

ETSF05/ETSF10 – Internet Protocols

SMTP

FTP

TFTP

DNS

SNMP

...

BOOTP

SCTP

TCP

UDP

Routing on the Internet

IGMP

ICMP

IP

ARP

RARP

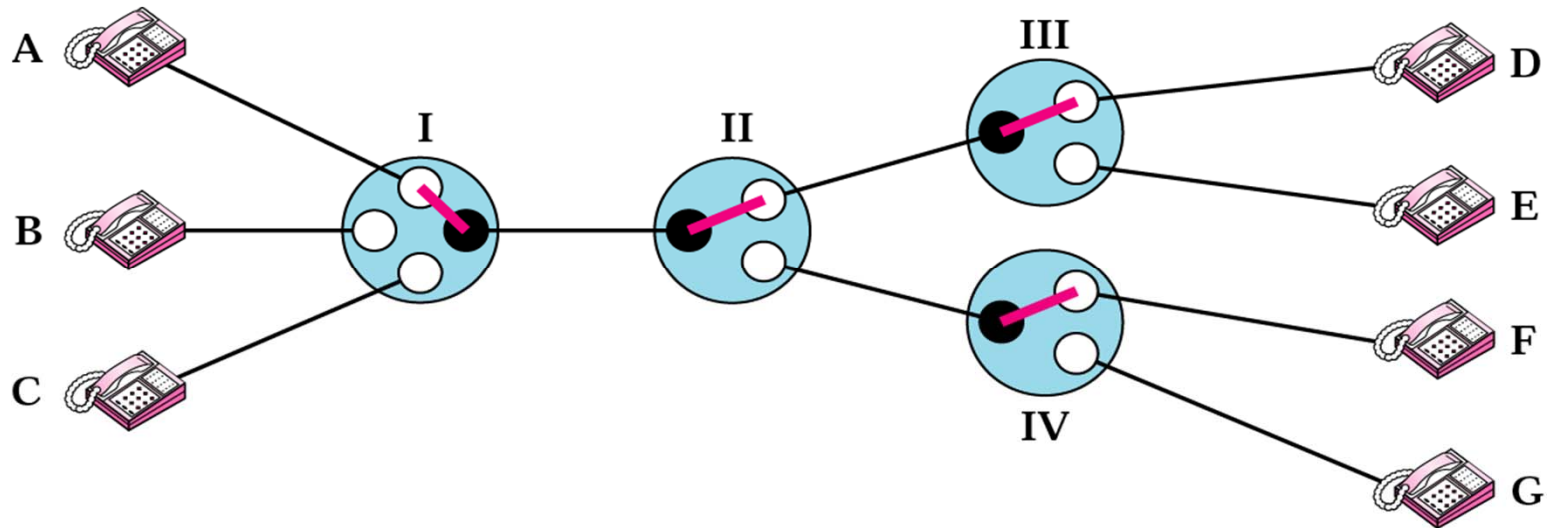
2014 extra föreläsning unicast routing

Underlying LAN or WAN
technology

Jens Andersson

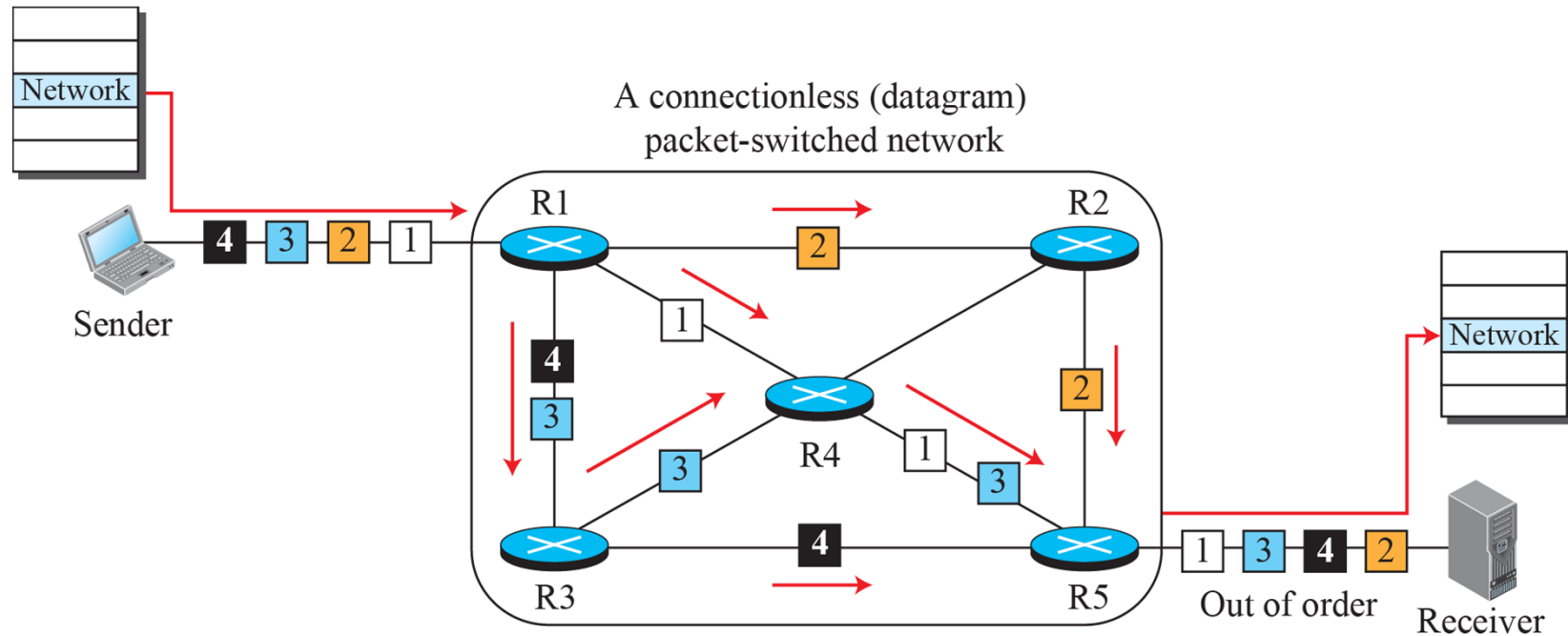


Circuit switched routing



Packet-switched Routing

- Choosing an optimal path
 - According to a cost metric
 - Decentralised: each router has full information

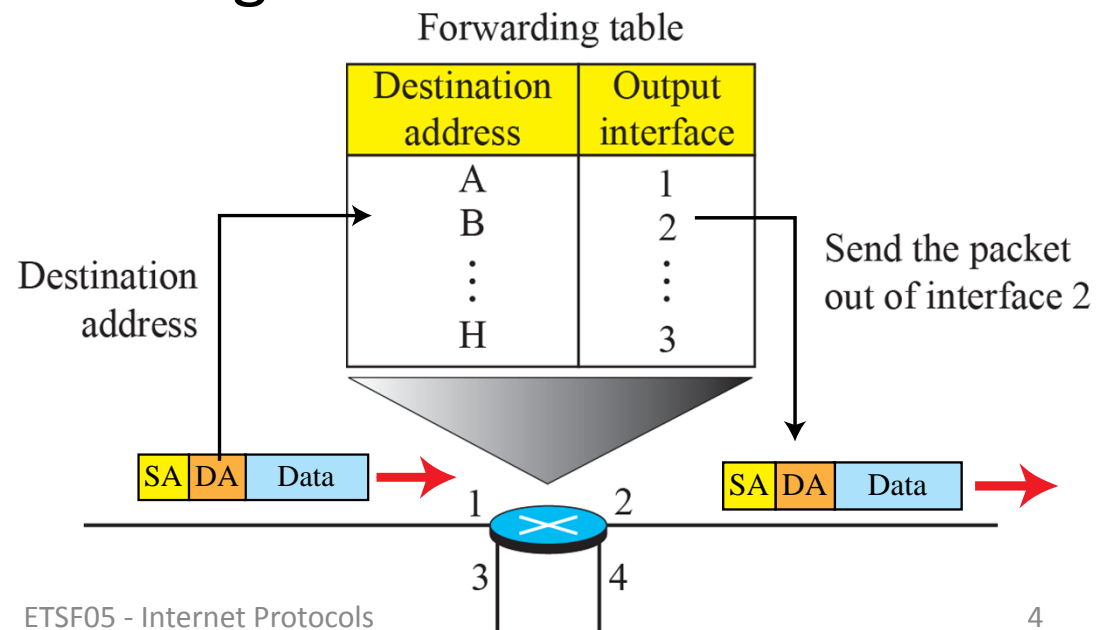


Router

- Internetworking device
 - Passes data packets between networks
 - Checks **Network Layer** addresses
 - Uses Routing/forwarding tables

Two functions:

- 1 Routing
- 2 Forwarding

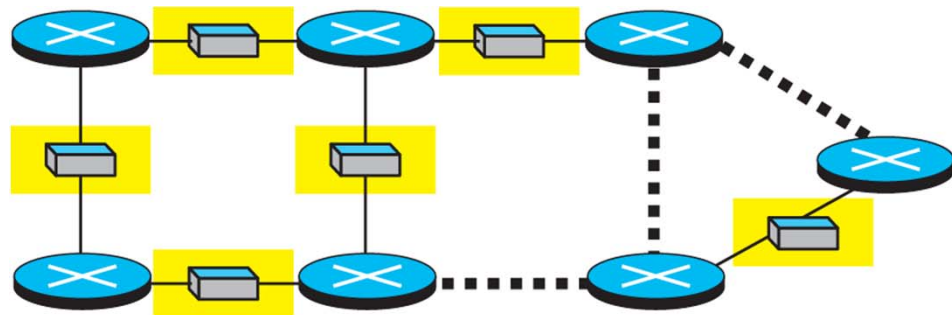
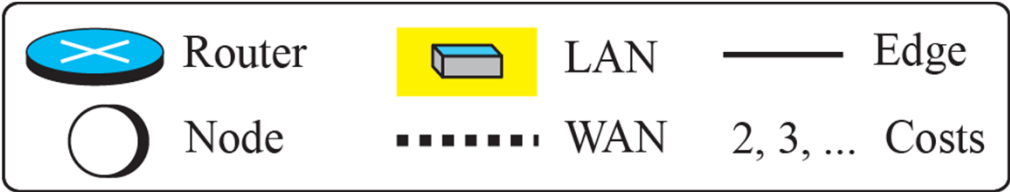


Common Cost Metrics

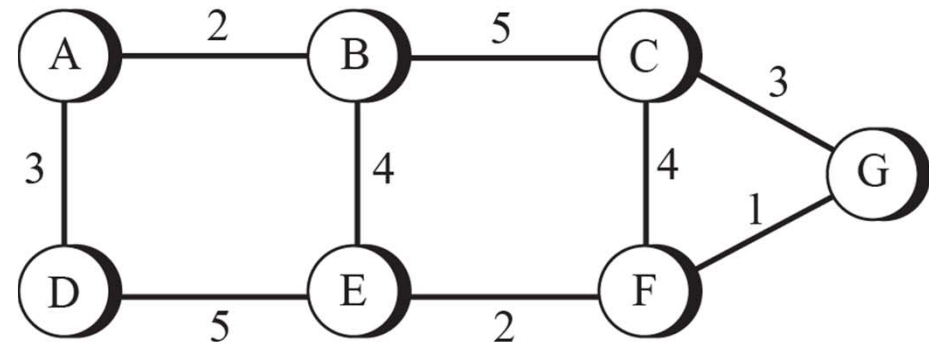
- Alternatives at the link level
 - Hop count
 - Inverse of the link bandwidth
 - Delay
 - Dynamically calculated
 - Administratively assigned
 - Combination
- Traffic monitored → metrics adjusted

Graphical representation of a net

Legend

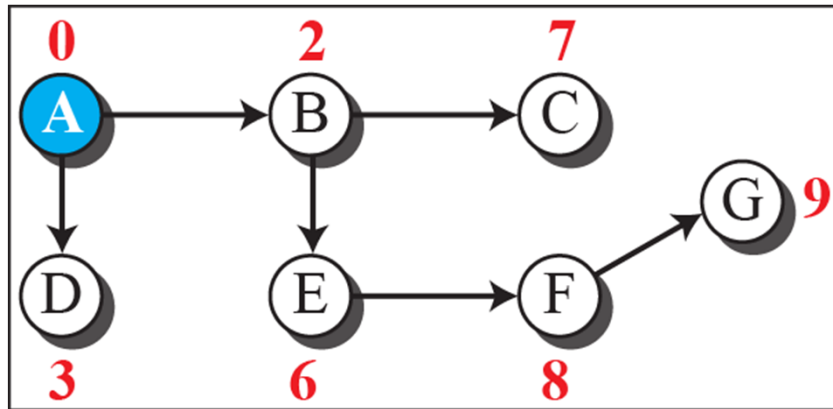


a. An internet





b. The weighted graph

What is an end node?



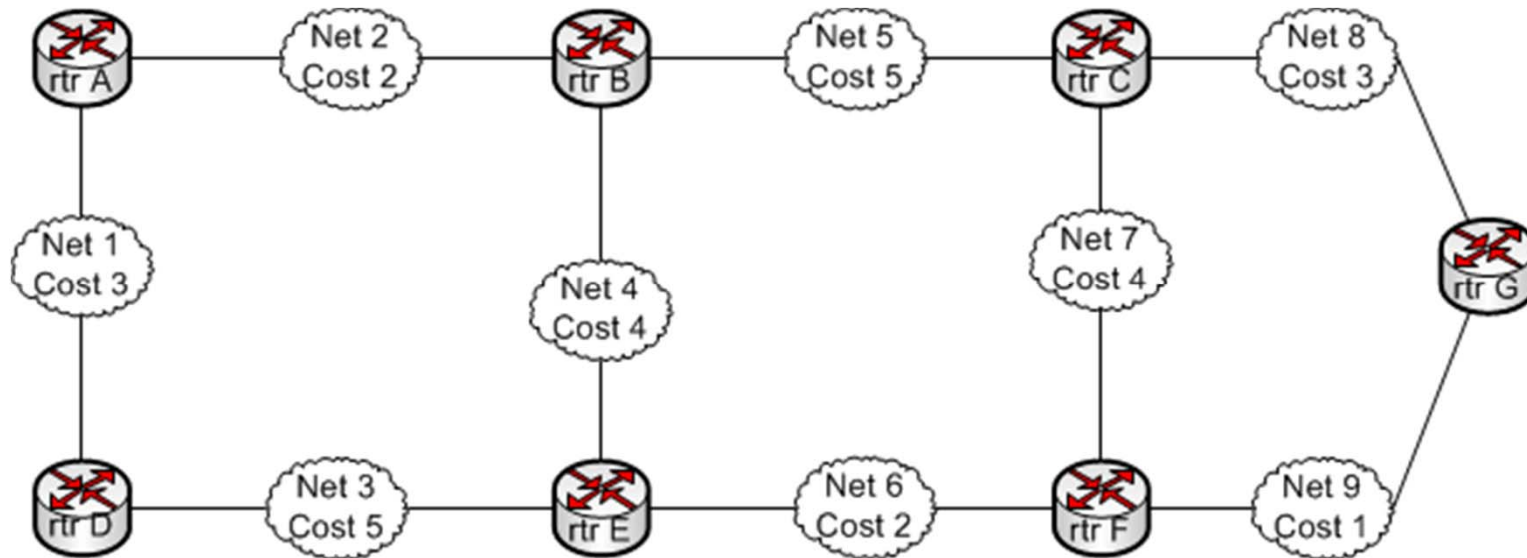
Legend

-  Root of the tree
-  Intermediate or end node
- 1, 2, ...** Total cost from the root

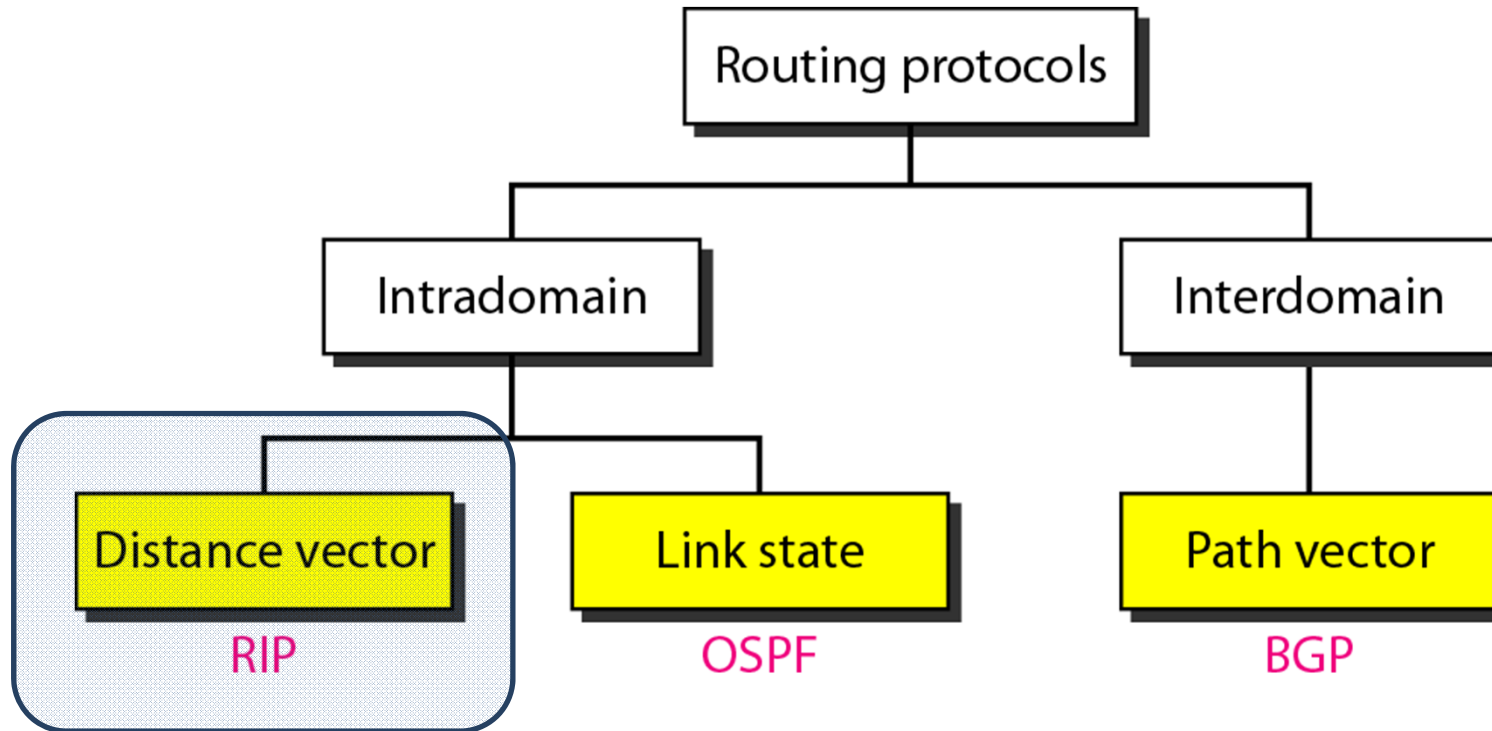
Problem: The LANs are our destinations/end nodes, not the routers

A more realistic representation

- Solution: Nets and routers are all nodes in the tree.
- Routers hold tables how to reach nets and what is the *next hop* for to get there



Routing Algorithms and Protocols



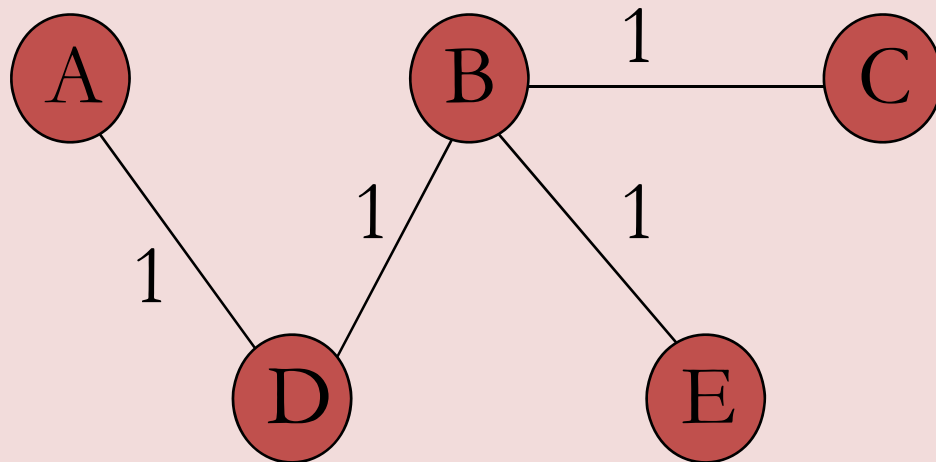
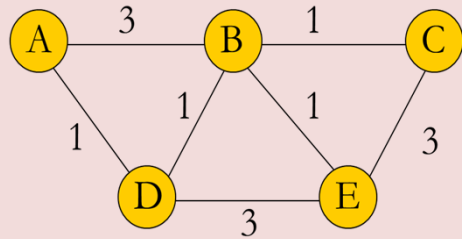
RIP (Routing Information Protocol)

- Included in BSD-UNIX Distribution in 1982
- Distance metric:
 - **# of hops** (max 15) to destination network
- Distance vectors:
 - exchanged among neighbours every 30" via Response Message (advertisement)
- Implementation:
 - Application layer protocol, uses UDP/IP

Distance Vector Routing

- Best path info **shared** locally
 - Periodically
 - Upon any change
- Routing tables **updated** for
 - New entries
 - Cost changes
- Metric
 - Not necessarily hop count!

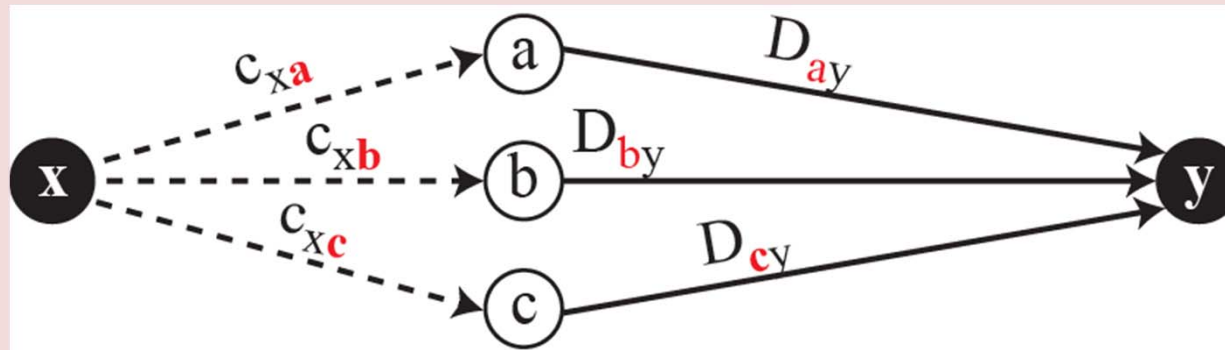
Nätgraf som träd



Distansvektor för A

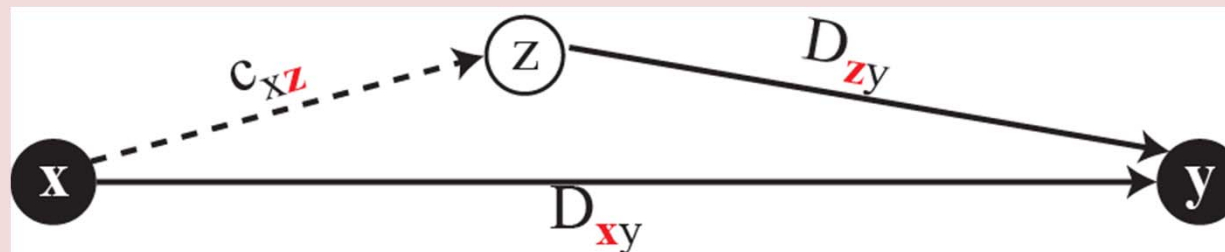
Nod	Dist
A	0
B	2
C	3
D	1
E	3

Bellman-Ford



a. General case with three intermediate nodes

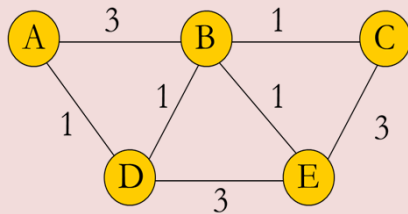
$$D_{xy} = \min\{(c_{xa} + D_{ay}), (c_{xb} + D_{by}), (c_{xc} + D_{cy}) \dots\}$$



b. Updating a path with a new route

$$D_{xy} = \min\{D_{xy}, (c_{xz} + D_{zy})\}$$

Uppdateringar



$$A[] = \min(A[], 3 + B[])$$

A, uppdaterad

Nod	Dist
A	0
B	3
C	4
D	1
E	4

A, ursprunglig

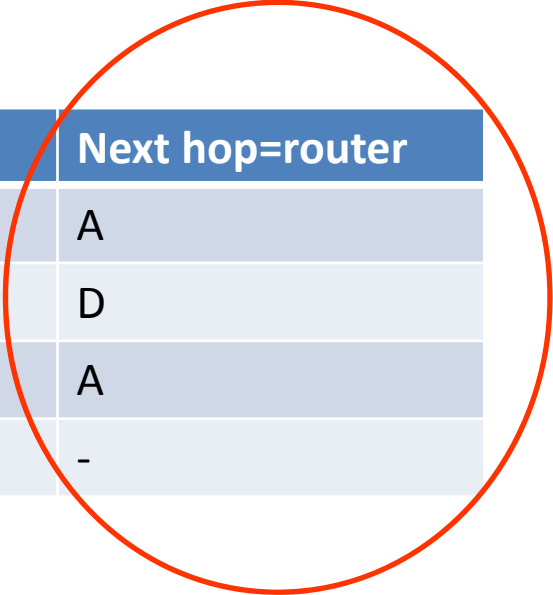
Nod	Dist
A	0
B	3
D	1

B

Nod	Dist
A	3
B	0
C	1
D	1
E	1

A RIP Forwarding/Routing Table

Destination=net	Cost	Next hop=router
123	3	A
32	5	D
16	3	A
7	2	-



RIP update message

- Contains the whole forwarding table
- Add 1 to cost in received message
- Change next hop to sending router
- Apply RIP updating algorithm

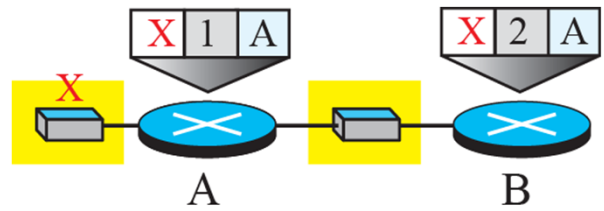
RIP Updating Algorithm (Bellman-Ford)

```
if (advertised destination not in table)
{
    add new entry // rule #1
}
else if (adv. next hop = next hop in table)
{
    update cost // rule #2
}
else if (adv. cost < cost in table)
{
    replace old entry // rule #3
}
```

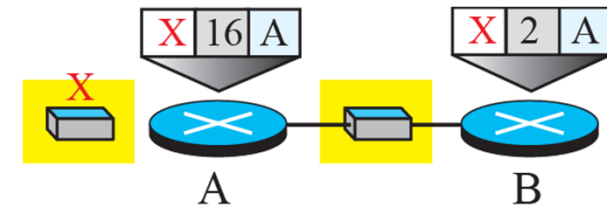
RIP Example



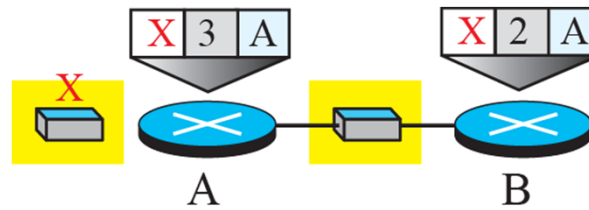
Two node instability/Count to infinity



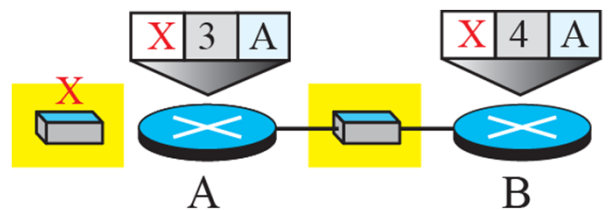
a. Before failure



b. After link failure

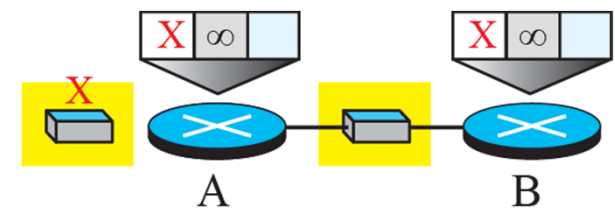


c. After A is updated by B



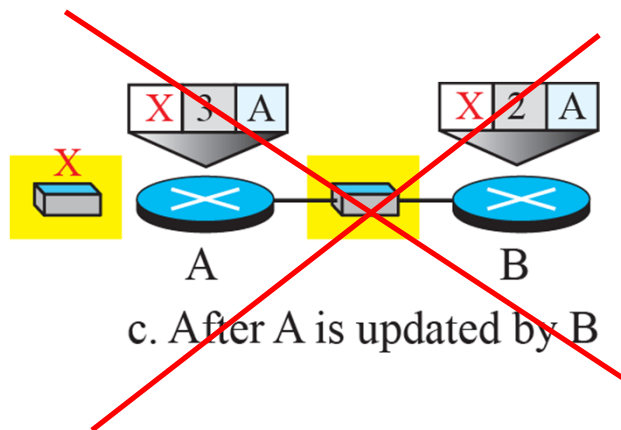
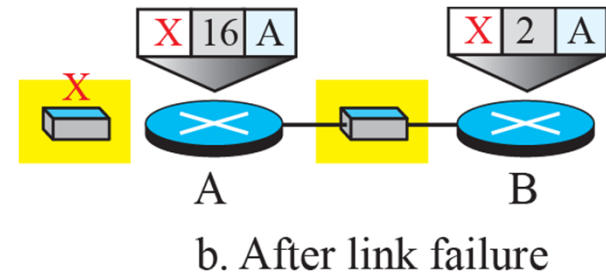
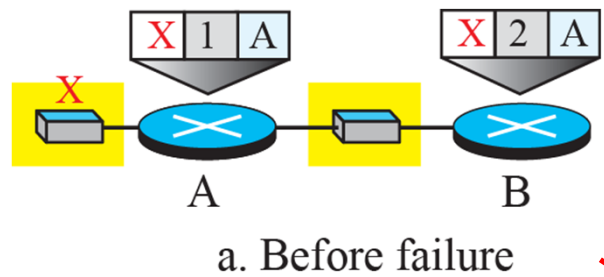
d. After B is updated by A

...



e. Finally

Split Horizon breaks Count to infinity

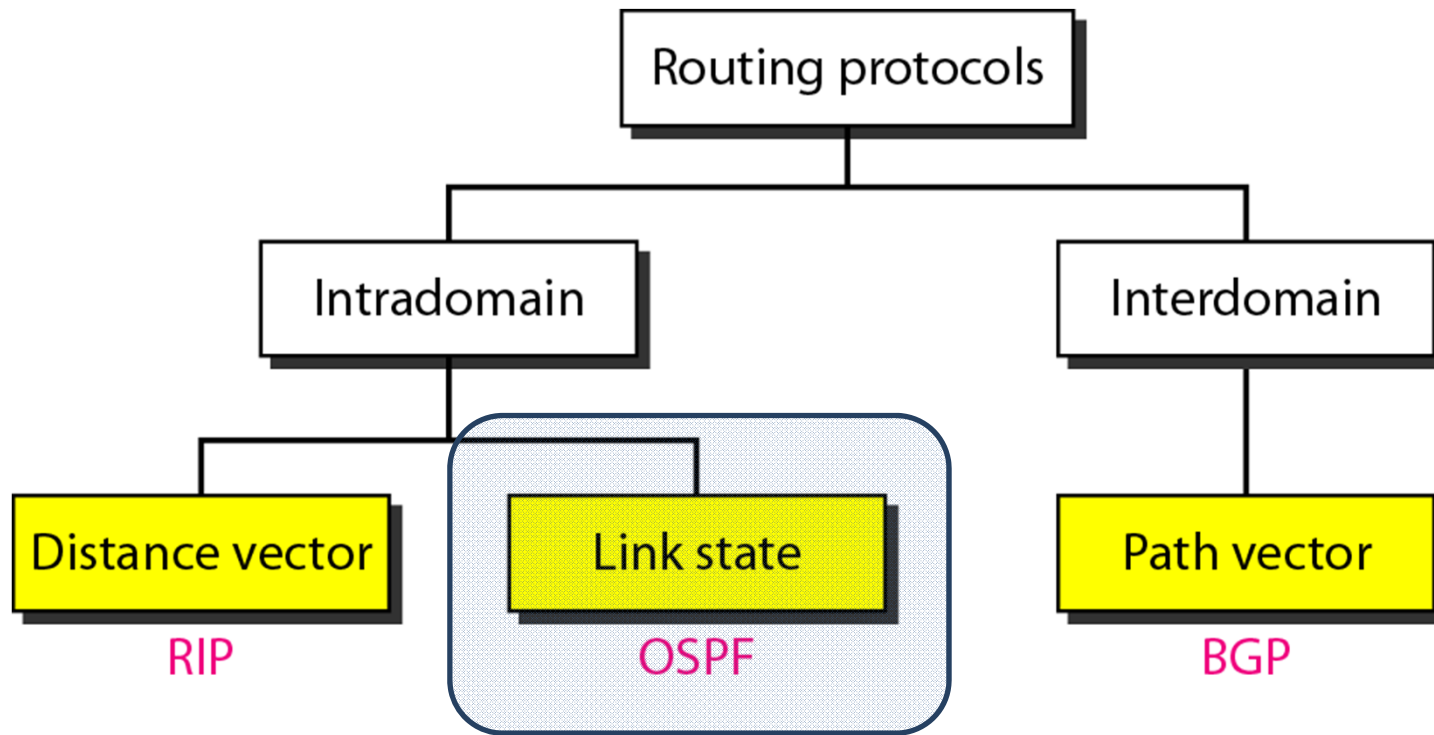


I have a route to X, but I got it from A so I won't tell A about it!

RIP: Link Failure and Recovery

- If no advertisement heard after 180”
 - Neighbour/link declared dead
 - Routes via neighbour invalidated (infinite distance = 16 hops)
 - New advertisements sent to neighbours (triggering a chain reaction if tables changed)
 - “Poison reverse” used to prevent count to infinity loops
 - “Good news travel fast, bad news travel slow”

Routing Algorithms and Protocols



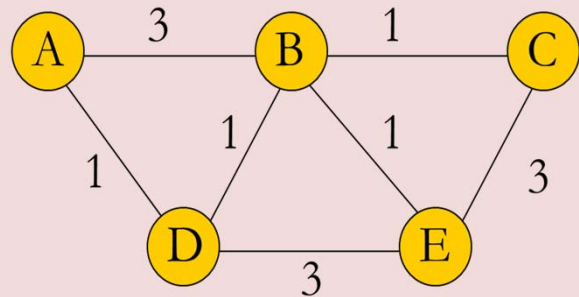
OSPF (Open Shortest Path First)

- Divides domain into areas
 - Limits flooding for efficiency
 - One "backbone" area connects all
- Distance metric:
 - Cost to destination network

Link State Routing

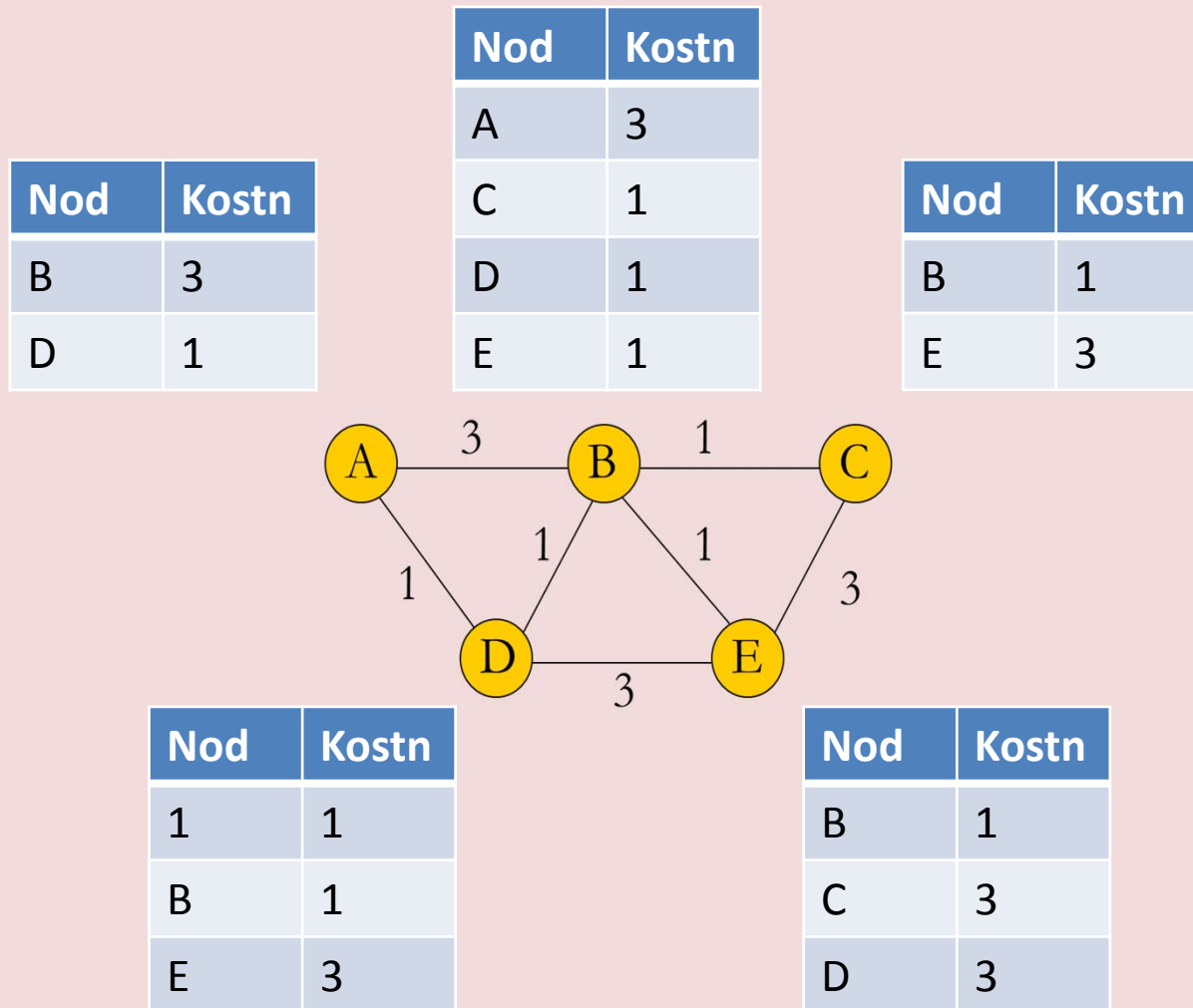
- Local topology info **flooded globally**
 - Periodically (very seldom ...)
 - Upon any change
- Routing tables **updated** for
 - Link state changes
 - Cost changes

Exempel på *link state* databas



	A	B	C	D	E
A	0	3	∞	1	∞
B	3	0	1	1	1
C	∞	1	0	∞	3
D	1	1	∞	0	3
E	∞	1	3	3	0

Link State Advertisements



Uppdateras vid förändring!

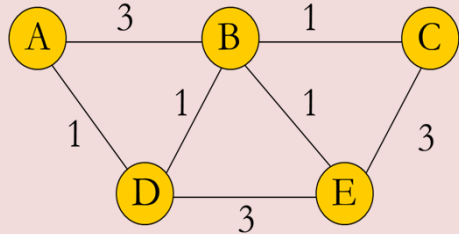
Tree Generation Algorithm (Dijkstra)

```
put yourself to tentative list
while tentative list not empty
{
  pick node which can be reached
    with least cumulative cost
  add it to your tree*
  put its neighbours to tentative list**
    with cumulative costs to reach them
}
```

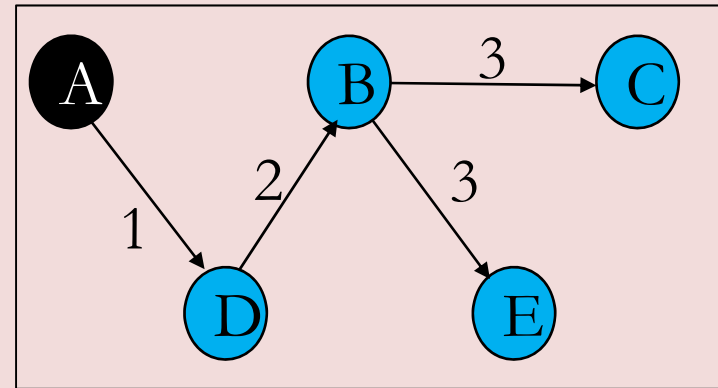
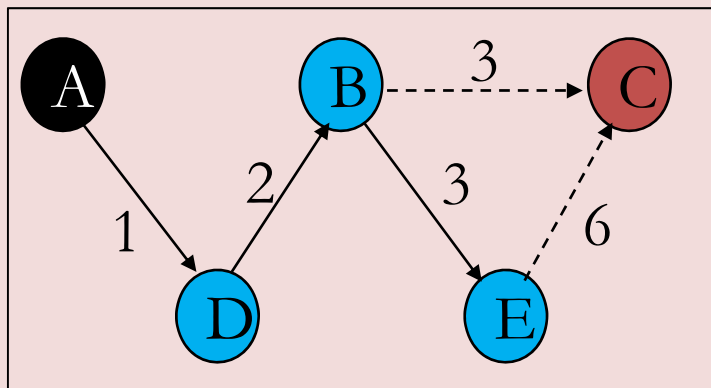
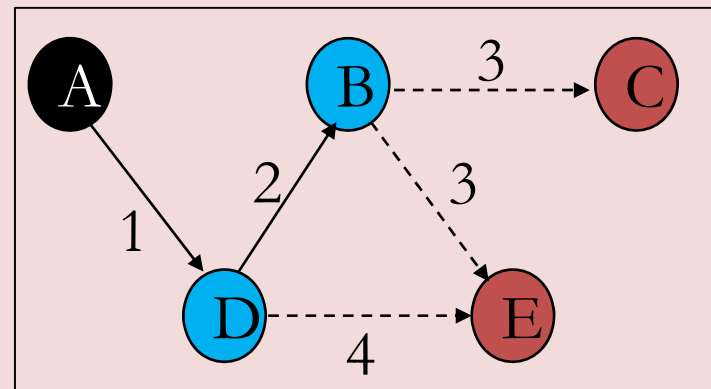
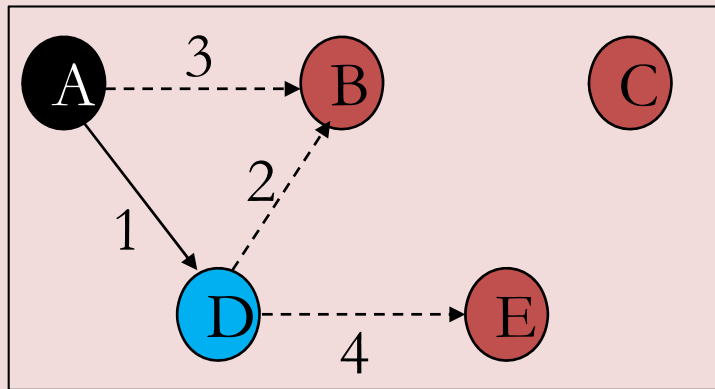
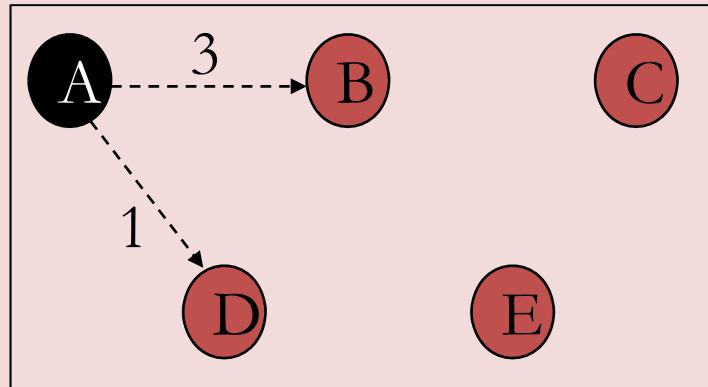
* (a. k. a. permanent list)

** (if not already there)

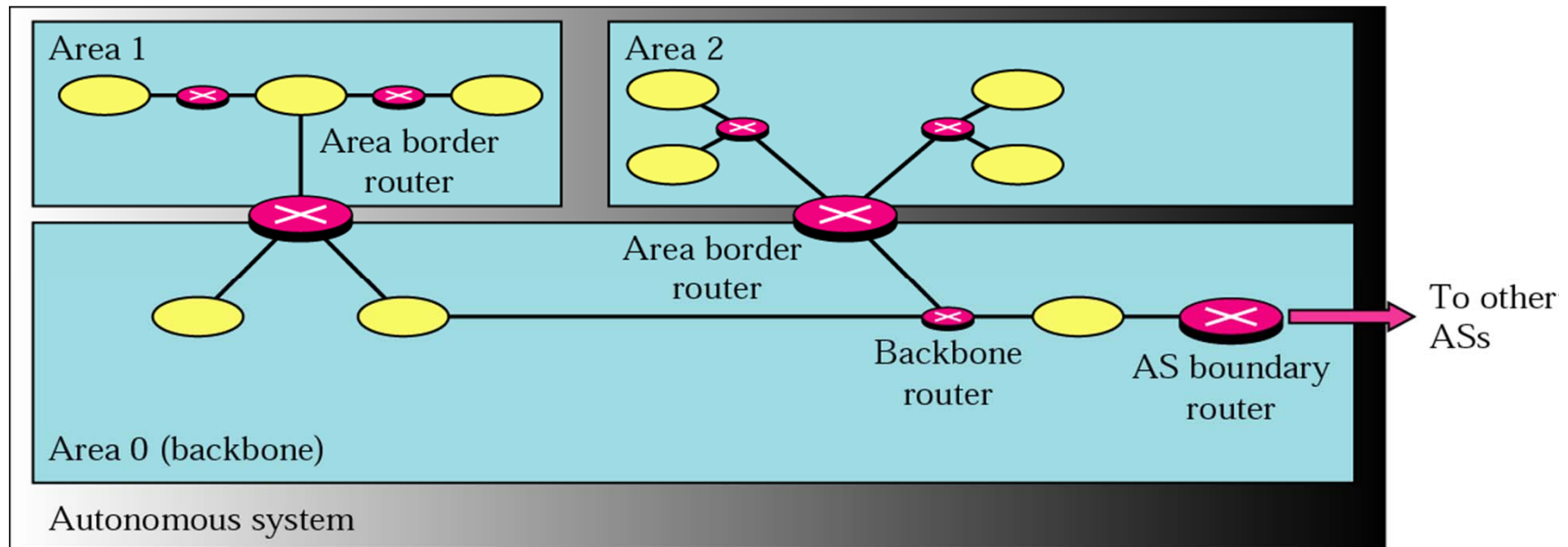
SPF exempel



- Rotnod
- Permanent nod
- Preliminär nod
- > Potentiell väg
- > Permanent väg

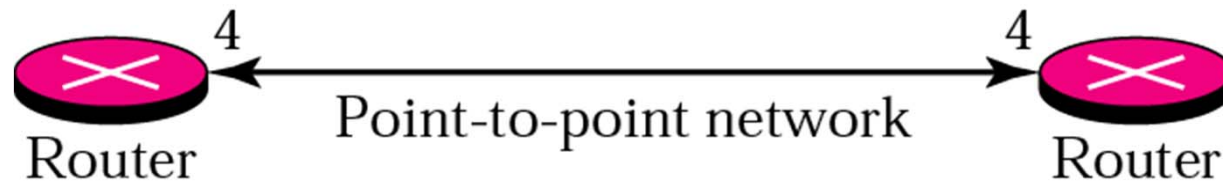
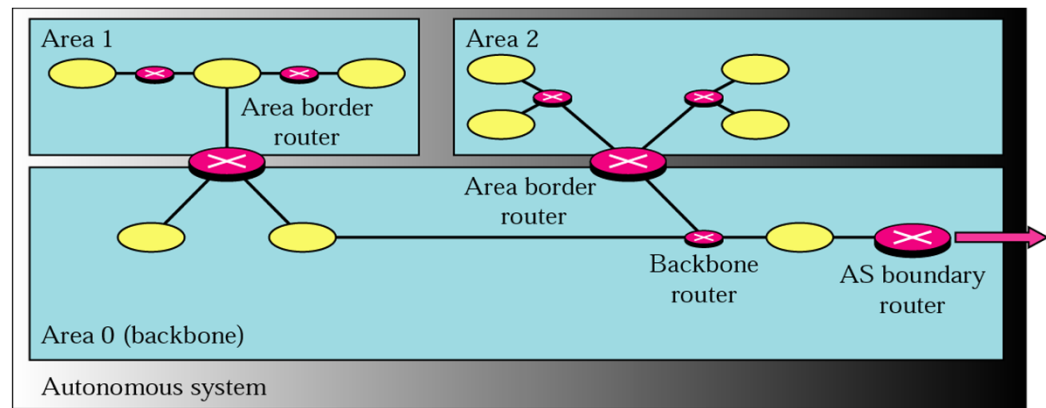


Areas, Router and Link Types

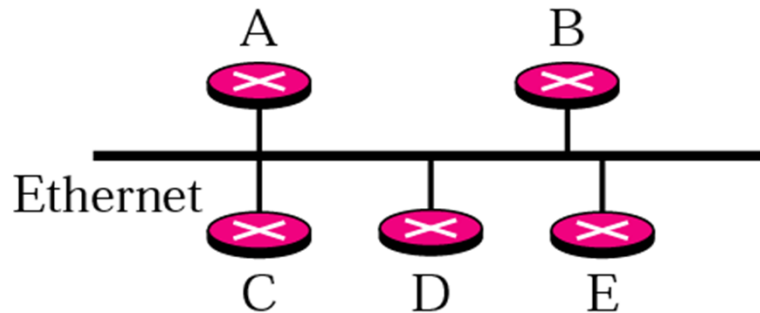
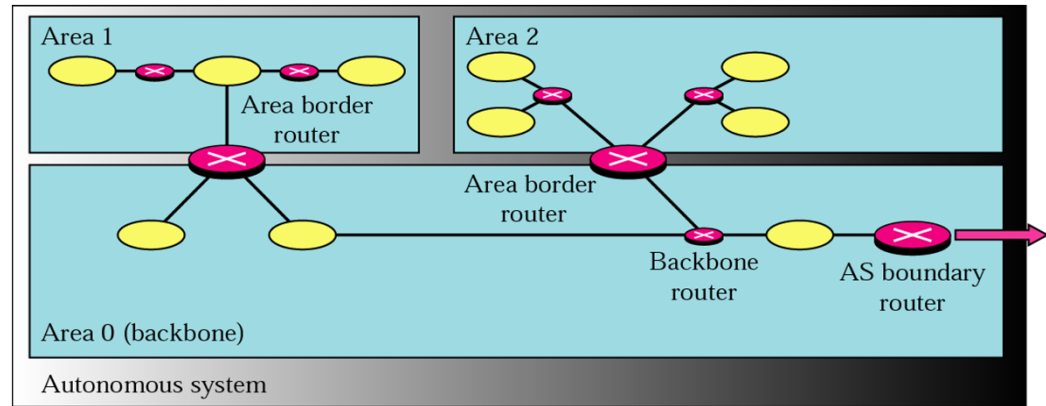


Point-to-Point Link

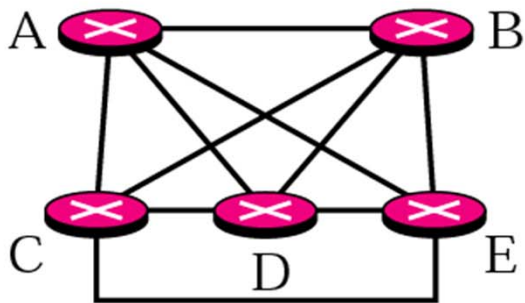
- Connects two routers
- No need for addresses



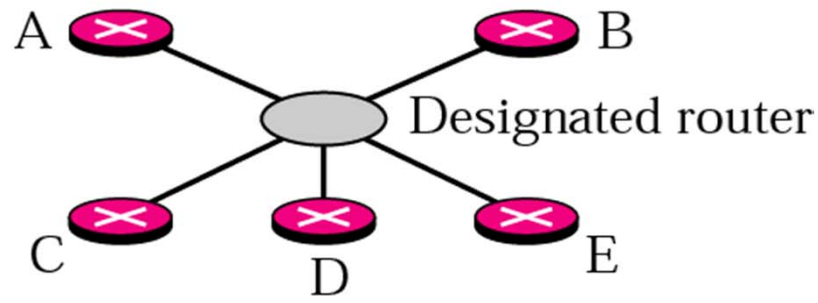
Transient Link



a. Transient network

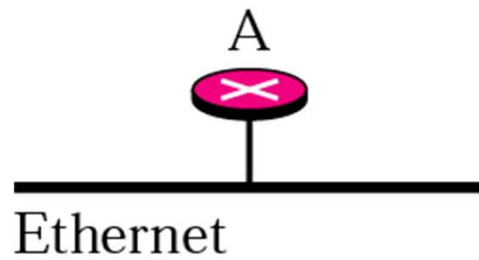
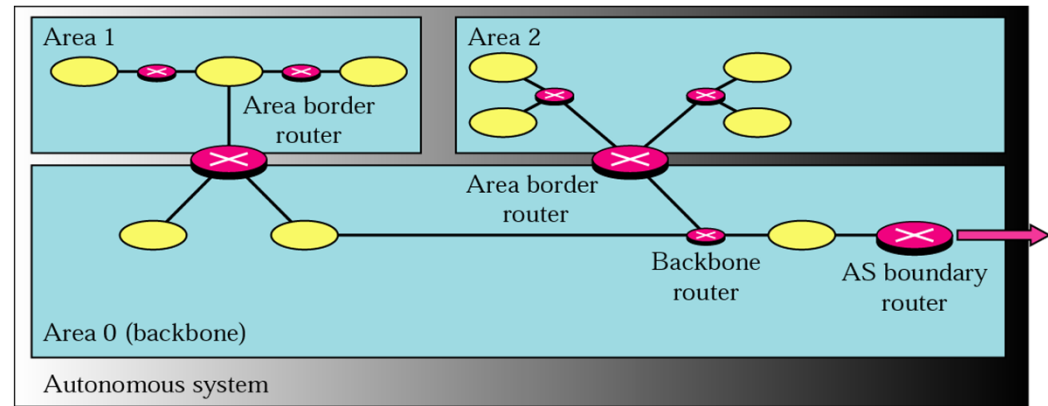


b. Unrealistic representation



c. Realistic representation

Stub Link



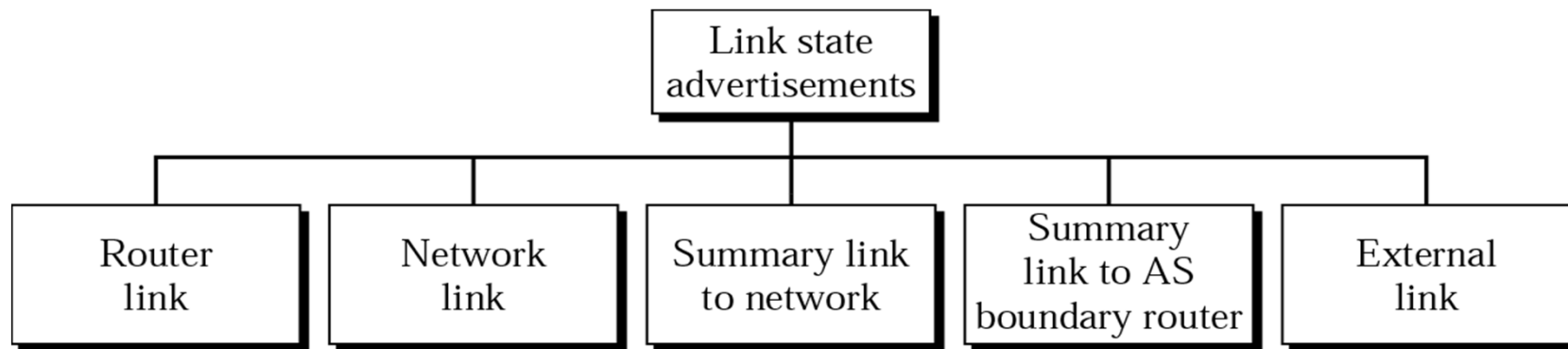
a. Stub network



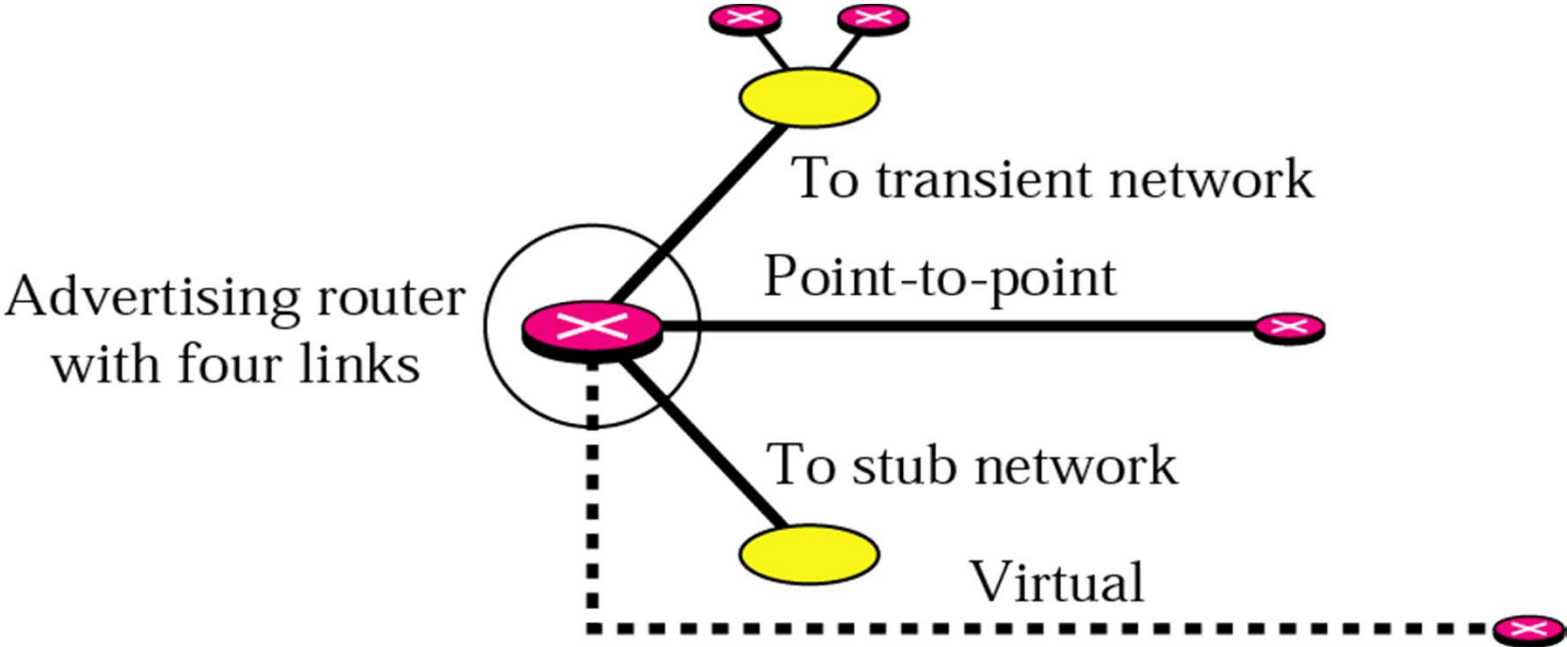
b. Representation

Link State Advertisements

- What to advertise?
 - Different entities as nodes
 - Different link types as connections
 - Different types of cost

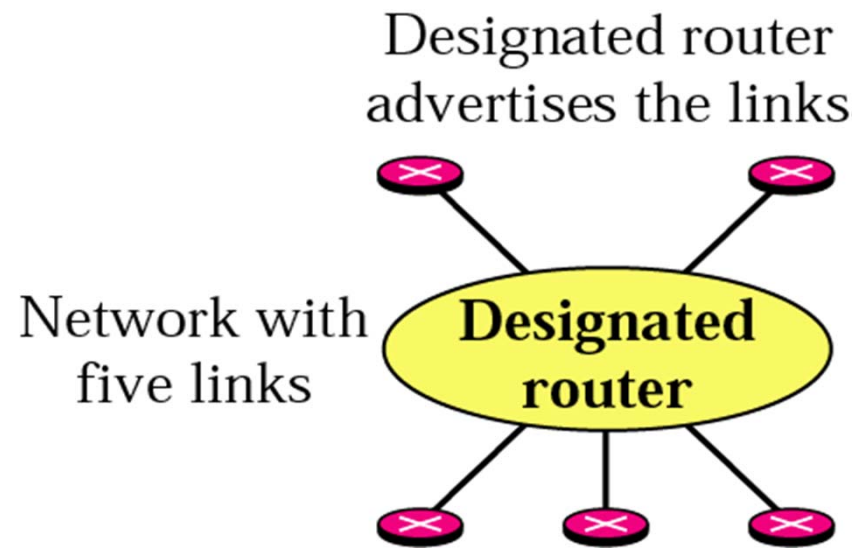


Router Link Advertisement



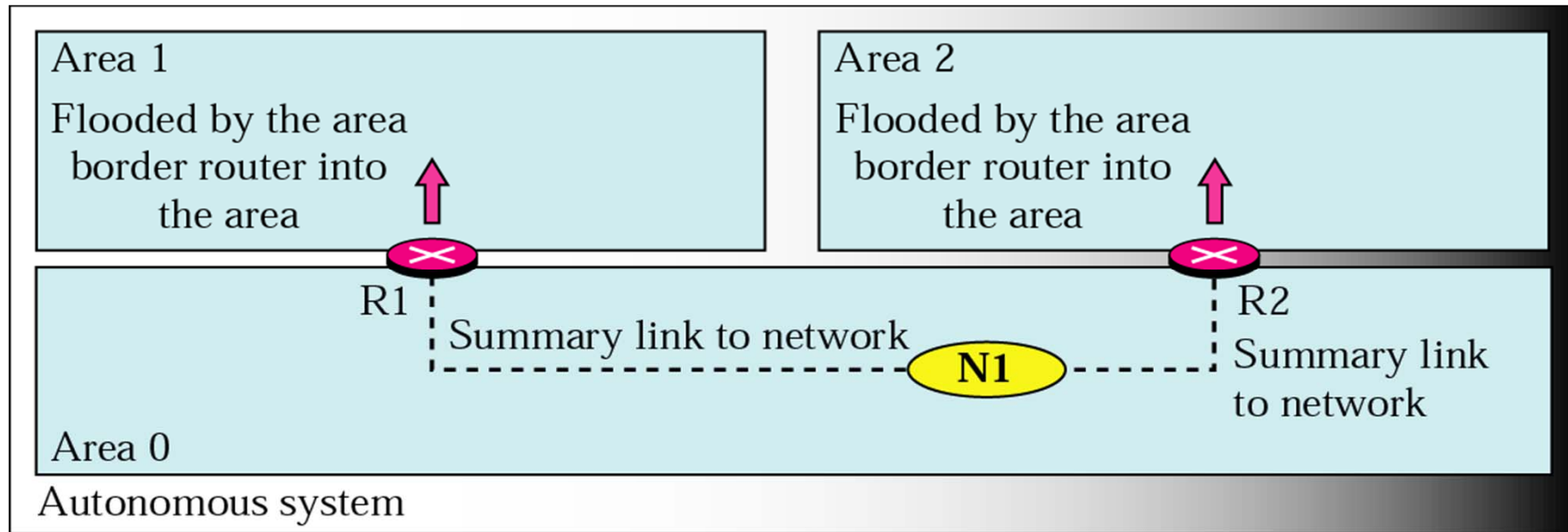
Network Link Advertisement

- Network is a passive entity
 - It cannot advertise itself



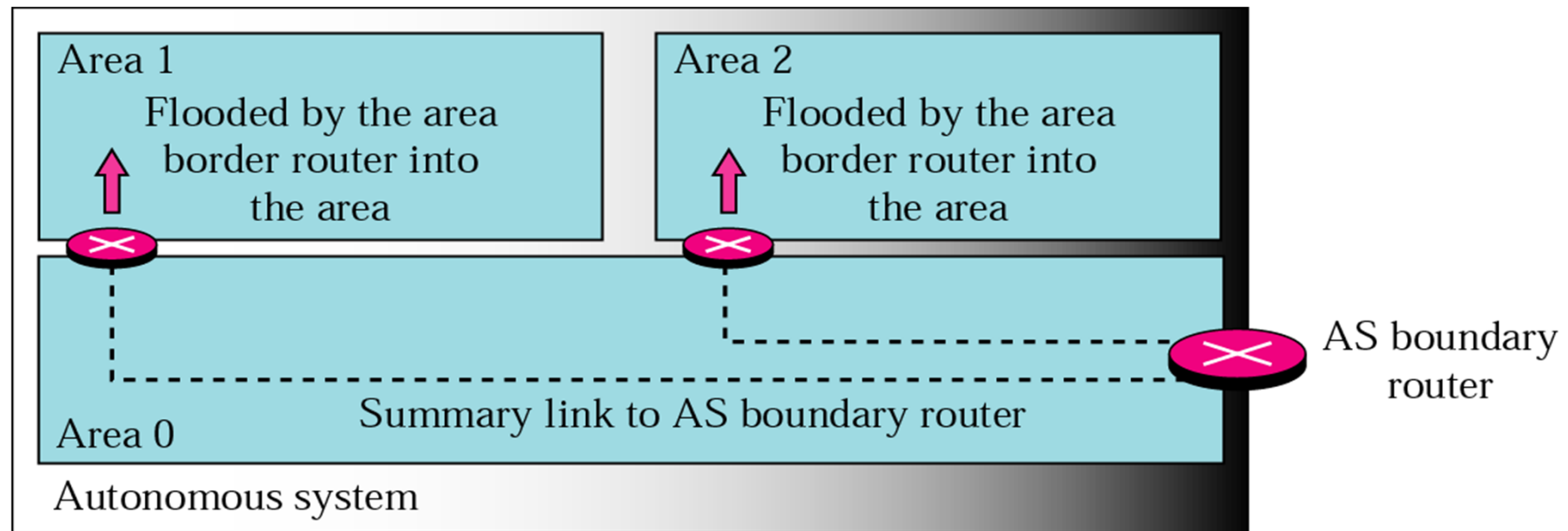
Summary Link to Network

- Done by area border routers
 - Goes through the backbone



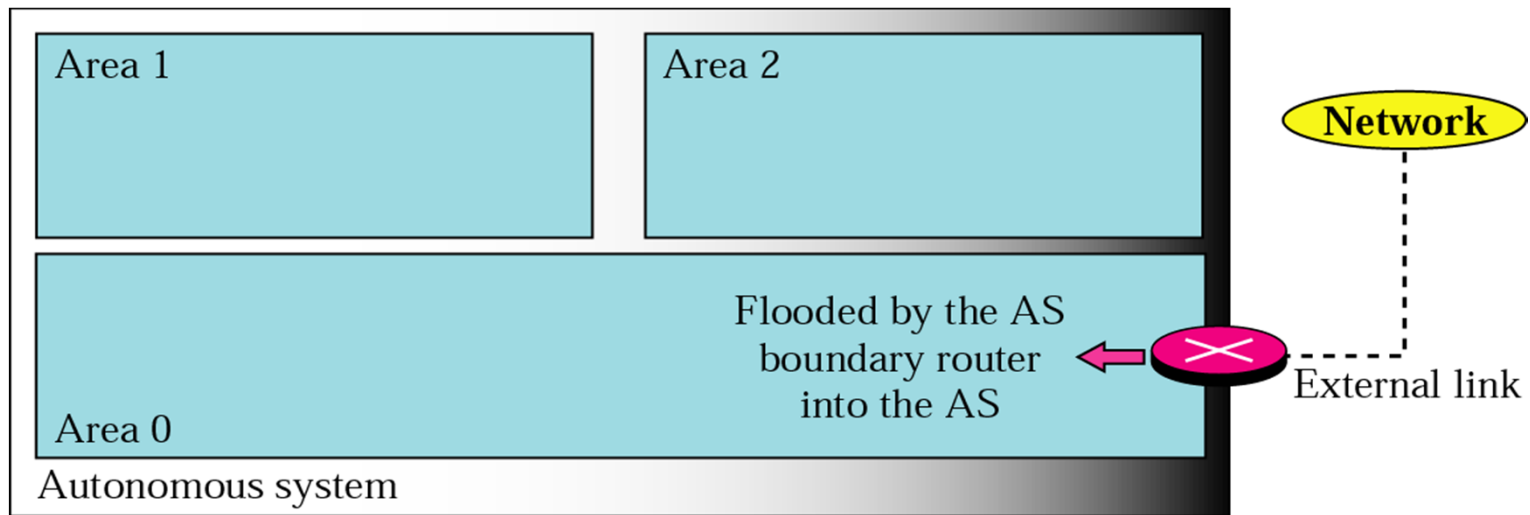
Summary Link to AS Boundary Router

- Links to other domains
"autonomous systems"



External Link Advertisement

- Link to a single network outside the domain



Hello message

- Find neighbours
- Keep contact with neighbours: I am still alive!
- Sent out periodically (typically every 10th second)
- If no hellos received during holdtime (typically 30 seconds), neighbour declared dead.
- Compare RIP update messages