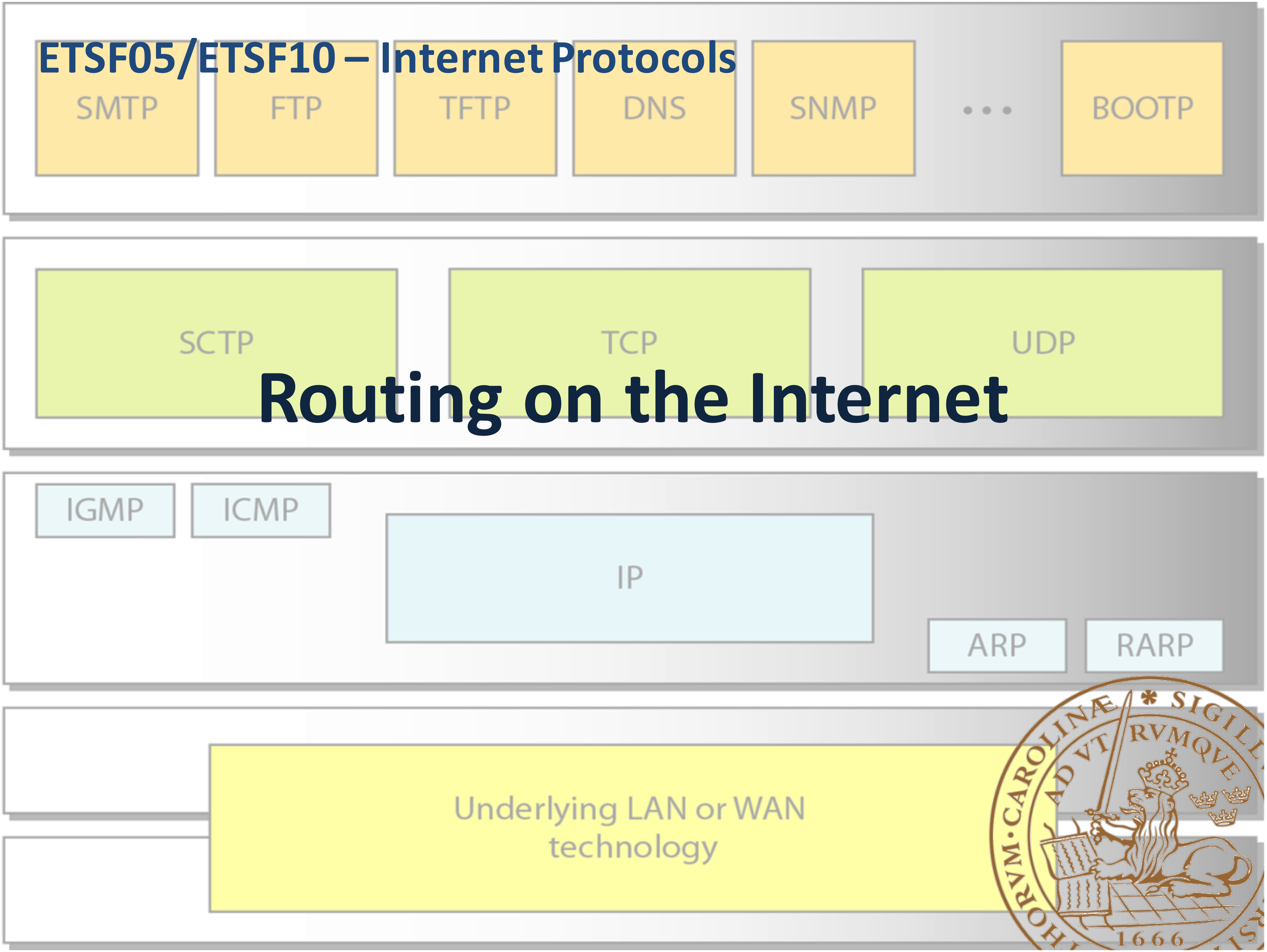
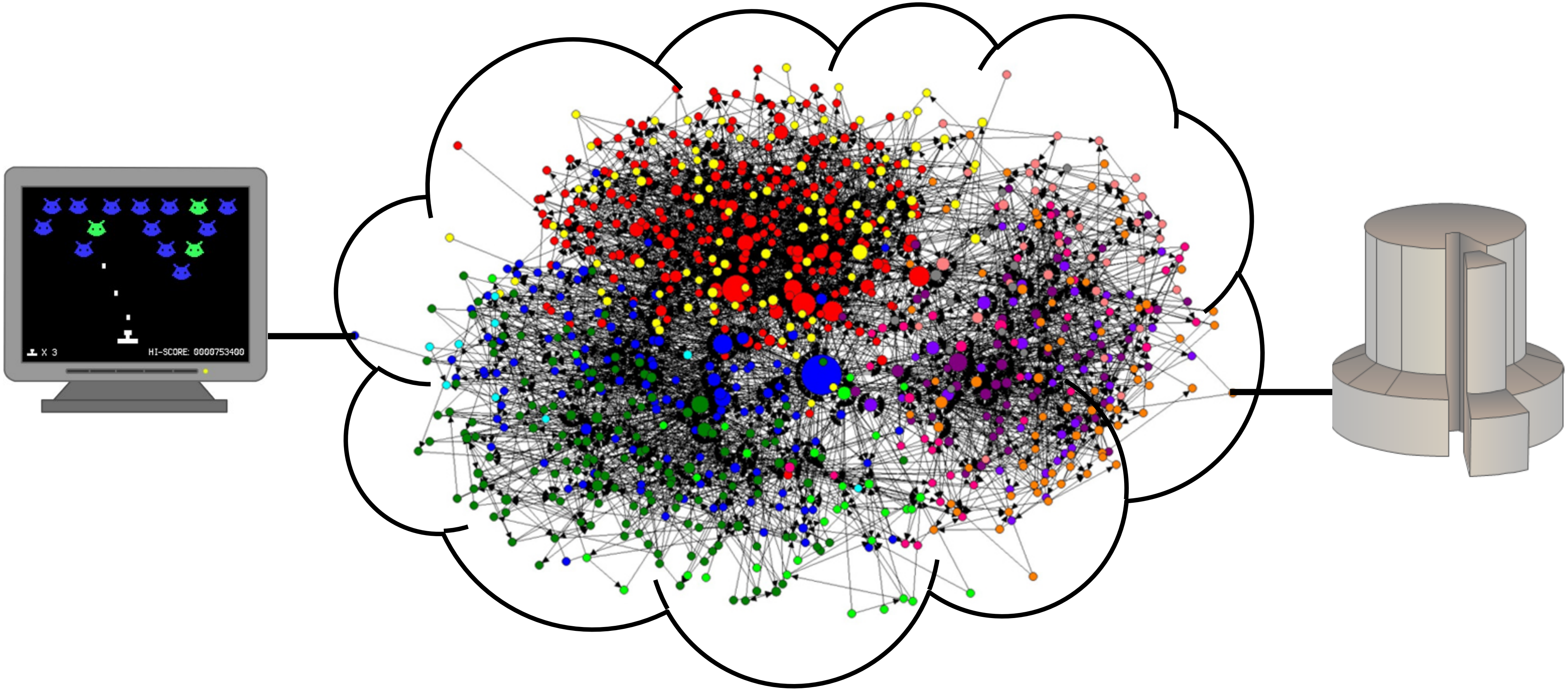
ETSF05/ETSF10 – Internet Protocols

Routing on the Internet



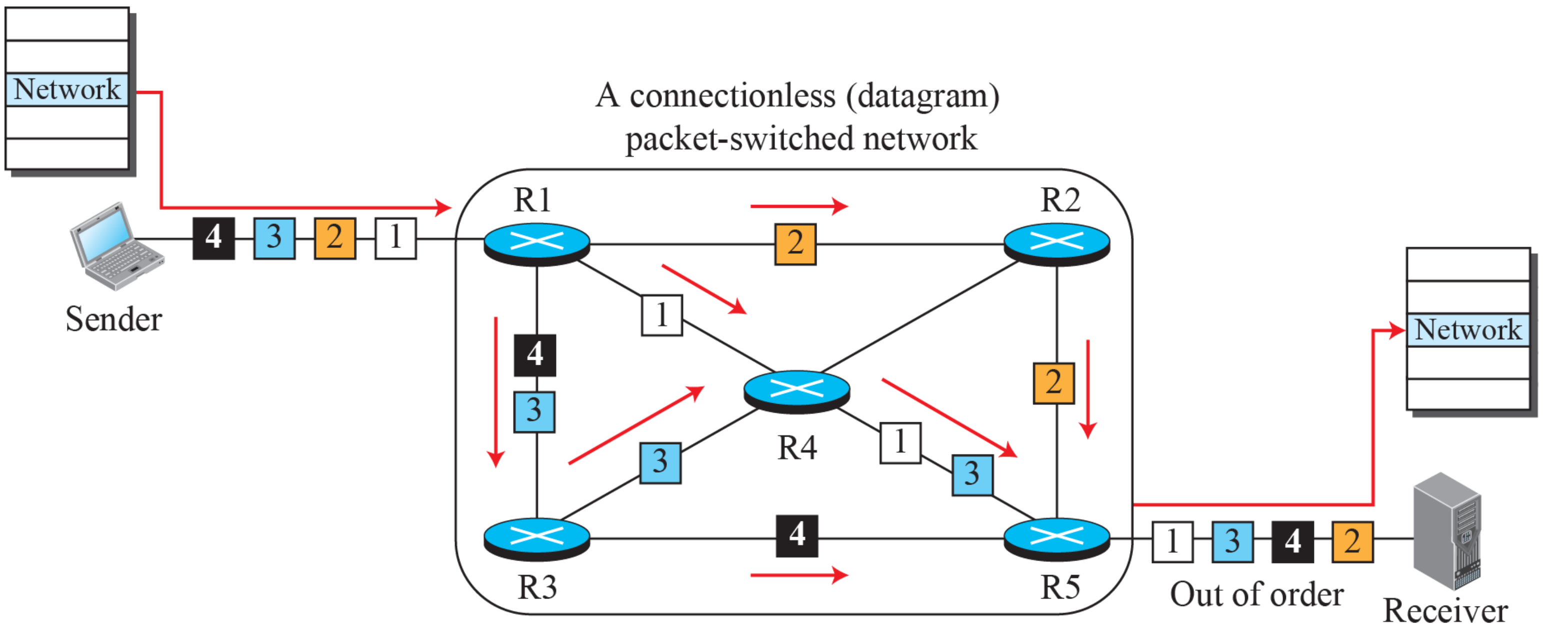
How does routing work?



Packet-switched Routing

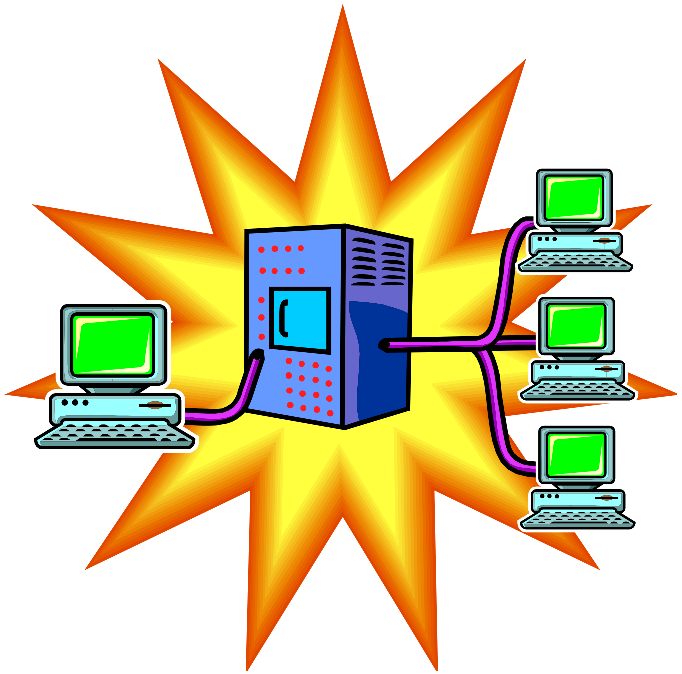
Choosing an optimal path

- According to a cost metric
- Decentralised forwarding
  - each router has full/necessary information

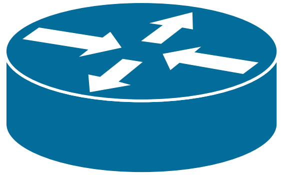


Routing in Packet Switching Networks

- Select route across network between end nodes
- Characteristics required:
  - Correctness
  - Simplicity
  - Robustness vs Stability
  - Fairness vs Optimality
  - Efficiency



Router

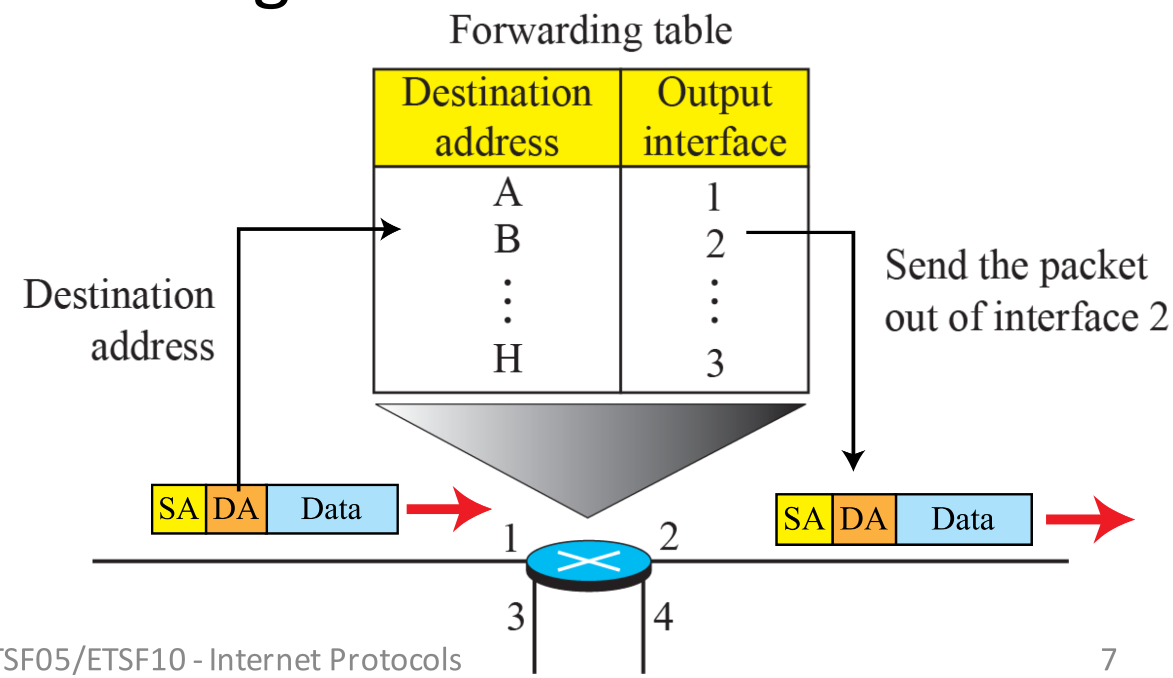


- Internetworking device
  - Passes data packets between networks
  - Checks ***Network Layer*** addresses
  - Uses Routing/forwarding tables

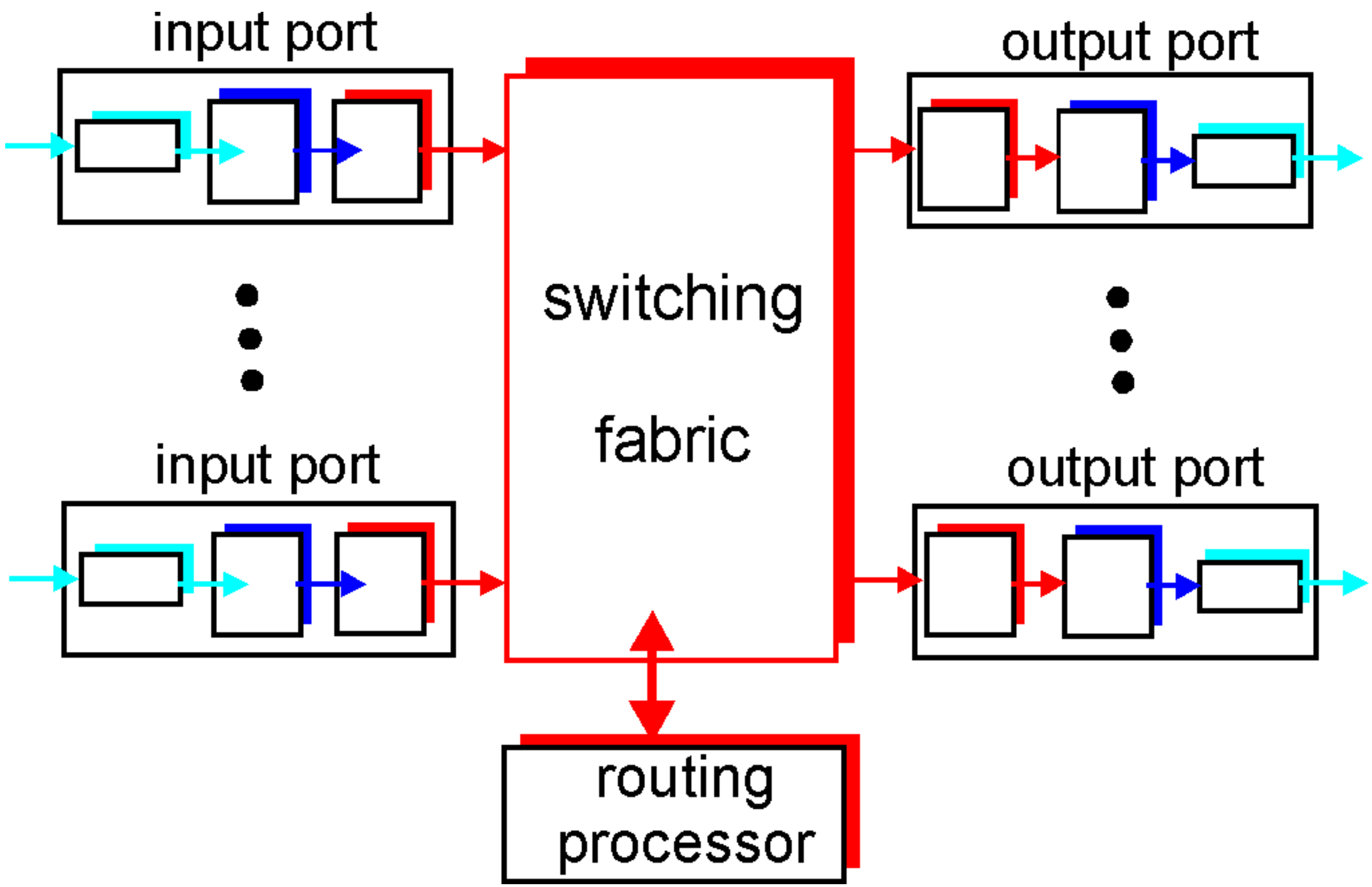
Two functions:

❶ Routing

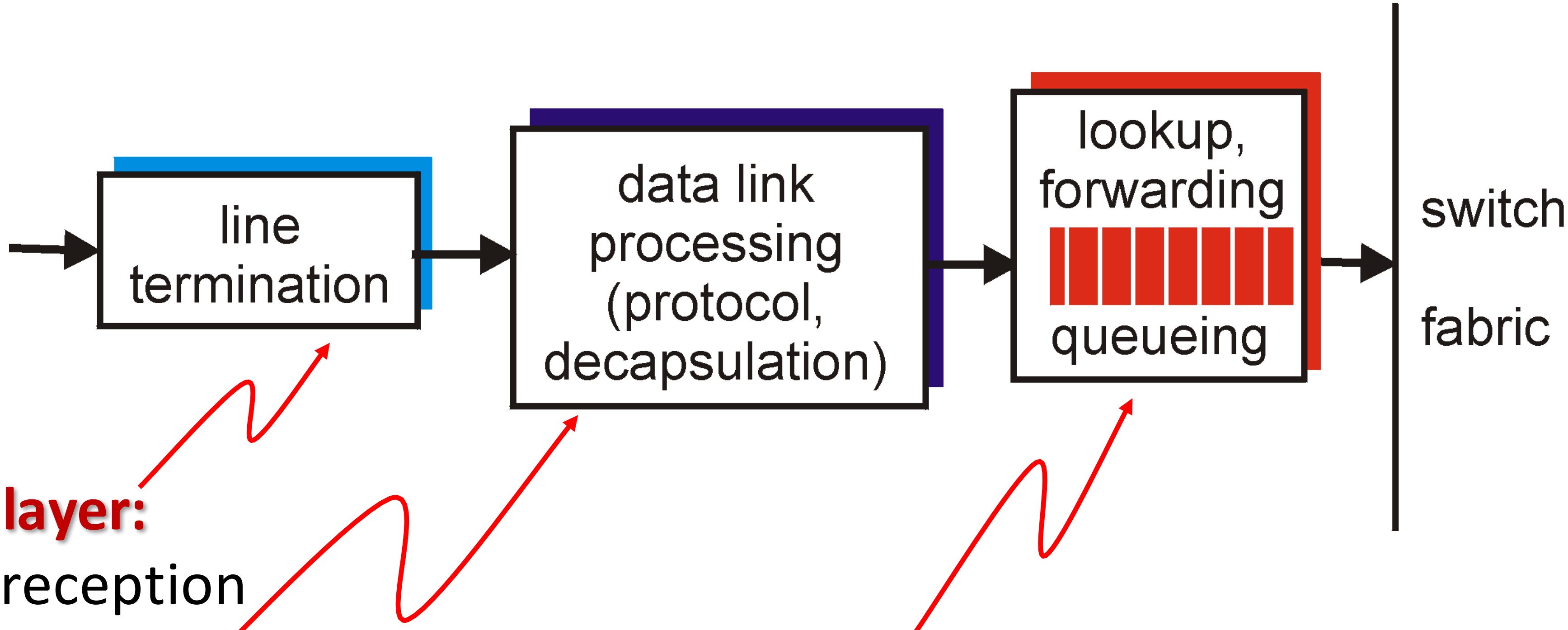
❷ Forwarding



Router Architecture Overview



Input Port



Physical layer:
bit-level reception

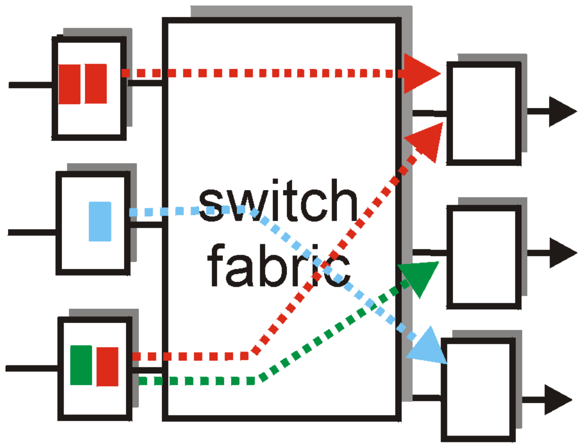
Data link layer:
e.g., Ethernet

Decentralized switching:

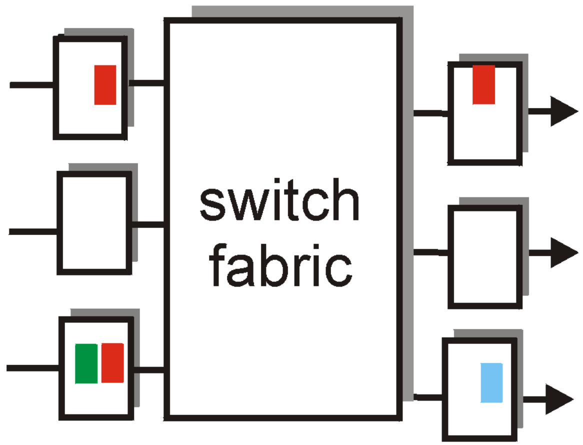
- Given destination, lookup output port using routing table
- Goal: complete input port processing at 'line speed'

Input Port Queuing

- Fabric slower that sum of input ports → **queuing**
- **Delay and loss** due to input buffer overflow
- **Head-of-the-Line (HOL) blocking**: Datagram at front of queue prevents others in queue from proceeding

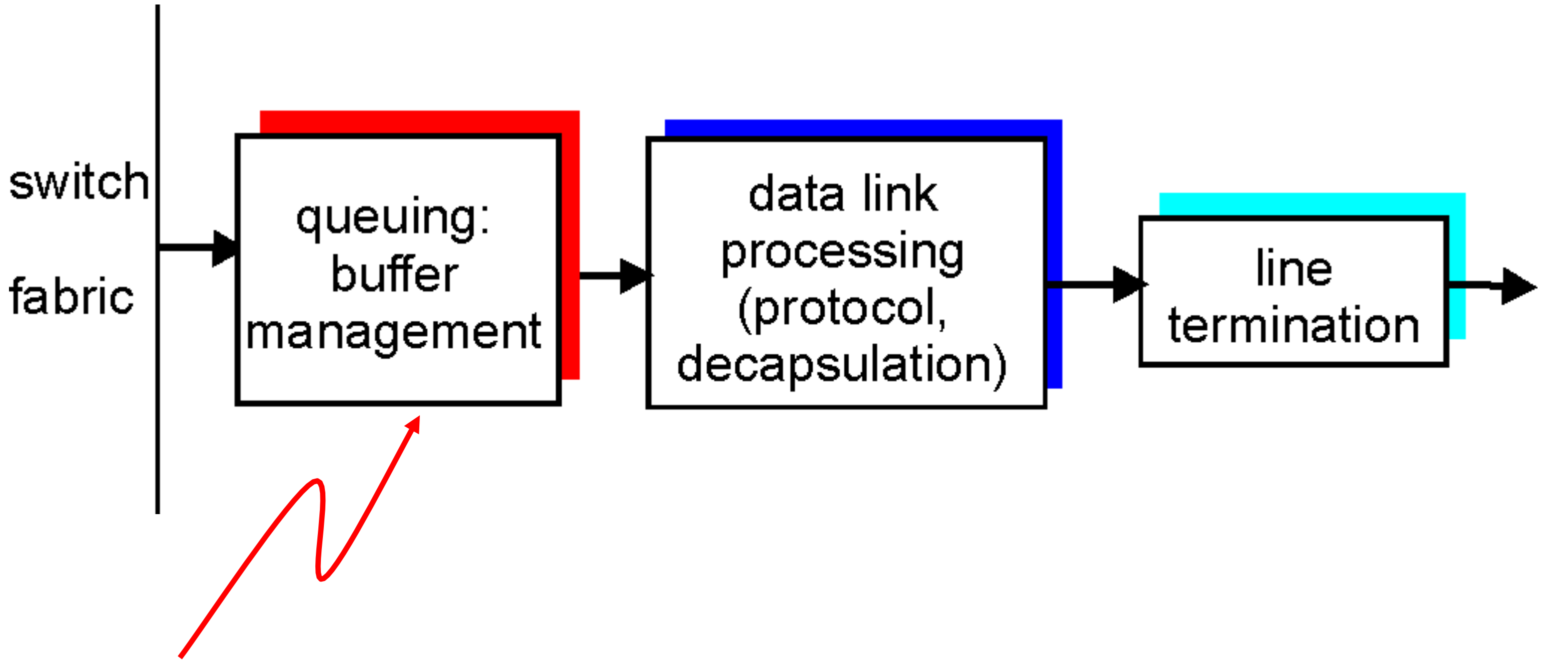


output port contention at time t - only one red packet can be transferred



green packet experiences HOL blocking

Output Port

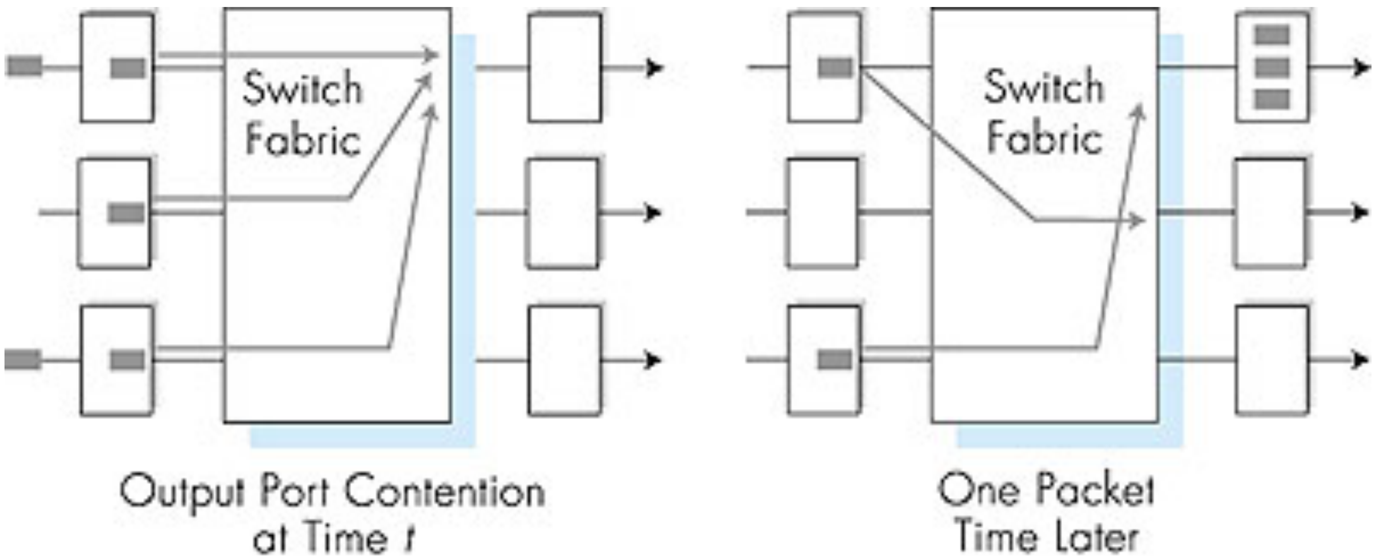


Priority Scheduling:

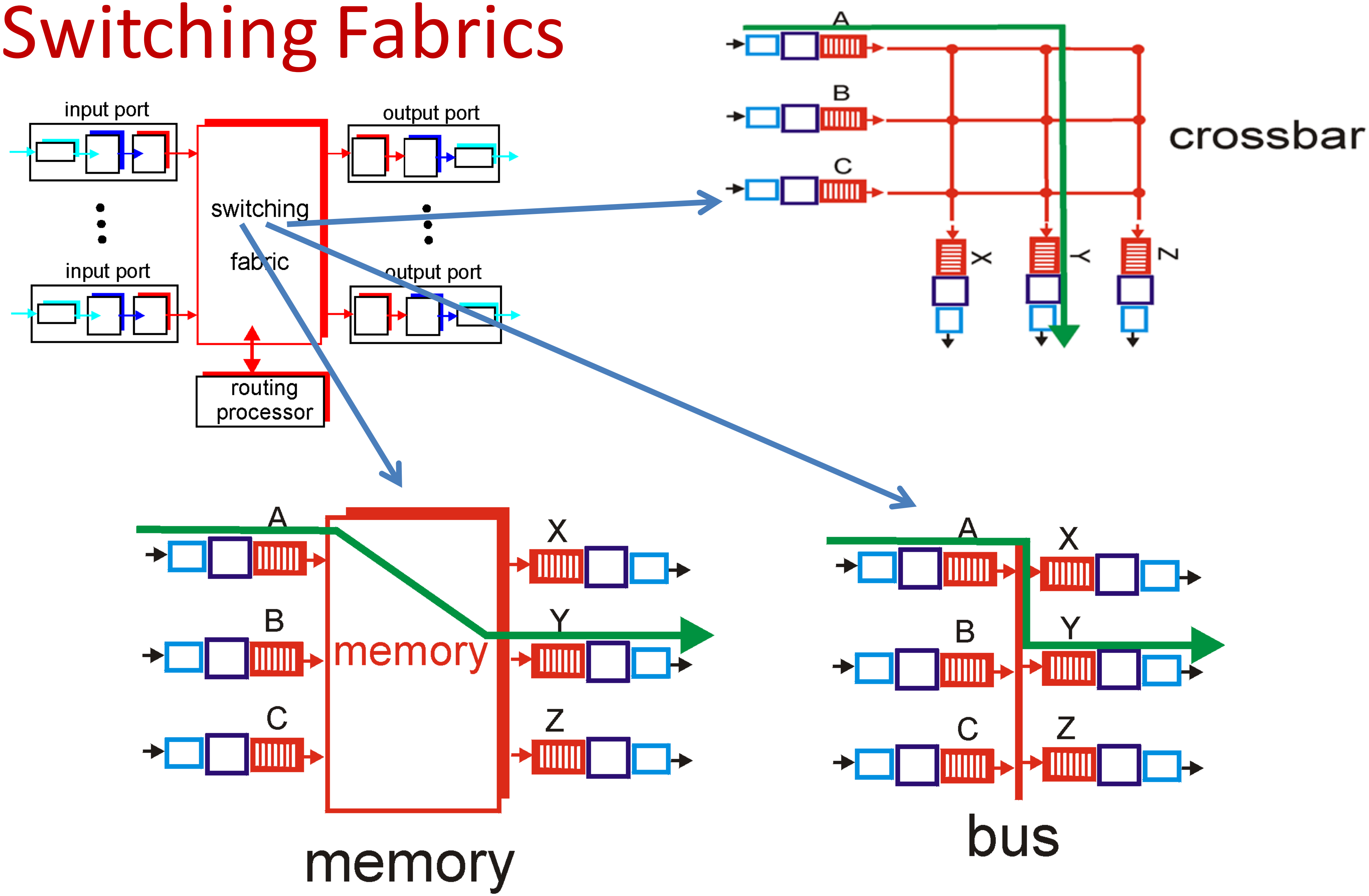
- Scheduling discipline may choose among queued datagrams for transmission

Output Port Queuing

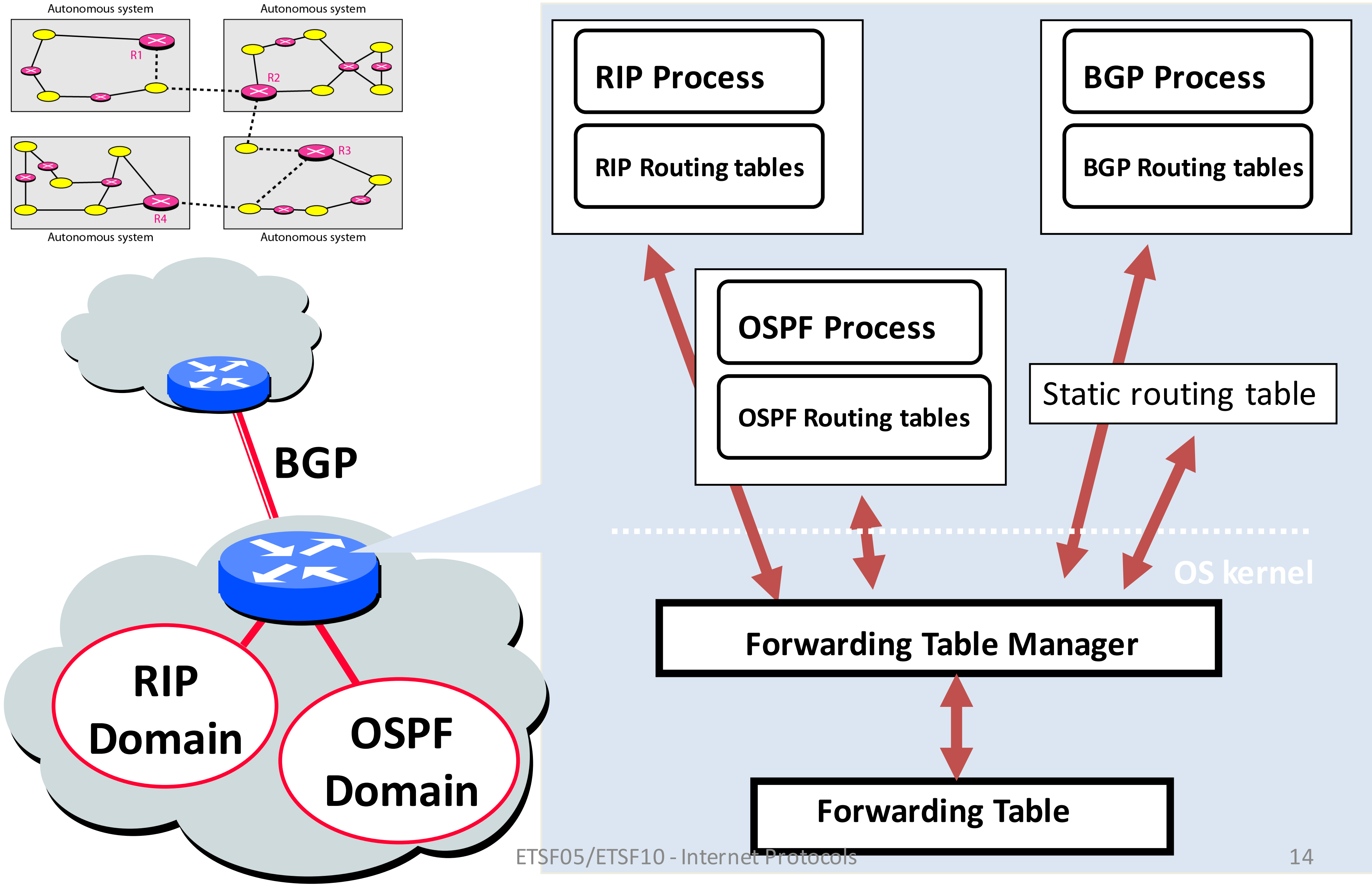
- Datagrams' arrival rate through the switch exceeds the transmission rate of the output line → buffering
- Delay and loss due to output port buffer overflow



Switching Fabrics



Routing Tables and Forwarding Table

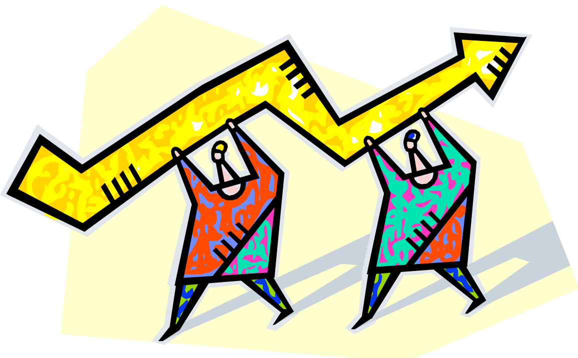


Router cache

- Save next hop for packet type (e.g. addr and TOS)
  - Keep packets within a session on the same path
  - Prohibits reordering
  - decreases delay variations
- Works in both directions
  - Reply take the same path as request
- Drawback: for long sessions (e.g. video) session continuity might be broken if link fails (e.g. mobility)
- Typical for user networks

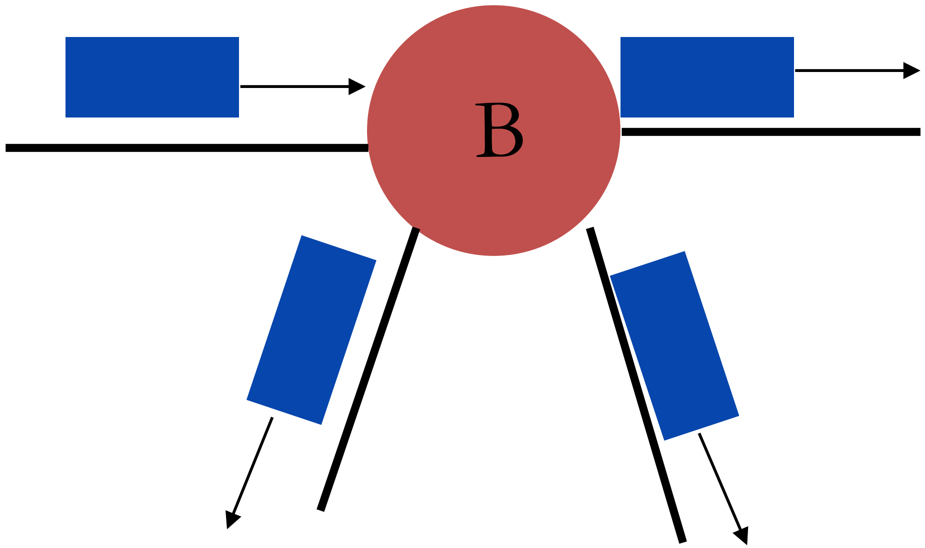
Performance Criteria

- Used for selection of route
- Simplest is to choose "**minimum hop**"
- Can be generalized as "**least cost**" routing
- Because "least cost" is more flexible it is more common than "minimum hop"



Flooding

- In Flooding an incomming packet is retransmitted on all outgoing links. A hop counter is used to prevent loops
- What are the alternatives to find the least cost path.



Best Path: Decision Time and Place

Decision time (when?)

- Packet or virtual circuit (session) basis
- Fixed or dynamically changing

Decision place (where?)

- Distributed - made by each node
  - More complex, but more robust
- Centralized – made by a designated node
- Source – made by source station

Network Information Source and Update Timing

- Routing decisions usually based on knowledge of network, traffic load, and link cost
  - Distributed routing
    - Using local knowledge, information from adjacent nodes, information from all nodes on a potential route
  - Central routing
    - Collect information from all nodes

Issue of update timing

- Depends on routing strategy
- Fixed - never updated
- Adaptive - regular updates

Routing Strategies - Fixed Routing

- Use a **single permanent** route for each source to destination pair of nodes
- Determined using a least cost algorithm
- **Route is fixed**
  - Until a change in network topology
  - Based on expected traffic or capacity
- Advantage is **simplicity**
- Disadvantage is **lack of flexibility**
  - Does not react to network failure or congestion

Routing Strategies - Adaptive Routing

- Used by almost all packet switching networks
- **Routing decisions change as conditions on the network change due to failure or congestion**
- **Requires information about network**

Disadvantages

- More complex
- Tradeoff between quality and overhead
- Too quick updates may lead to oscillations
- Too slow updates may lead to outdates information

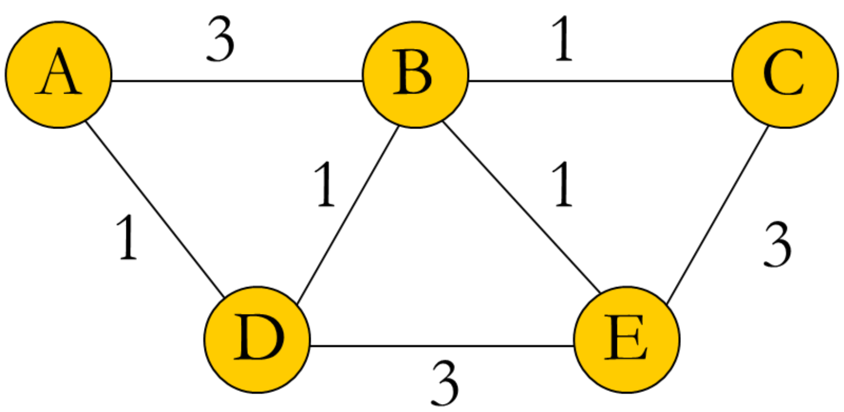
Link cost

- A *cost function* describes the cost for transmitting a packet over a link
- The link cost can depend on e.g.
  - Data rate
  - Load
  - Length
  - Transmission media
  - etc

Graf

A network can be described by a graph, consisting of nodes (*N*) and adges (*E*) with weights *w(e)*, i.e. costs.

Example (undirected graph)



N={A,B,C,D,E}

E	*w(e)*
AB	3
AD	1
BC	1
BD	1
BE	1
CE	3
DE	3

ARPANET Routing Strategies 1st Generation

Distance Vector Routing

- **1969**
- Distributed adaptive using **estimated delay**
  - Queue length used as estimate of delay
- Version of **Bellman-Ford** algorithm
- **Node exchanges delay vector with neighbors**
- **Update routing table based on incoming information**
- **Doesn't consider line speed**, just queue length and responds slowly to congestion

Least cost alg 1
Bellman-Ford

- Find the shortes path from one source node s to the others.
- Let $d(n)$ be the cost from s to n

Init:

$d(s) = 0$

$d(n) = \infty,\ n \neq s$

for $\mathrm{i} = 1\ to |N| - 1$

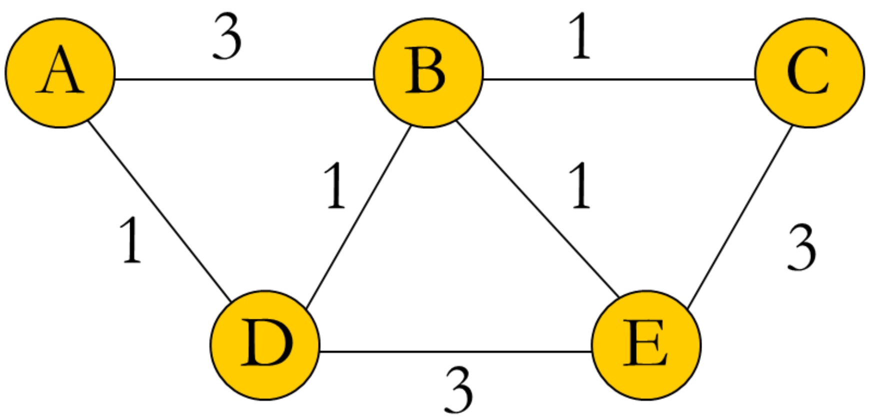
for each $n \in N$

$d(n) = \min_{u \in N}\{d(u) + w(u, n)\}$ // Find the shortest path from

// node u to node n in one step

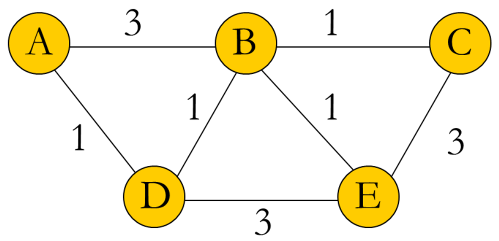
- Addition: Keep track of the path!!

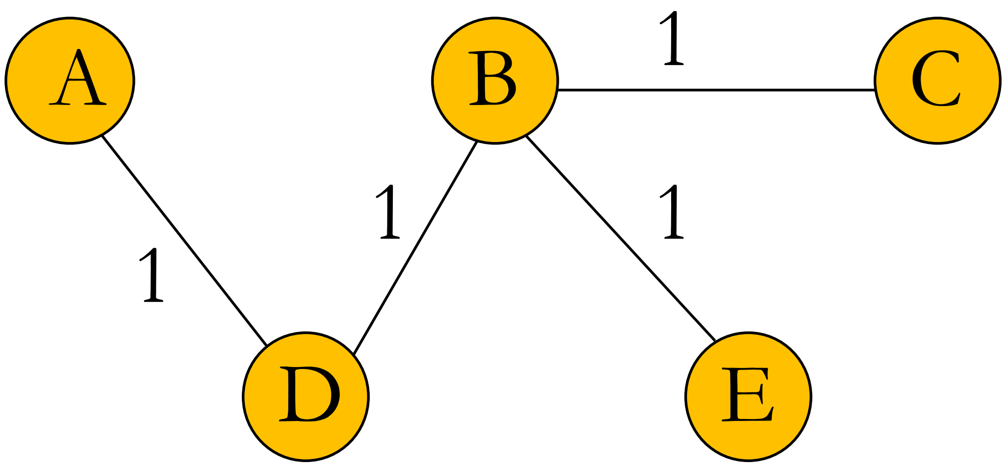
Example Bellman-Ford



Nod	A	B	C	D	E
init	0	∞	∞	∞	∞
i=1	0	3	∞	1	∞
i=2	0	2	4	1	4
i=3	0	2	3	1	3
i=4	0	2	3	1	3

Net graph as a tree

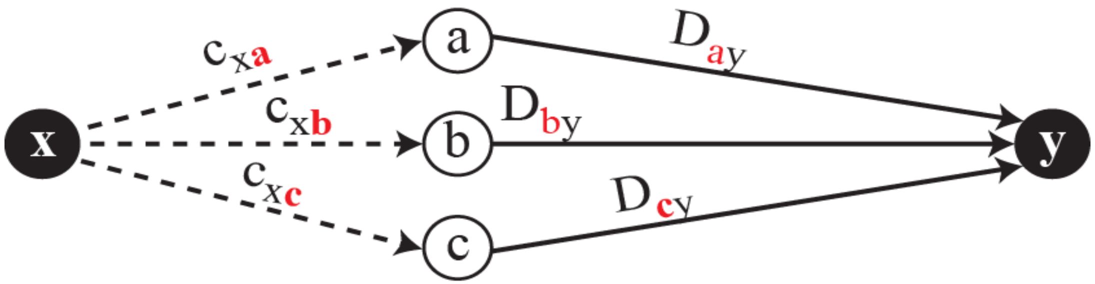




Distant vector for A when the algorithm converged

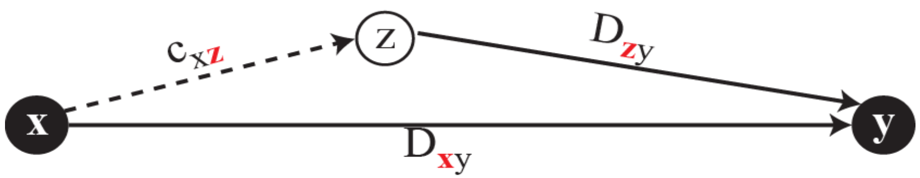
Nod	Dist
A	0
B	2
C	3
D	1
E	3

Bellman-Fords algoritm grafiskt



a. General case with three intermediate nodes

$$D_{xy} = \min\{(c_{xa} + D_{ay}), (c_{xb} + D_{by}), (c_{xc} + D_{cy}) \dots\}$$



b. Updating a path with a new route

$$D_{xy} = min\{D_{xy}, (c_{xz} + D_{zy})\}$$

Jmf Stallings kap 19.2

ARPANET Routing Strategies 2nd Generation

Link-State Routing

- **1979**
- Distributed adaptive using **delay** criterion
  - Using timestamps of arrival, departure and ACK times
- Re-computes average delays every 10 seconds
- **Any changes are flooded to all other nodes**
- Re-computes routing using **Dijkstra's algorithm**
- Good under light and medium loads
- Under heavy loads, little correlation between reported delays and those experienced

Least cost alg 2
Dijkstra

- Find the shortes path from one source node s to the others.
- Let $d(n)$ be the cost from s to n

Init:

$d(s) = 0$

$d(n) = \infty, \; n \neq s$

$V = \emptyset$ // Visited nodes

while $V \subset E$

$u = arg \min_{u \notin V} d(u)$ // Find least weight to unvisted node

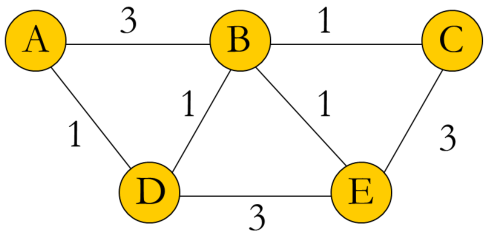
$V = V \cup u$ // Add u to visited nodes

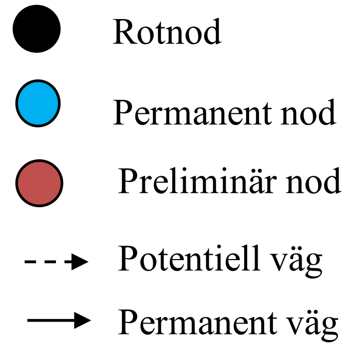
for $n \notin V$

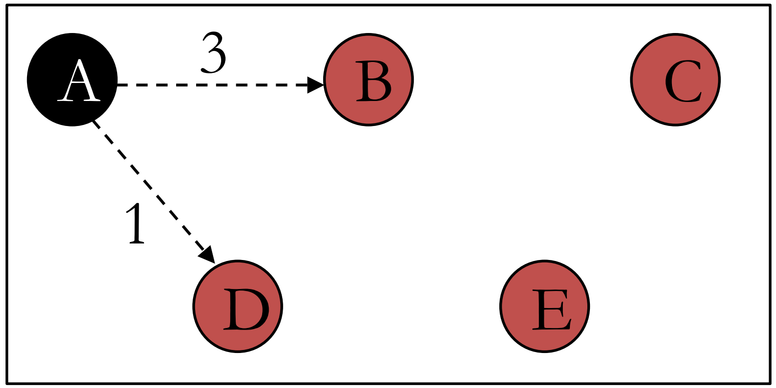
$d(n) = \min\{d(n), d(u) + w(u, n)\}$ // Less cost to go to n via u?

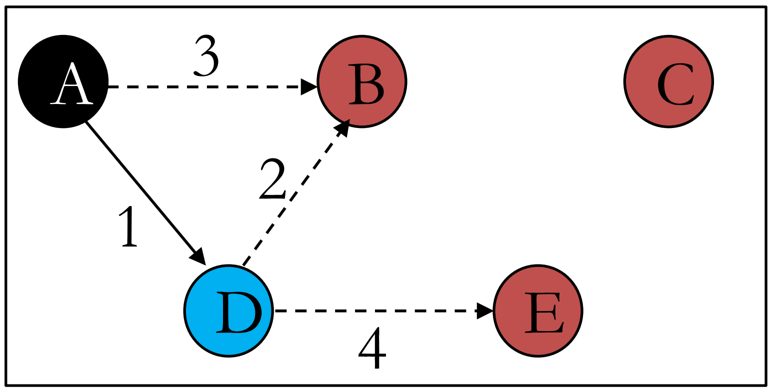
- Addition: Keep track of the path!!

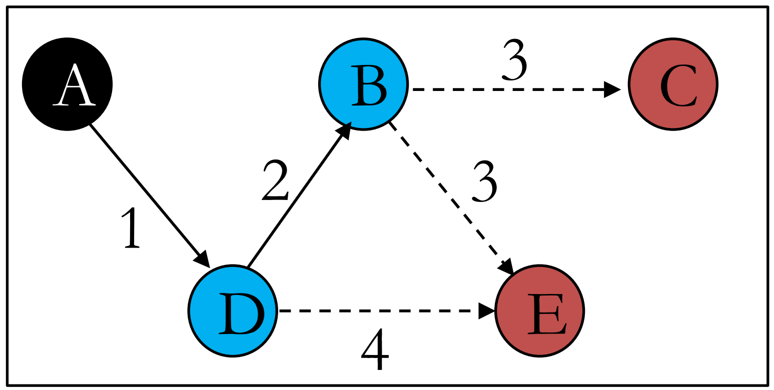
SPF exempel med graf

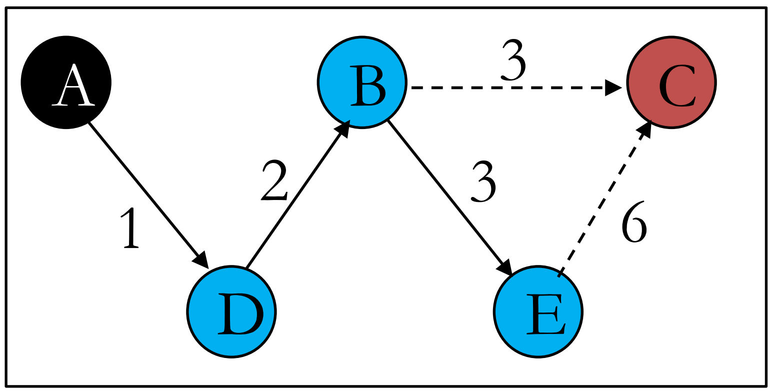


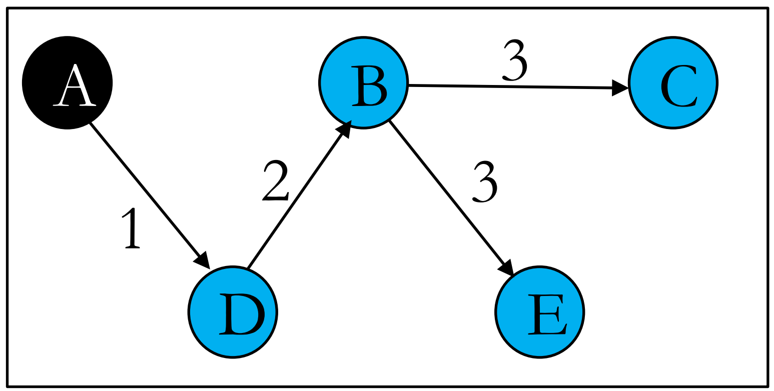












Dijkstra tabell

Besökt	L(A)	L(B)	L(C)	L(D)	L(E)
ϕ	**0**	∞	∞	∞	∞
{A}		3:A	∞	**1:A**	∞
{A,D}		**2:D**	∞		4:D
{A,D,B}			**3:B**		3:B
{A,D,B,C}					**3:B**
{A,D,B,C,E}					

ARPANET Routing Strategies 3rd Generation

- **1987**
- Link cost calculation changed
  - Dampen routing oscillations
  - Reduce routing overhead
- Measure average delay over last 10 seconds and transform into link utilization estimate
- Calculate average utilization based on current value and previous average

$$U(n+1) = \frac{1}{2}\rho(n) + \frac{1}{2}\ U(n)$$

- Use as link cost a function based n the average utilization

Autonomous Systems (AS)

- Exhibits the following characteristics:
  - Is a set of routers and networks managed by a single organization
  - Consists of a group of routers exchanging information via a common routing protocol
  - Except in times of failure, is connected (in a graph-theoretic sense); there is a path between any pair of nodes

Interior Router Protocol (IRP) Interior Gateway Protocol (IGP)

- Shared routing protocols passes routing information between routers **within an AS**
- Custom tailored to specific applications and requirements
- Examples:
  - Routing Information Protocol (RIP)
  - Open Shortest Path First (OSPF)

Exterior Router Protocol (ERP) Exterior Gateway Protocol (EGP)

- Protocol used to pass routing information **between routers in different ASs**
- Will need to pass less information than an IRP
  - To transmit a datagram from a host in one AS to a host in another AS, a router in the first system need only determine the target AS and devise a route to get into it
  - Once the datagram enters the target AS, the routers within that system can cooperate to deliver the datagram
  - The ERP is not concerned with details of the route
- Examples:
  - Border Gateway Protocol (BGP)
  - Open Shortest Path First (OSPF)

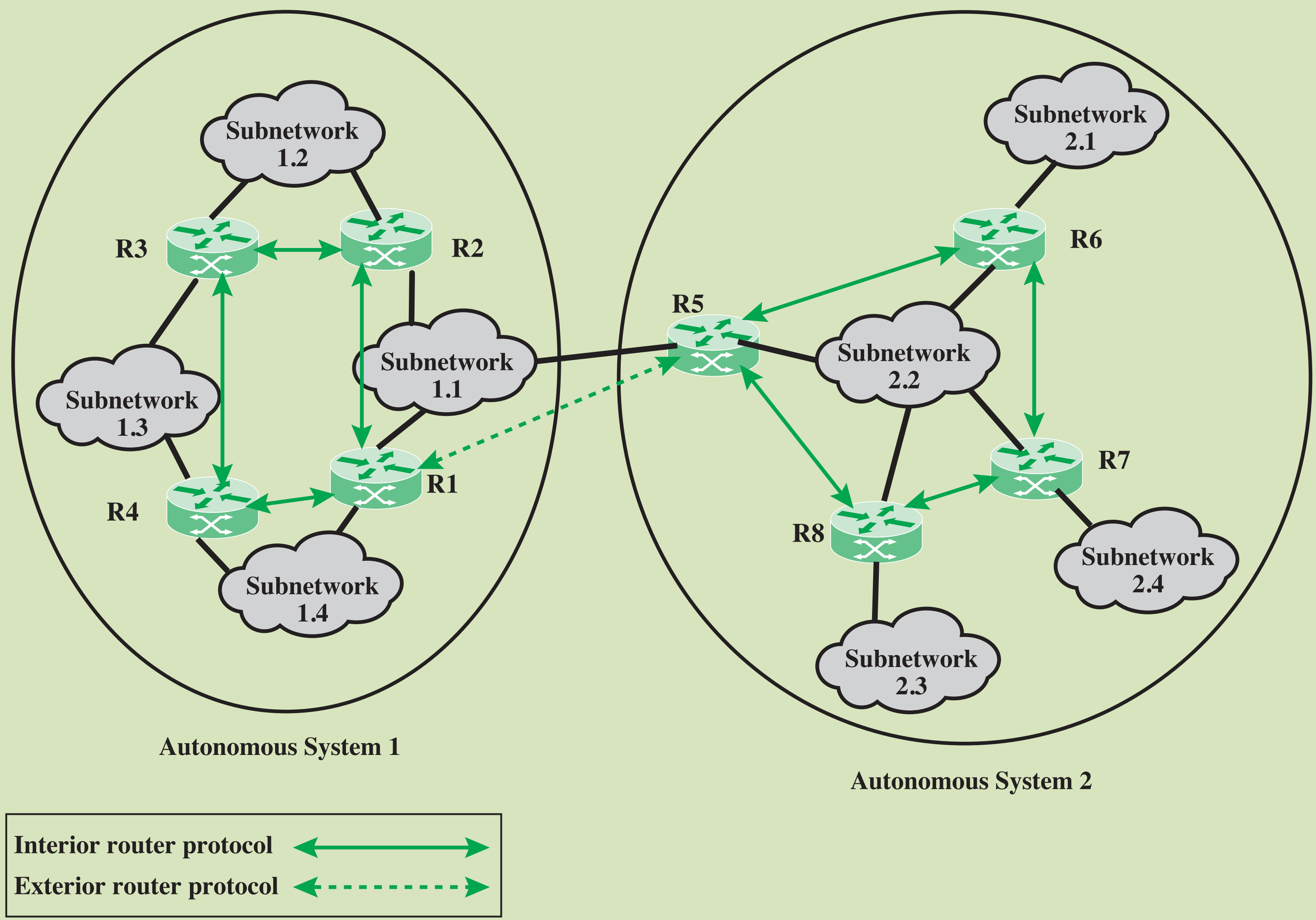
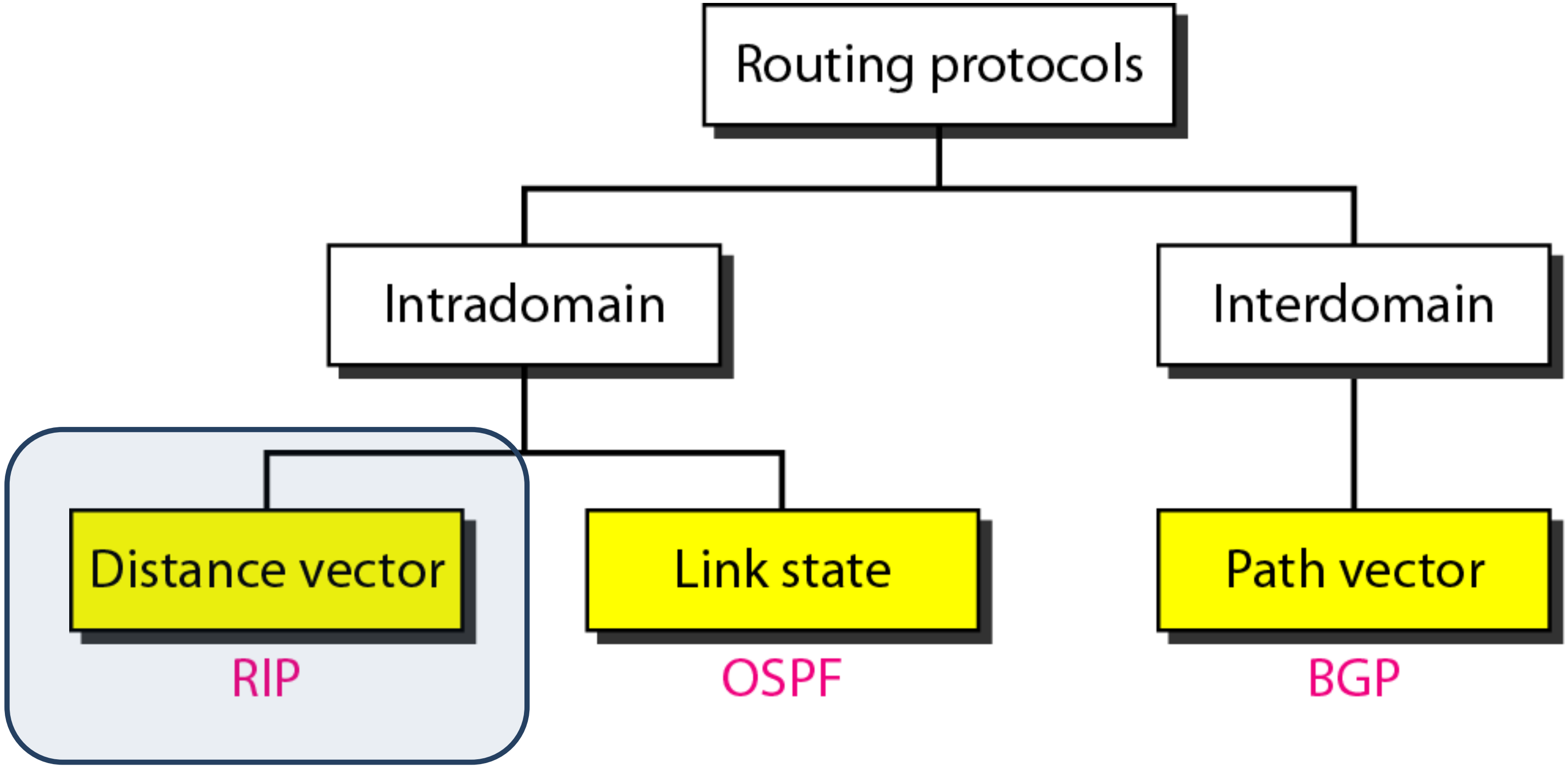


Figure 19.9 Application of Exterior and Interior Routing Protocols

Routing Algorithms and Protocols



Distance-Vector Routing

- Requires that each node exchange information with its neighboring nodes
  - Two nodes are said to be neighbors if they are both directly connected to the same network
- Used in the first-generation routing algorithm for ARPANET
- Each node maintains a vector of link costs for each directly attached network and distance and next-hop vectors for each destination
- Routing Information Protocol (RIP) uses this approach

RIP (Routing Information Protocol)

- Included in BSD-UNIX Distribution in 1982
- Distance metric:
  - **# of hops** (max 15) to destination network
- Distance vectors:
  - exchanged among neighbours every 30 second via Response Message (advertisement)
- Implementation:
  - Application layer protocol, uses UDP/IP

RIP update message

- Contains the whole forwarding table
- Action on reception:
  - Add 1 to cost in received message
  - Change next hop to sending router
  - Apply RIP updating algorithm

- Received update msgs identify neighbours!

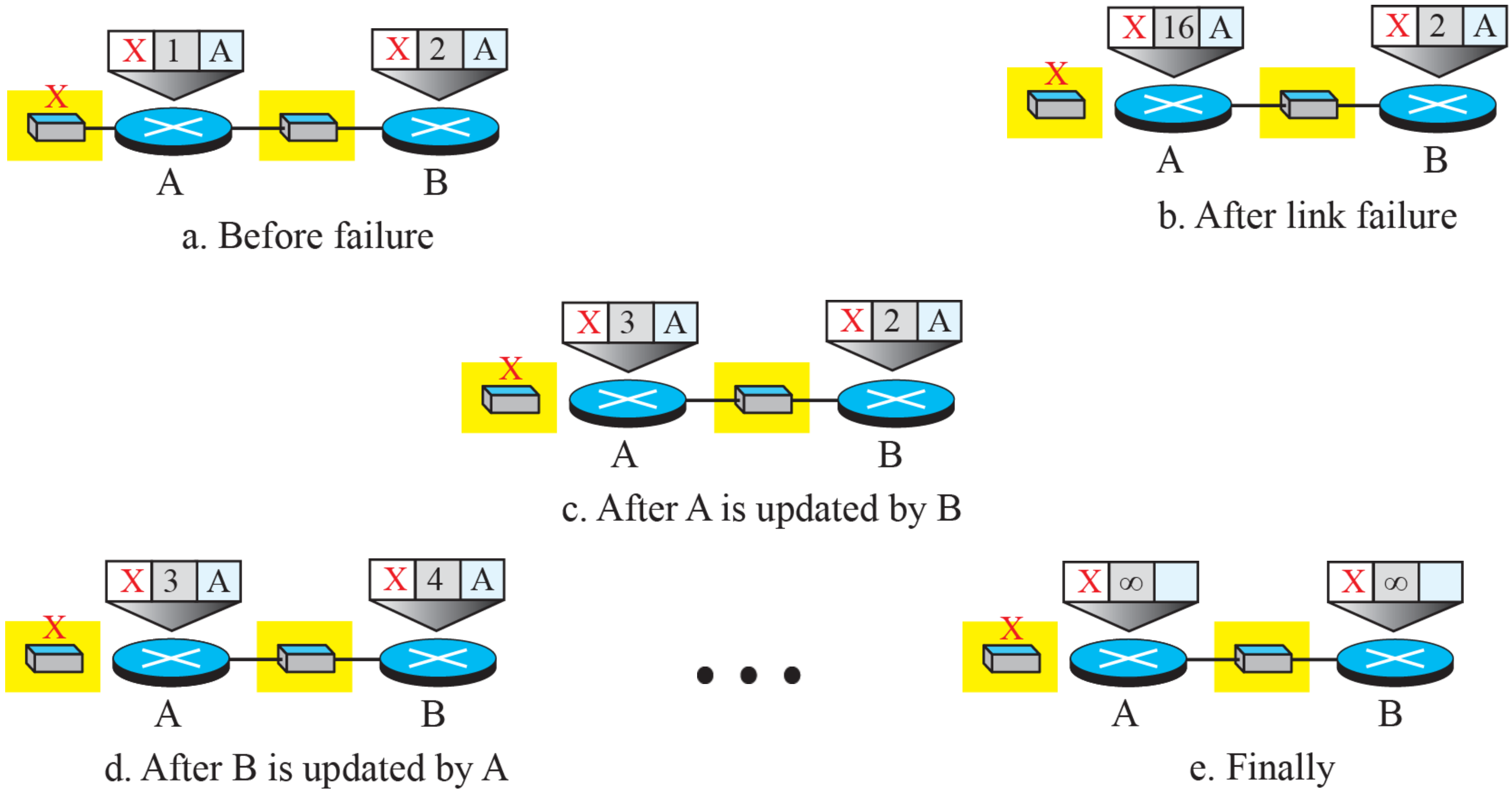
RIP Updating Algorithm (Bellman-Ford)

```
if (advertised destination not in table)
    {
    add new entry // rule #1
    }
else if (adv. next hop = next hop in table)
    {
    update cost // rule #2
    }
else if (adv. cost < cost in table)
    {
    replace old entry // rule #3
    }
```

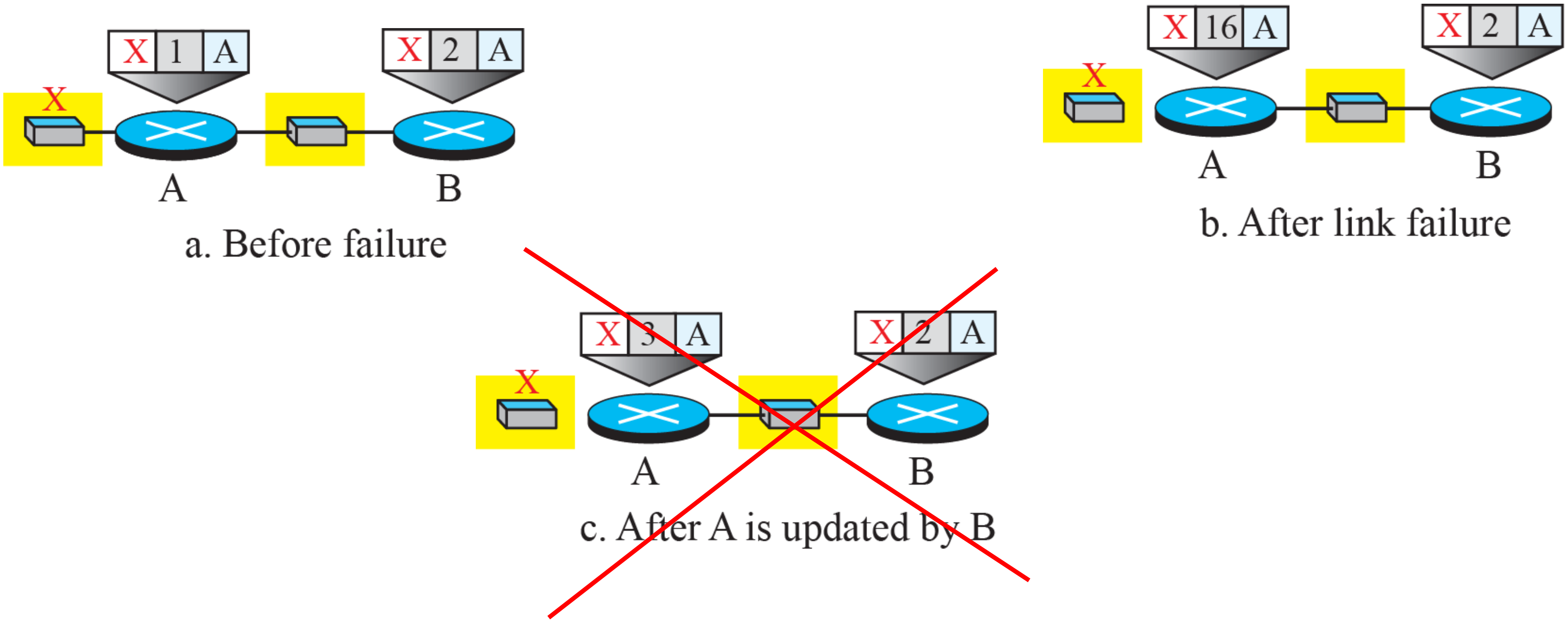
RIP Example


New R1
Des.
N. R.
Cost
N1
1
N2
1
N3
1
N4
R2
2
N5
R2
2
Old R1
Des.
N. R.
Cost
N1
1
N2
1
N3
1
R2 Seen by R1
Des.
N. R.
Cost
N3
R2
2
N4
R2
2
N5
R2
2

Two node instability/Count to inifinity



Split Horizon breaks Count to inifinity

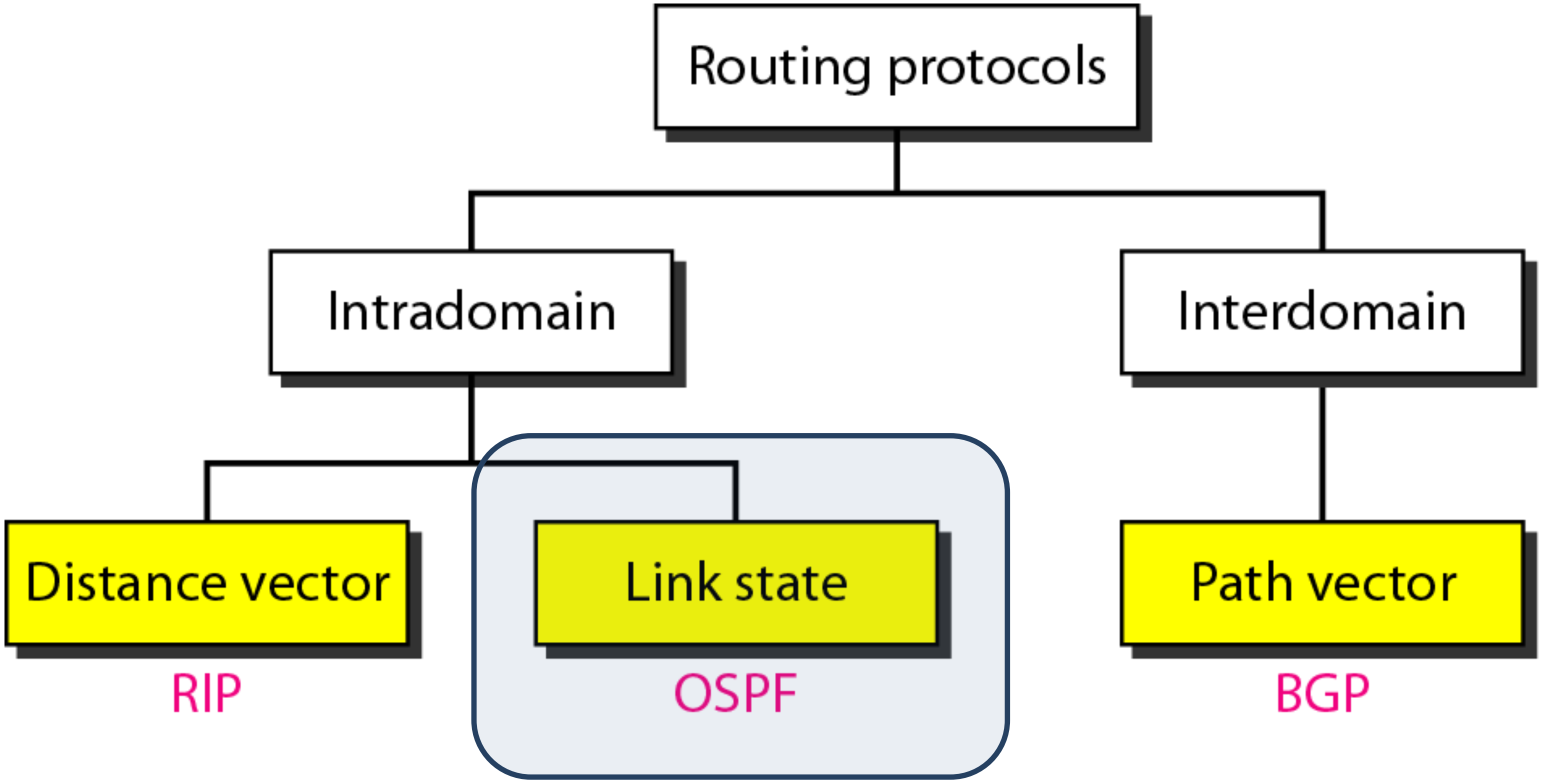


I have a route to X, but I got it from A so I won't tell A about it!

RIP: Link Failure and Recovery

- If no advertisement heard after 180 seconds
  - Neighbour/link declared dead
  - Routes via neighbour invalidated (infinite distance = 16 hops)
  - New advertisements sent to neighbours (triggering a chain reaction if tables changed)
  - “Poison reverse” used to prevent count to infinity loops
  - “Good news travel fast, bad news travel slow”

Routing Algorithms and Protocols



Link-State Routing

- When a router is initialized, it determines the link cost on each of its network interfaces
- The router then advertises this set of link costs to all other routers in the internet topology, not just neighboring routers
- From then on, the router monitors its link costs
- When there is a significant change, the router advertises its link costs to all other routers

- The OSPF protocol is an example
- The second-generation routing algorithm for ARPANET also uses this approach

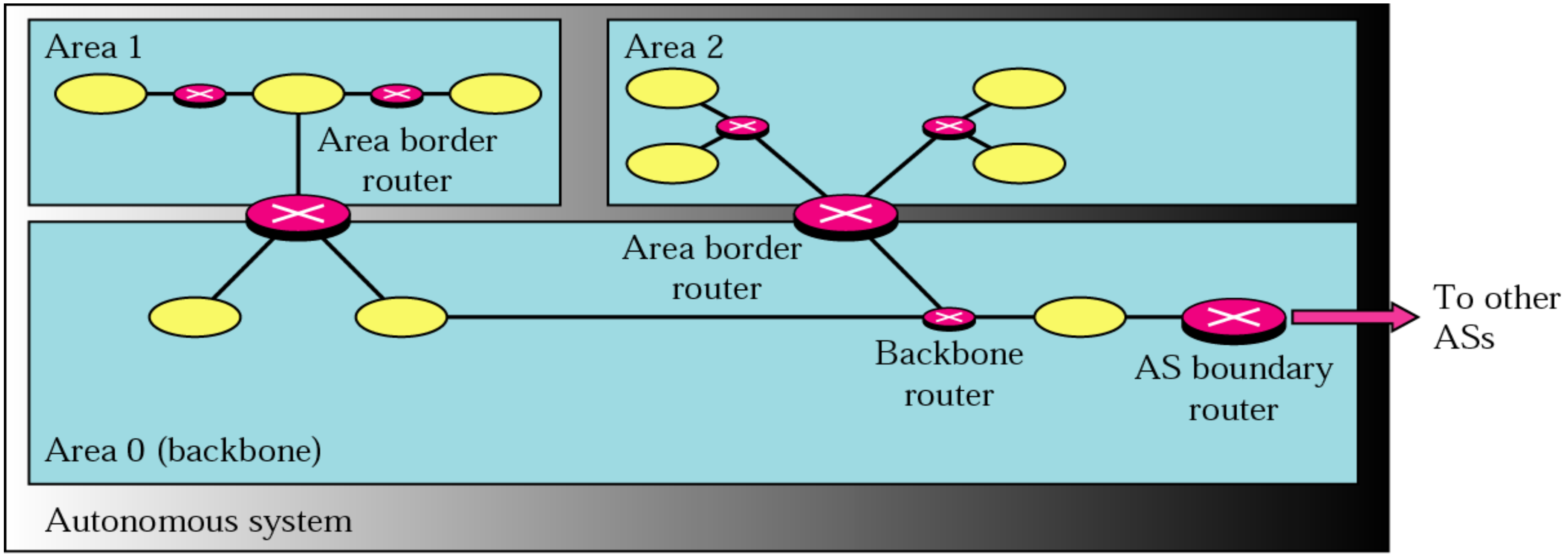
Open Shortest Path First (OSPF) Protocol

- RFC 2328 (Request For Comments)
- Used as the interior router protocol in TCP/IP networks
- Computes a route that incurs the least cost based on a user-configurable metric
- Is able to balance loads over multiple equal-cost paths

OSPF (Open Shortest Path First)

- Divides domain into areas
  - Limits flooding for efficiency
  - One ”backbone” area connects all
- Distance metric:
  - Cost to destination network

Areas, Router and Link Types



Graph

Network topology expressed as a graph

- Routers
- Networks
  - Transit, passing data through
  - Stub, not transit
- Edges
  - Direct, router to router
  - Indirect, router to network

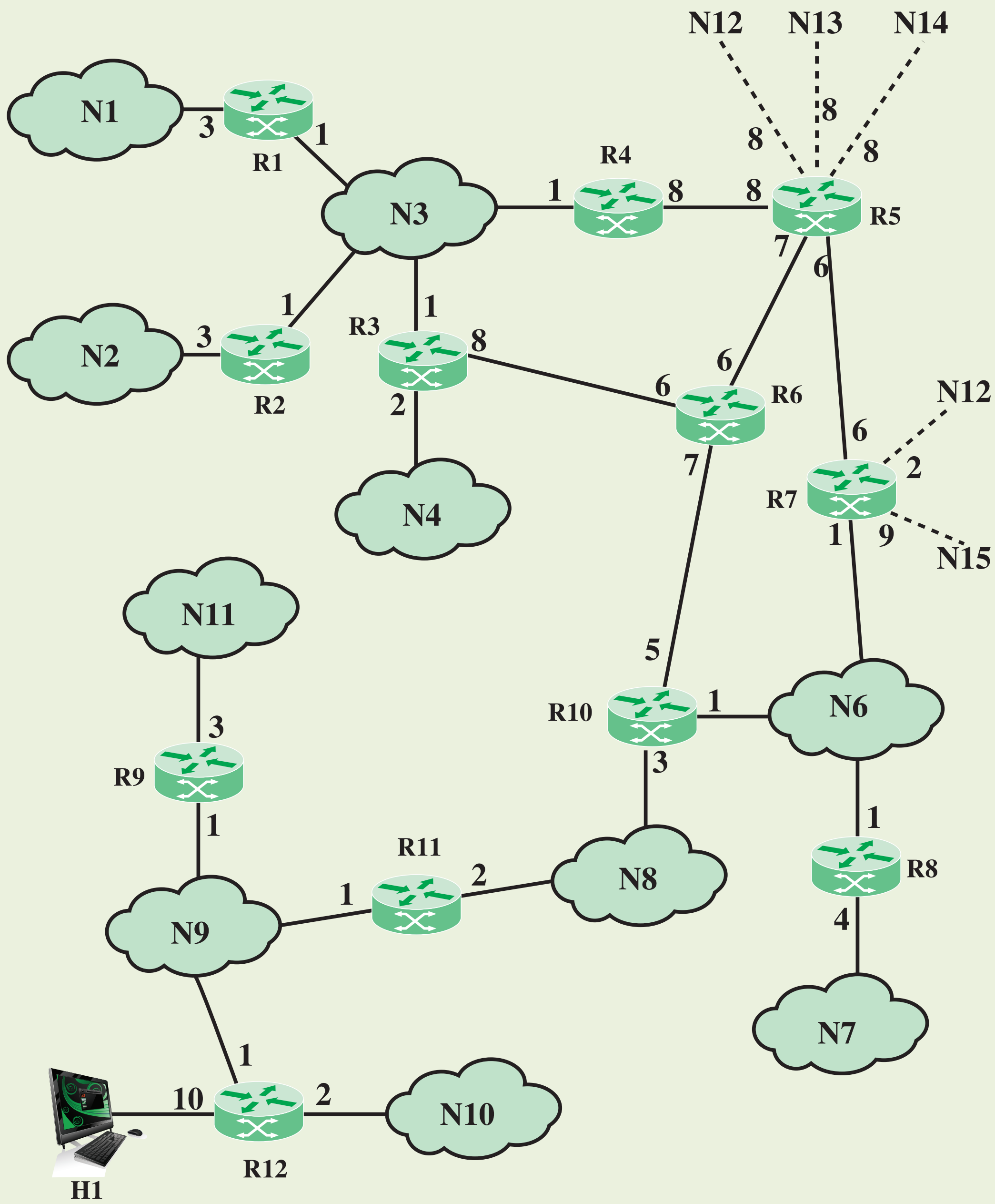


Figure 19.11 A Sample Autonomous System

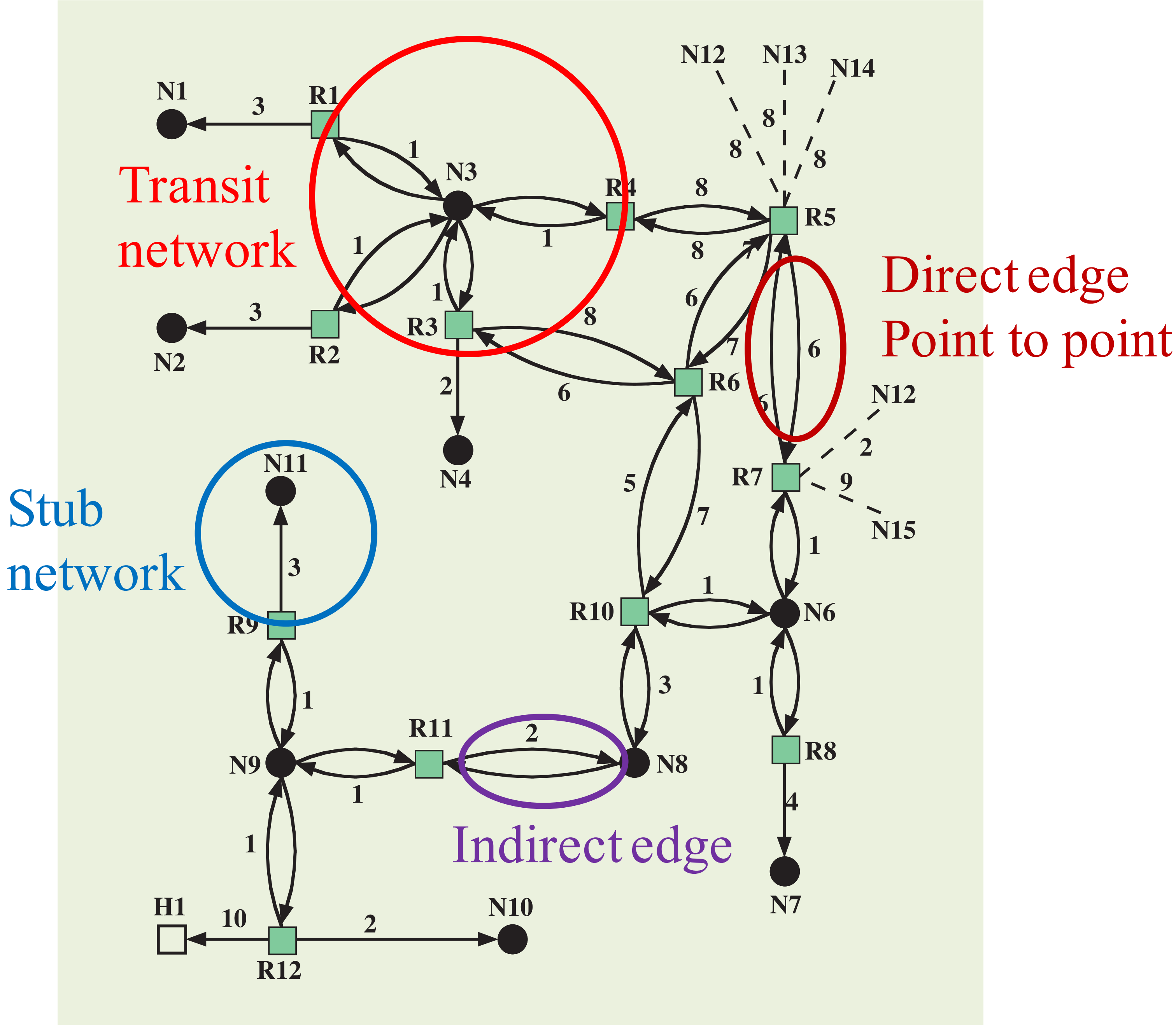
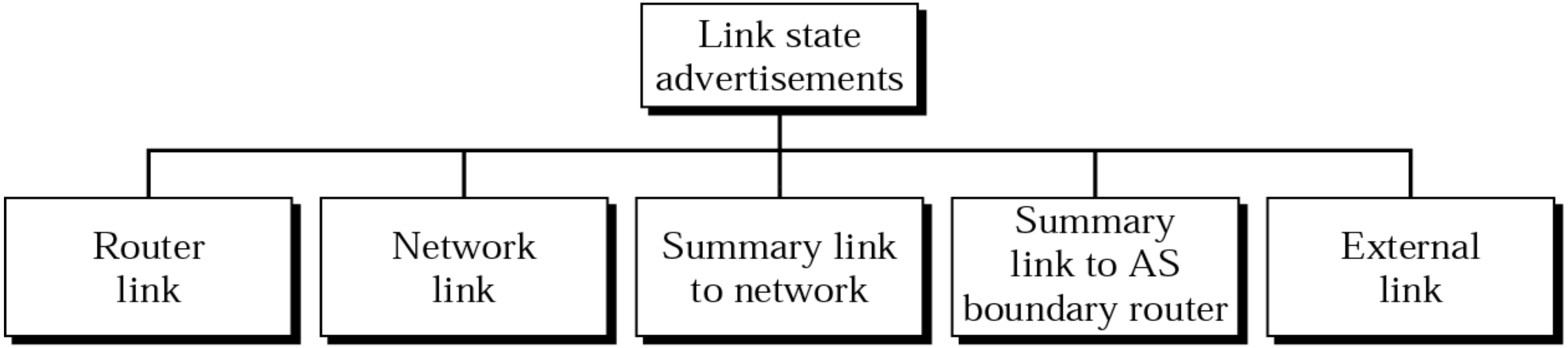


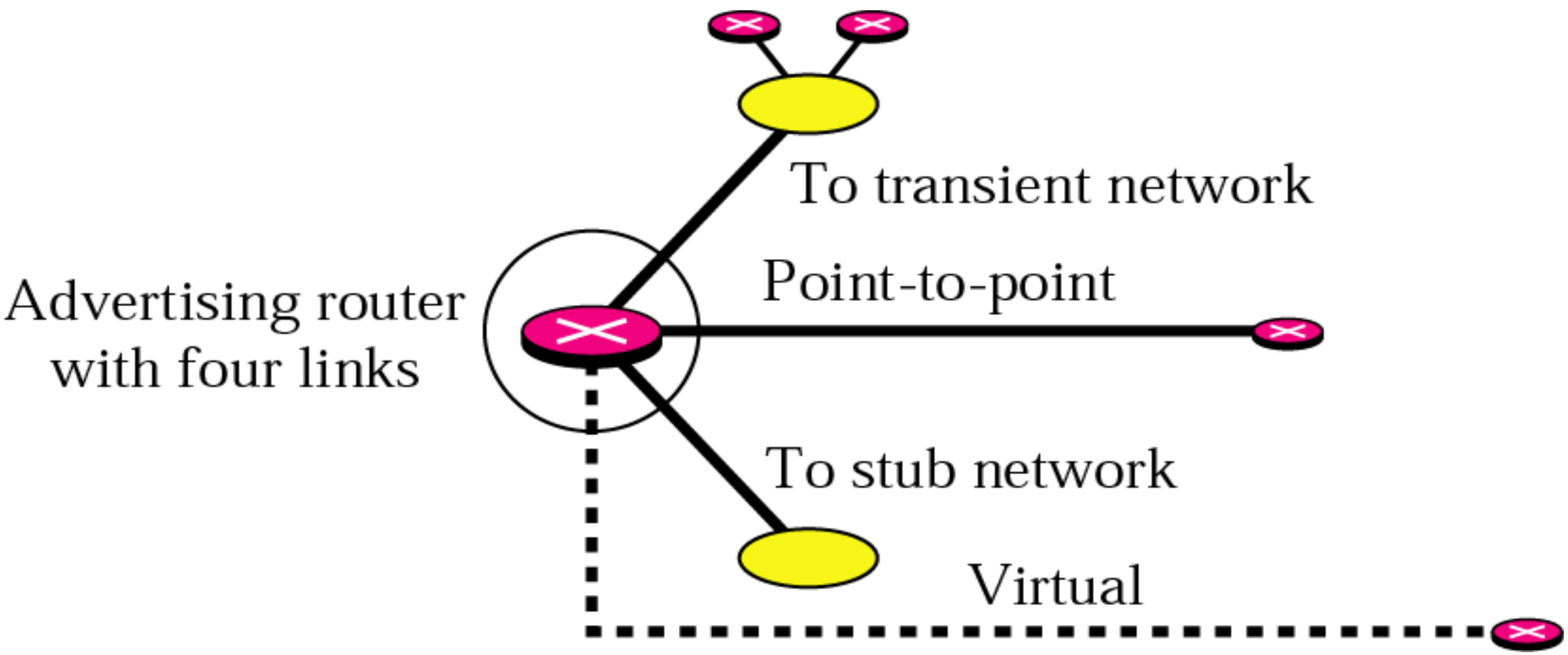
Figure 19.12 Directed Graph of Autonomous System of Figure 19.11

Link State Advertisements

- What to advertise?
  - Different entities as nodes
  - Different link types as connections
  - Different types of cost

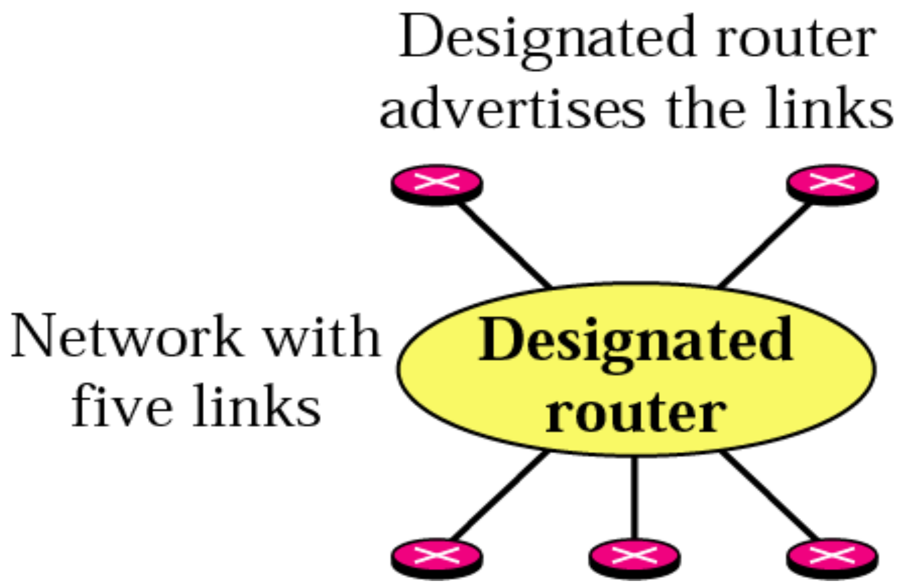


Router Link Advertisement



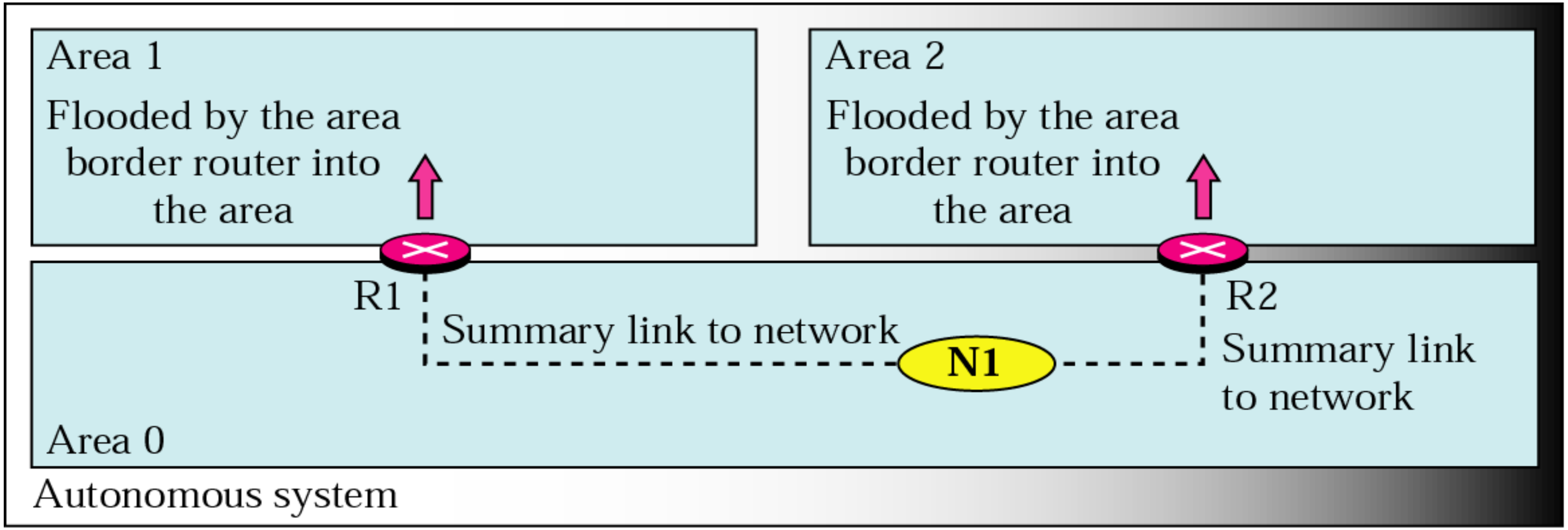
Network Link Advertisement

- Network is a passive entity
  - It cannot advertise itself



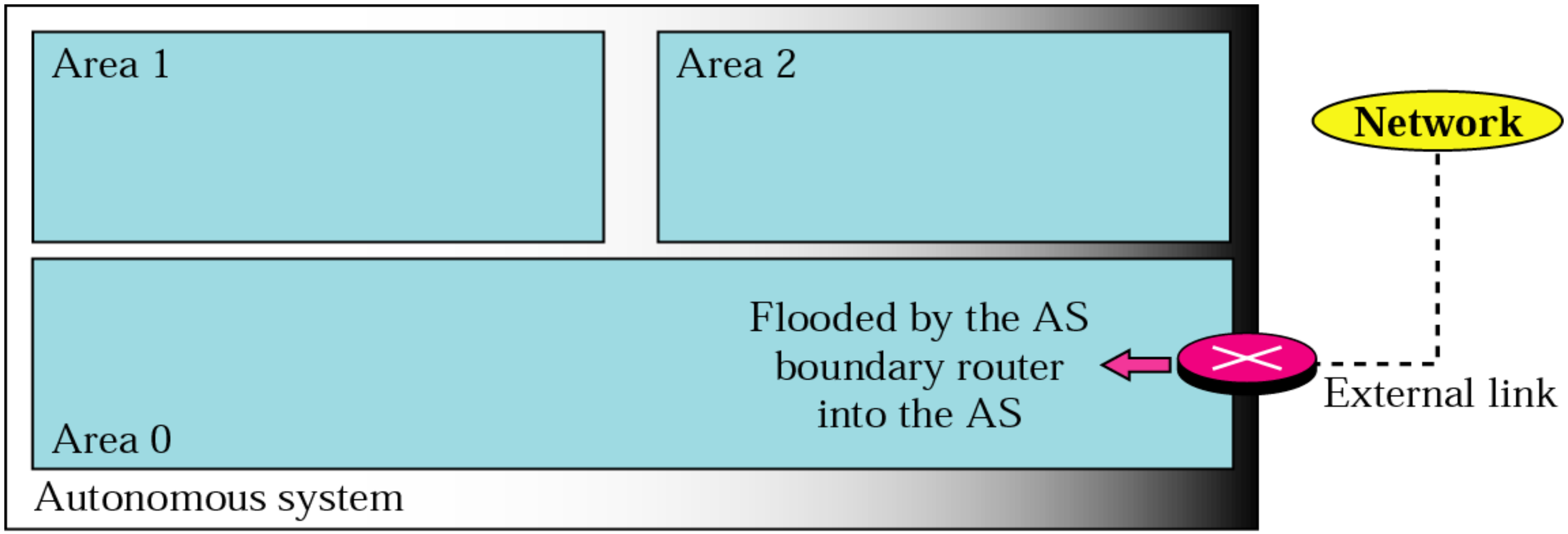
Summary Link to Network

- Done by area border routers
  - Goes through the backbone



External Link Advertisement

- Link to a single network outside the domain



Hello message

- Find neighbours
- Keep contact with neighbours: I am still alive!
- Sent out periodically, typically every 10 second
- If no hellos received during holdtime (typically 30 seconds), neighbour declared dead.

- Compare RIP update messages

Routing Algorithms and Protocols

- Interior and Exterior Router Protocols
- Distance vector
  - Bellman-Ford
  - Announce whole table to neighbors
  - RIP
- Link State
  - Dijkstra
  - Announce neighbor connections to whole network
  - OSPF