

Exam March 2018

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Q2:

$$b = 350 \text{ kB} \\ = 350 \times 10^3 \text{ B}$$

$$r = 1.5 \text{ MB/s} \\ = 1.5 \times 10^6 \text{ B/s}$$

$$M = 20 \text{ MB/s} \\ = 20 \times 10^6 \text{ B/s}$$

Maximum burst length

$$S = M \times \frac{b}{M-r} \text{ octets} \\ = 20 \times 10^6 \times \frac{350 \times 10^3}{20 \times 10^6 - 1.5 \times 10^6} \\ = 20 \times \frac{350 \times 10^3}{18.5} \\ = 378.38 \times 10^3 \text{ B} \\ = 378.38 \text{ kB}$$

Q3:

Symbols: $\{0, 1\}$

$$P(\text{detect } 1 | \text{transmit } 1) \\ = P(\text{detect } 0 | \text{transmit } 0) \\ = 0.72$$

$$P(\text{transmit } 1) = 0.2$$

$$P(\text{symbol correct} | \text{detect } 1)$$

$$= P(\text{transmit } 1 \mid \text{detect } 1)$$

Conditional probability:

$$P(A|B) = \frac{P(A \cap B)}{P(B)}$$

$$P(B|A) = \frac{P(A \cap B)}{P(A)}$$

$$\Rightarrow P(A|B) = \frac{P(B|A) P(A)}{P(B)} \quad (\text{Baye's law})$$

$$\therefore P(\text{transmit } 1 \mid \text{detect } 1)$$

$$= \frac{P(\text{detect } 1 \mid \text{transmit } 1) \times P(\text{transmit } 1)}{P(\text{detect } 1)}$$

$$= \frac{0.72 \times 0.2}{P(\text{detect } 1)}$$

$$= \frac{0.144}{P(\text{detect } 1)}$$

$$P(\text{detect } 1) = P(\text{transmit } 1 \text{ and detect } 1) + P(\text{transmit } 0 \text{ and detect } 1)$$

$$= 0.2 \times 0.72 + (1-0.2) \times (1-0.72)$$

$$= 0.144 + 0.8 \times 0.28$$

$$= 0.144 + 0.224$$

$$= 0.368$$

$$\therefore P(\text{transmit } 1 \mid \text{detect } 1)$$

$$= \frac{0.144}{0.368}$$

$$= 0.39.$$

Q4:

See random access tutorial solutions.

$$\sum_{k=1}^{\infty} p^{k-1} (1-p) \times k = \frac{1}{1-p}$$

Q5:

Flow 2 gets 4 times bandwidth
of Flow 1: $4 = \frac{8}{2}$

Flow 3 gets 2.5 times bandwidth
of Flow 1: $2.5 = \frac{5}{2}$

Flow 1 gets 1 times bandwidth
of Flow 1: $1 = \frac{2}{2}$

Weights:

$$w_1 = 1$$

$$w_2 = 4$$

$$w_3 = 2.5$$

P	Size	Flow	Weighted size	Start	Finish
1	50	1	50	0	50
2	80	1	80	50	130
3	30	1	30	130	160
4	100	2	25	0	25
5	150	2	37.5	25	62.5
6	125	2	31.25	62.5	93.75
7	480	3	192	0	192
8	200	3	80	192	272

Transmission order:

4, 1, 5, 6, 2, 3, 7, 8

Q6:

2 nodes \Rightarrow 2 mini-slots per frame

4 data slots per node \Rightarrow 8 data slots per frame

(a) P(a given node reserves 4 slots)

= P(at least 4 packets arrived at the node during the previous frame)

= 1 - P(less than 4 packets arrived)

$$= 1 - \sum_{k=0}^3 \frac{\lambda^k e^{-\lambda}}{k!}$$

$$= 1 - e^{-\lambda} - \lambda e^{-\lambda} - \frac{\lambda^2 e^{-\lambda}}{2} - \frac{\lambda^3 e^{-\lambda}}{6}$$

Each node has the same probability and they are independent, so for both nodes we have:

P(all slots reserved)

$$= \left(1 - e^{-\lambda} - \lambda e^{-\lambda} - \frac{\lambda^2 e^{-\lambda}}{2} - \frac{\lambda^3 e^{-\lambda}}{6} \right)^2$$

(b) Expected number of slots used:

$$K = \sum_{k=0}^7 k \times \frac{(2\lambda)^k e^{-2\lambda}}{k!} + \sum_{k=8}^{\infty} 8 \times \frac{(2\lambda)^k e^{-2\lambda}}{k!}$$

$$= \sum_{k=0}^7 k \times \frac{(2\lambda)^k e^{-2\lambda}}{k!} + 8 \left(1 - \sum_{k=0}^7 \frac{(2\lambda)^k e^{-2\lambda}}{k!} \right)$$

$\lambda = 2$:

$$K = \sum_{k=0}^7 k \times \frac{4^k e^{-4}}{k!} + 8 \left(1 - \sum_{k=0}^7 \frac{4^k e^{-4}}{k!} \right)$$

$$\begin{aligned}
 &= 3.557 + 8 \times 0.051 \\
 &= 3.557 + 0.409 \\
 &= 3.966
 \end{aligned}$$

$$\begin{aligned}
 \text{Utilisation} &= \frac{3.966}{\underbrace{0.1 \times 2}_{\text{mini-slots}} + \underbrace{8}_{\text{data slots}}} \\
 &= \frac{3.966}{8.2} \\
 &= 0.48
 \end{aligned}$$

Q7:

Video calls: Poisson, $\frac{1}{\lambda} = 1.4 \text{ s}$

GBR, $480 \times 10^3 \text{ bps}$ per cell
 Call length average: $4 \times 60 + 30 = 270 \text{ s}$

Web: 60 concurrent sessions on average
 non-GBR

Total capacity: $100 \times 10^6 \text{ bps}$

$$(a) \frac{100 \times 10^6}{480 \times 10^3} = 208$$

(b) Erlang system, $N = 208$

$$h = 270 \text{ s}$$

$$\lambda = \frac{1}{1.4}$$

$$= 0.714 \text{ calls/s}$$

$$A = \lambda h$$

$$= 270 \times 0.714$$

$$= 192.857$$

$$P_B = \frac{A^N}{N!}$$

$$\frac{n!}{\sum_{j=0}^n \frac{A^j}{j!}}$$

$$= \frac{((192 \cdot 857)^{208}) / 208!}{\sum_{j=0}^{208} \frac{(192 \cdot 857)^j}{j!}}$$

$$(\approx 0.016)$$

(c) Carried video call traffic:

$$A_c = A(1 - P_B)$$

$$= 192 \cdot 857 (1 - 0.016)$$

$$= 189.77 \text{ simultaneous calls on average}$$

Average capacity used for video calls:

$$= A_c \times 480 \times 10^3$$

$$= 189.77 \times 480 \times 10^3$$

$$= 91090.286 \times 10^3$$

$$= 91.090286 \times 10^6 \text{ bps}$$

Average capacity available for web traffic

$$= 100 \times 10^6 - 91.090286 \times 10^6$$

$$= 8.91 \times 10^6 \text{ bps}$$

$$= 8.91 \text{ Mbps} = 0.15 \text{ Mbps per connection}$$

(d) Increase, as then some capacity will be reserved for web traffic.

Q9:

Max allowed delay for control data:

10 ms

Logging data $\lambda_d = 400$ pk/s
avg. length = 2000 B

Control traffic: $\lambda_c = 500$ pk/s
avg length = 200 B

Data rate per channel: 8×10^6 kbits/s

(a) Logging channel:

non-persistent CSMA

$$S = \frac{G}{1+G} \quad (\text{Normalised offered load})$$

Normalised offered load

$$G = \frac{2000 \times 8 \times 400}{8 \times 10^6}$$

$$= 0.8$$

Normalised throughput:

$$S = \frac{0.8}{1+0.8}$$

$$= \frac{0.8}{1.8}$$

$$= 0.44$$

Control channel:

1-persistent CSMA

$$G = \frac{200 \times 8 \times 500}{8 \times 10^6}$$

$$= 0.1$$

$$S = \frac{G e^{-G} (1+G)}{G + e^{-G}}$$

$$= \frac{0.1 \times e^{-0.1} (1 + 0.1)}{0.1 + e^{-0.1}}$$

$$= \frac{0.0995}{1.0048}$$

$$= 0.099$$

(b) Data rate to centralized controller: 5×10^6 b/s
 Arrival rate for control packets:

$$S = 0.099 \Rightarrow \lambda'_c = \frac{0.099 \times 8 \times 10^6}{200 \times 8}$$

$$= 495 \text{ ph/s}$$

Arrival rate for logging traffic:

$$S = 0.99 \Rightarrow \lambda'_l = \frac{0.99 \times 8 \times 10^6}{2000 \times 8}$$

$$= 220 \text{ ph/s}$$

FCFS:

$$\lambda = \lambda'_l + \lambda'_c$$

$$= 495 + 220$$

$$= 715 \text{ ph/s}$$

Average packet length

$$L = 2000 \times \frac{220}{715} + 200 \times \frac{495}{715}$$

$$= 615.38 + 138.46$$

$$= 753.846$$

$$N = \frac{5 \times 10^6 \text{ bs}^{-1}}{753.8846 \times 8 \text{ b}}$$

$$= 829.082$$

$$\rho = \frac{\lambda}{\mu}$$

$$= \frac{715}{829.082}$$

$$= 0.86$$

Priority queuing:

$$\rho_c = \frac{\lambda_c}{\mu_c}$$

$$\mu_c = \frac{5 \times 10^6 \text{ bs}^{-1}}{200 \times 8 \text{ b}}$$

$$= 3125 \text{ pkts/s}$$

$$\rho_c = \frac{495}{3125}$$

$$= 0.1584$$

(c) Delay band for M/M/1:

$$P(T_R \leq t) = 1 - e^{-\mu(1-\rho)t}$$

$$t = 10 \times 10^{-3} \text{ s}$$

$$\rho = 0.86$$

$$\mu = 829.082 \text{ pkts/s}$$

$$P(T_R \leq 10 \times 10^{-3})$$

$$= 1 - e^{-829.082(1-0.86) \times 10 \times 10^{-3}}$$

$$= 1 - e^{-1.16}$$

$$= 0.69$$

Priority queuing:

$$\rho = 0.1584$$

$$\mu = 3125 \text{ h/s}$$

$$P(T_R \leq 10 \times 10^{-3})$$

$$= 1 - e^{-3125(1-0.1599) \times 10 \times 10^{-3}}$$

$$= 1 - e^{-26.3}$$

$$= 1 - 3.8 \times 10^{-12}$$

$$\approx 1.0$$