameters: $(\phi(n), d, p, q) = (16072072, 10714715, 4007, 4013)$

c)

$$D(E(M)) = C^d = (M^e)^d = M^{ed} = M^{1+k\phi(n)} = M \cdot (M^k)^{\phi(n)} = [\text{Euler}] = M \pmod n.$$
