



LUND
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Information Theory

Lecture 8

Capacity of DMC

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Discrete memoryless channel (DMC)

Definition

A **discrete channel** is a system $(\mathcal{X}, P(y|x), \mathcal{Y})$, with

- input alphabet \mathcal{X} of size N
- output alphabet \mathcal{Y} of size M
- transition probability distribution $P(y|x)$

The channel is **memoryless** if the probability distribution is independent of previous input symbols.



Channel capacity

Definition

The **information channel capacity** of a DMC is

$$C = \max_{p(x)} I(X; Y)$$

where the maximum is taken over all input distributions.

Theorem (Channel Coding Theorem)

A code rate is achievable if and only if

$$R < C = \max_{p(x)} I(X; Y)$$

DMC

Definition

A discrete memoryless channel is **uniformly dispersive** if the distribution for the outgoing branches has the same set of probabilities.

Theorem

If a channel is uniformly dispersive, then

$$H(Y|X) = H(\mathbf{r})$$

where \mathbf{r} is the set of probabilities for outgoing branches. The capacity of the channel can be written

$$C = \max_{p(x)} H(Y) - H(\mathbf{r})$$

DMC

Definition (Symmetric channels)

A discrete memoryless channel is **symmetric** if

- for all X , the branches leaving the code symbol have the same set of probabilities, p_1, p_2, \dots, p_M .
- for all Y , the branches entering the received symbol have the same set of probabilities, p_1, p_2, \dots, p_N .

Seen from the transition matrix $[p(y|x)]$ all rows are permutations of each other and all columns are permutations of each other.



DMC

Definition (Weakly symmetric channels)

A discrete memoryless channel is **weakly symmetric** if

- for all X , the branches leaving the code symbol have the same set of probabilities, p_1, p_2, \dots, p_M .
- for all Y , the branches entering the received symbol have the same sum, $\sum_{i=1}^N p_i$.

Seen from the transition matrix $[p(y|x)]$ all rows are permutations of each other and all columns have the same sum $\sum_x p(y|x)$.

DMC

Theorem

If a channel is symmetric or weakly symmetric, the channel capacity is

$$C = \log M - H(\mathbf{r})$$

where \mathbf{r} is the set of probabilities labeling branches leaving a code symbol X , or, viewed in a transition matrix, one row of the transition matrix.