

Exercises:

Part I

1. Represent 305 with 10 bits
2. Represent -305 with 10 bits
3. Represent -197 with 9 bits
4. Represent 197 with 9 bits
5. Represent "0101010101" in hexadecimal
6. Represent "1100111101" in hexadecimal
7. Represent "11001100" as decimal (unsigned representation)
8. Represent "11001100" as decimal (signed representation)

Part II

1. Assume *b* is a variable of size 4 bytes and is stored in a byte addressable memory at address 0xA80. If the processor's endianness is little-endian, and the processor writes the value 0xA155F0D3 in the variable *b* which bytes will be written to each memory address.
2. Assume *x* is a variable of type pointer that points to a single byte. Further, assume *b* is a variable of size 4 bytes and is stored at memory address 0xA80. Given that the processor uses big-endian, evaluate the new value of *b* after the following code is executed:

```
b=0x2F552;  
x=0xA81;  
b=b*x;
```

3. Given a variable *b* which is assumed to have a value in the range [0..7], write the necessary statements in C to ensure that the bit at bit position *b* in another variable *c* is set to one.
4. Write a statement in C such that for a given variable *b* the bit at position 3 is set to 0, the bit at position 5 is set to 1, the bit at bit position 2 is inverted. Assume that the variable *b* is of size 1 byte.

Answers:

Part I

1.	$305/2 = 152$	1	↑	
	$152/2 = 76$	0		
	$76/2 = 38$	0		
	$38/2 = 19$	0		
	$19/2 = 9$	1		
	$9/2 = 4$	1		
	$4/2 = 2$	0		
	$2/2 = 1$	0		
	$1/2 = 0$	1		
	$0/2 = 0$	0		

(0100110001)

2. $\sim 0100110001 = 1011001110$

	$+ \quad \quad \quad 1$	
	$= 1011001111$	(-305)

3. $-197 = (-1) \cdot 2^8 + X = -256 + X \rightarrow (1 \text{ --- } -) \quad X=59 \text{ with 8 bits}$

59/2 = 29	1	↑	
29/2 = 14	1		
14/2 = 7	0		
7/2 = 3	1		
3/2 = 1	1		
1/2 = 0	1		
0/2 = 0	0		
0/2 = 0	0		

$\rightarrow (-197) = (100111011)$

4. $\sim 100111011 = 011000100$

$$\begin{array}{r} + \quad \quad \quad 1 \\ \hline = 011000101 \end{array} \quad (197)$$

5. 0101010101- unsigned 0001|0101|0101 → 0x155

0101010101- signed 0001|0011|0011 → 0x155

6. 1100111101- unsigned 0011|0011|1101 → 0x33D

1100111101- signed 1111|0011|1101 → 0xF3D

7. $11001100 \rightarrow 1 \cdot 2^7 + 1 \cdot 2^6 + 0 \cdot 2^5 + 0 \cdot 2^4 + 1 \cdot 2^3 + 1 \cdot 2^2 + 0 \cdot 2^1 + 0 \cdot 2^0 = 204$

8. $11001100 \rightarrow (-1) \cdot 2^7 + 1 \cdot 2^6 + 0 \cdot 2^5 + 0 \cdot 2^4 + 1 \cdot 2^3 + 1 \cdot 2^2 + 0 \cdot 2^1 + 0 \cdot 2^0 = -52$

Part II

1.

Address	Value
0xA80	0xD3
0xA81	0xF0
0xA82	0x55
0xA83	0xA1

2. First, the value 0x2F552 needs to be extended to 32 bits, i.e. $b = 0x0002F552$. This variable will be stored in memory as

Address	Value
0xA80	0x00
0xA81	0x02
0xA82	0xF5
0xA83	0x52

As x points to memory address 0xA81, the expression $*x$ is evaluated as 0x02 (see table above).

The new value of b is then

$$\begin{array}{r} 0x0002F552 \\ + \quad \underline{0x00000002} \\ \hline 0x0002F554 \end{array}$$

3. $c = (1 \ll b) | c;$

4. $b = ((b \& 0b11110111) | 0b00100000) \wedge 0b00000100;$