EITF35: Introduction to Structured VLSI Design

Part 2.2.1: Sequential circuit

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Lab computer log in

Username: group name (svd01-svd16)

Password: 15Change

You may change your password after first log-in

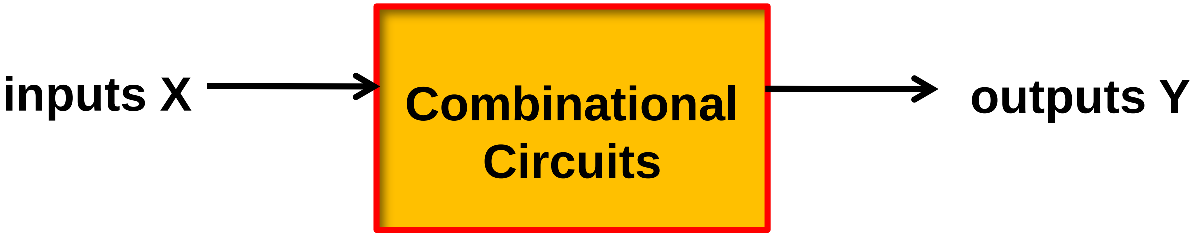
Come to me if the door access is not working

Outline

- □ Sequential vs. Combinational
- □ Synchronous vs. Asynchronous
- □ Basic Storage Elements
- □ Timing
- □ Folding & Pipeline

Sequential vs. Combinational

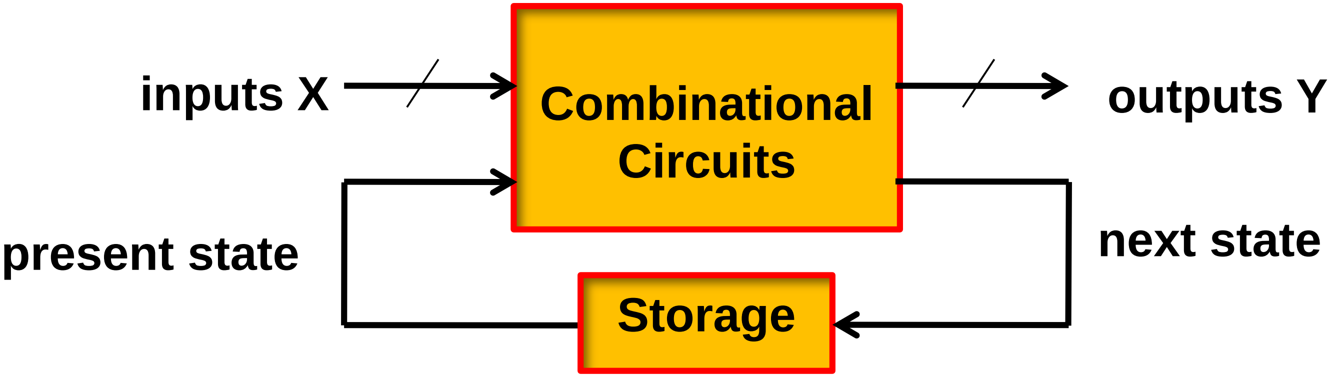
- A combinational circuit:



- At any time, outputs depend only on present inputs
  - Changing inputs changes outputs
- No regard for previous inputs
  - No memory (history)
- Time is “ignored” !
  - Time-independent circuit

Sequential vs. Combinational

- A sequential circuit:

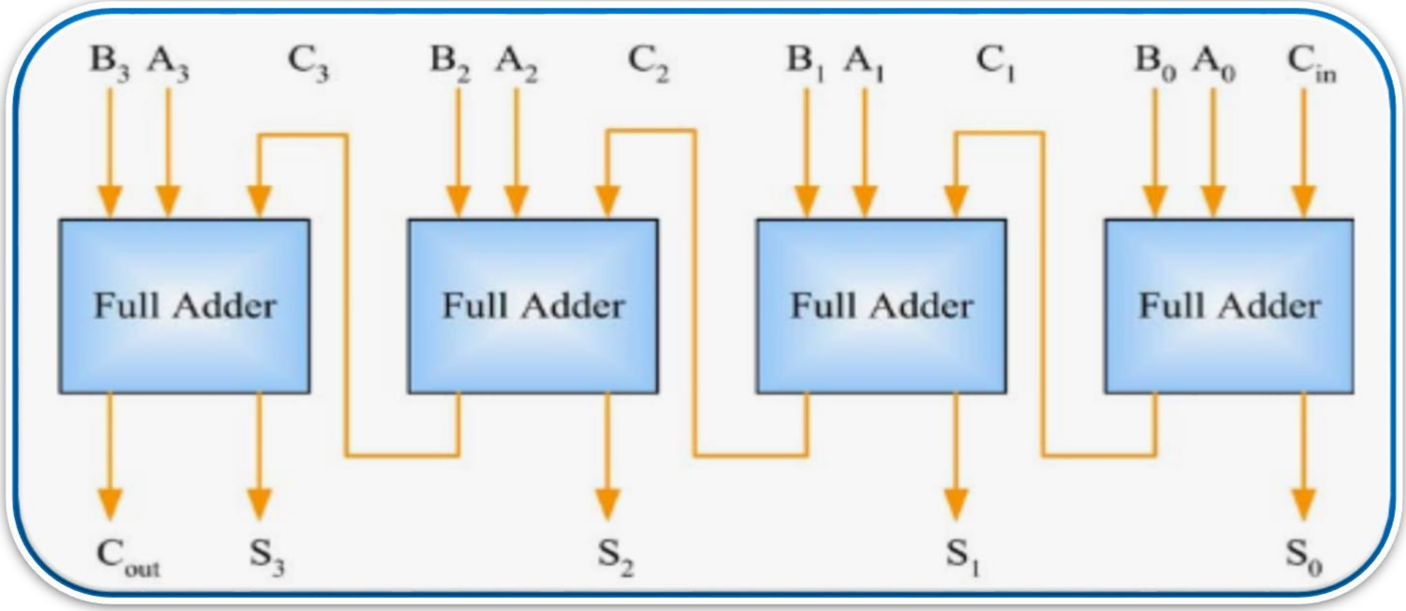


- **Outputs depends on inputs and past history of inputs**
  - Previous inputs can be stored into **storage elements**
  - **Input order matters**

Sequential vs. Combinational: adders

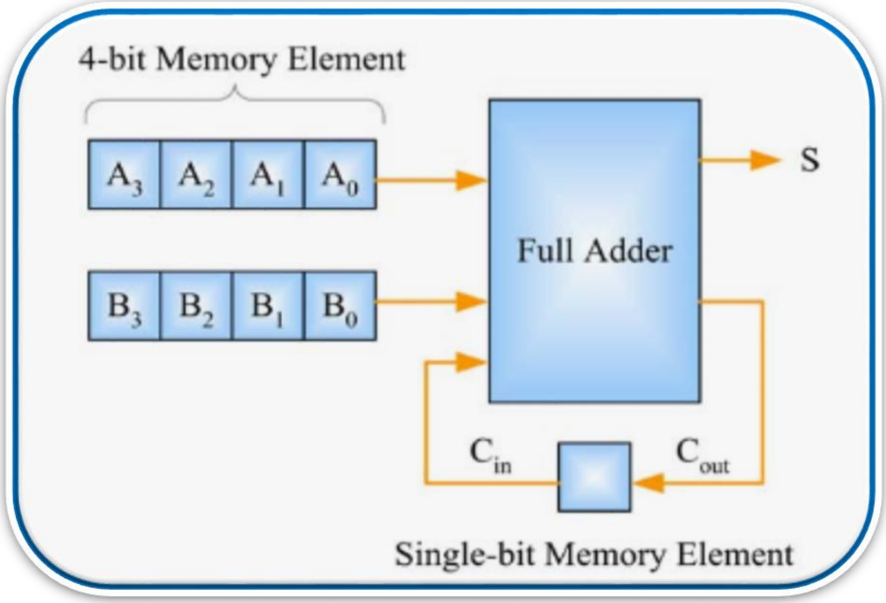
- Calculate $A_3A_2A_1A_0 + B_3B_2B_1B_0$

- Combinational adder
  - 4 full adders are required
  - One adder is active at a time slot



What we can do with storage elements?

□Sequential Adder



□Folding!

- One full adder
- 1-bit memory for carry
- Two 4-bit memory for operators

□4 clock cycles to get the output

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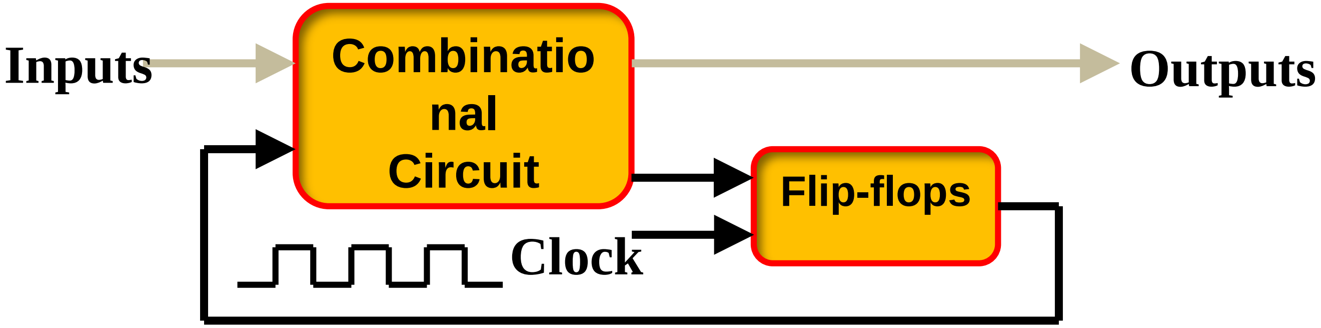
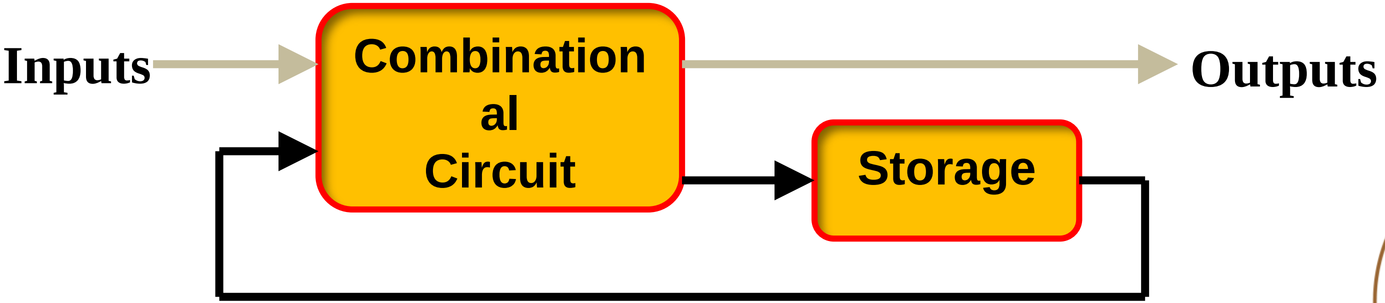
Synchronous vs. Asynchronous

□Two types of sequential circuits:

- Synchronous: The behavior of the circuit depends on the input signal at discrete instances of time (also called **clocked**)



- Asynchronous: The behavior of the circuit depends on the input signals at ***any instance of time***



Synchronous vs. Asynchronous

- When you have a clock
- You know that washer takes 1 hour
- You put the laundry in the washer and leave
- Dry 1hour later

Synchronous vs. Asynchronous

- What if you don’t have a clock …

Synchronous or Asynchronous?

□ **Sync. Advantages: Simplicity to design, debug, and test**

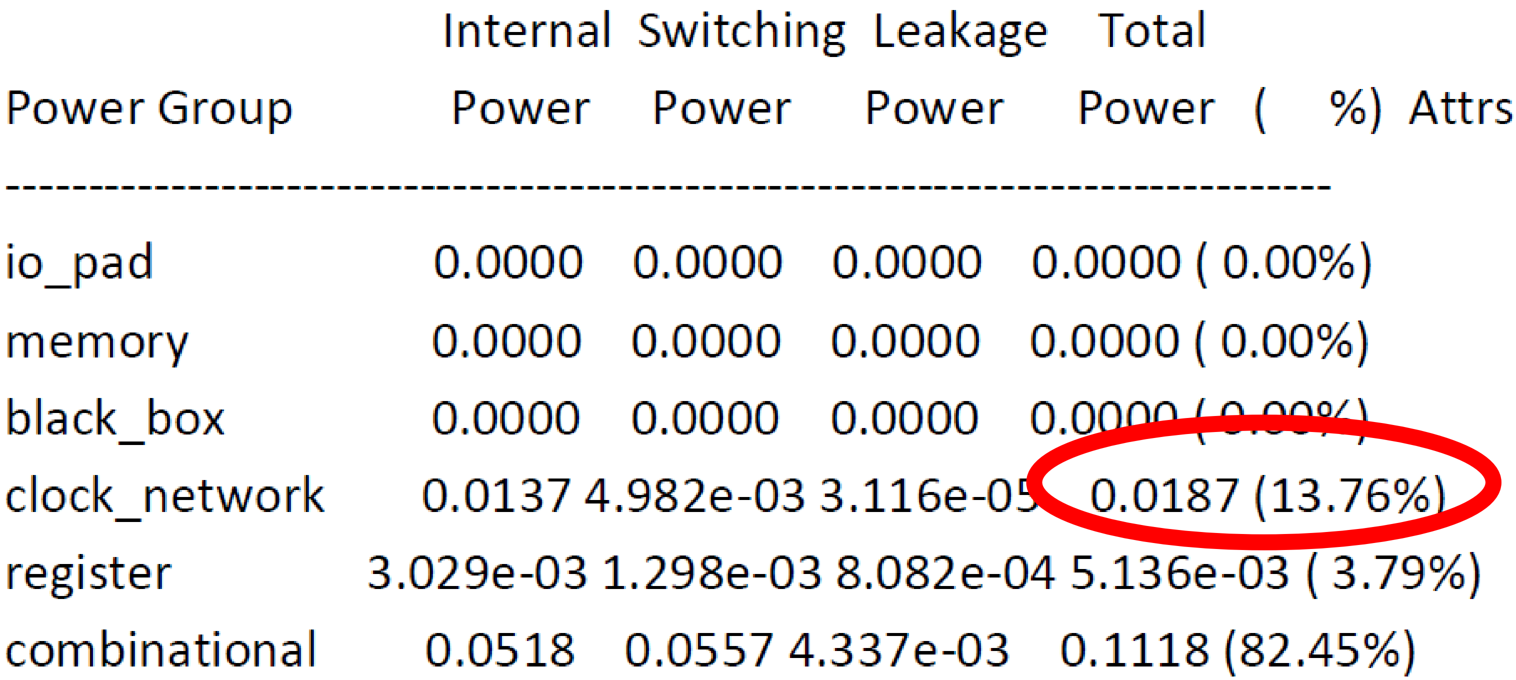
- Timing is controlled by one simple clock
- No hand-shake circuits
- Well supported by EDA tools
- Recommended for **VLSI**

□ **Sync. Disadvantages:**

- Performance constrained by worst-case: critical path
- Overhead for clock network
- Less power efficient

We will focus on synchronous circuits in this course

Power Example



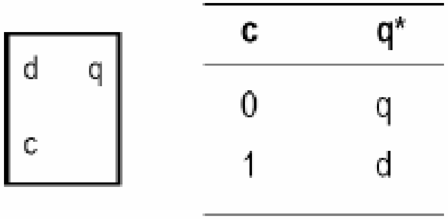
```
                     Internal  Switching  Leakage   Total
Power Group           Power     Power     Power     Power  (   %)  Attrs
--------------------------------------------------------------------
io_pad               0.0000    0.0000    0.0000   0.0000 ( 0.00%)
memory               0.0000    0.0000    0.0000   0.0000 ( 0.00%)
black_box            0.0000    0.0000    0.0000   0.0000 ( 0.00%)
clock_network       0.0137 4.982e-03 3.116e-05    0.0187 (13.76%)
register         3.029e-03 1.298e-03 8.082e-04 5.136e-03 ( 3.79%)
combinational       0.0518    0.0557 4.337e-03    0.1118 (82.45%)
```

Outline

- □Sequential vs. Combinational
- □Synchronous vs. Asynchronous
- □**Basic Storage Elements**
- □Timing
- □Folding & Pipeline

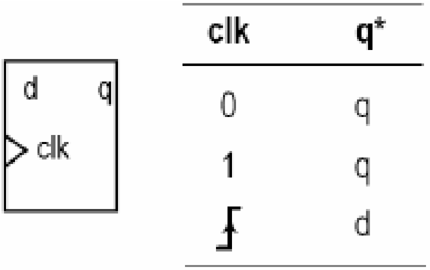
Basic storage element

- □D latch: level sensitive
- □D flip-flop (D-FF): edge sensitive



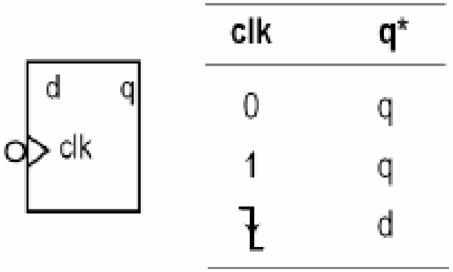
c	q*
0	q
1	d

D latch



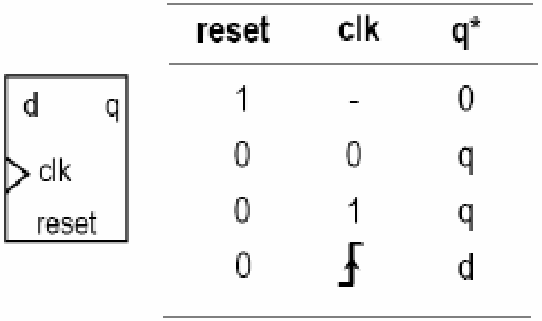
clk	q*
0	q
1	q
↑	d

pos-edge triggered D-FF



clk	q*
0	q
1	q
↓	d

neg-edge triggered D-FF

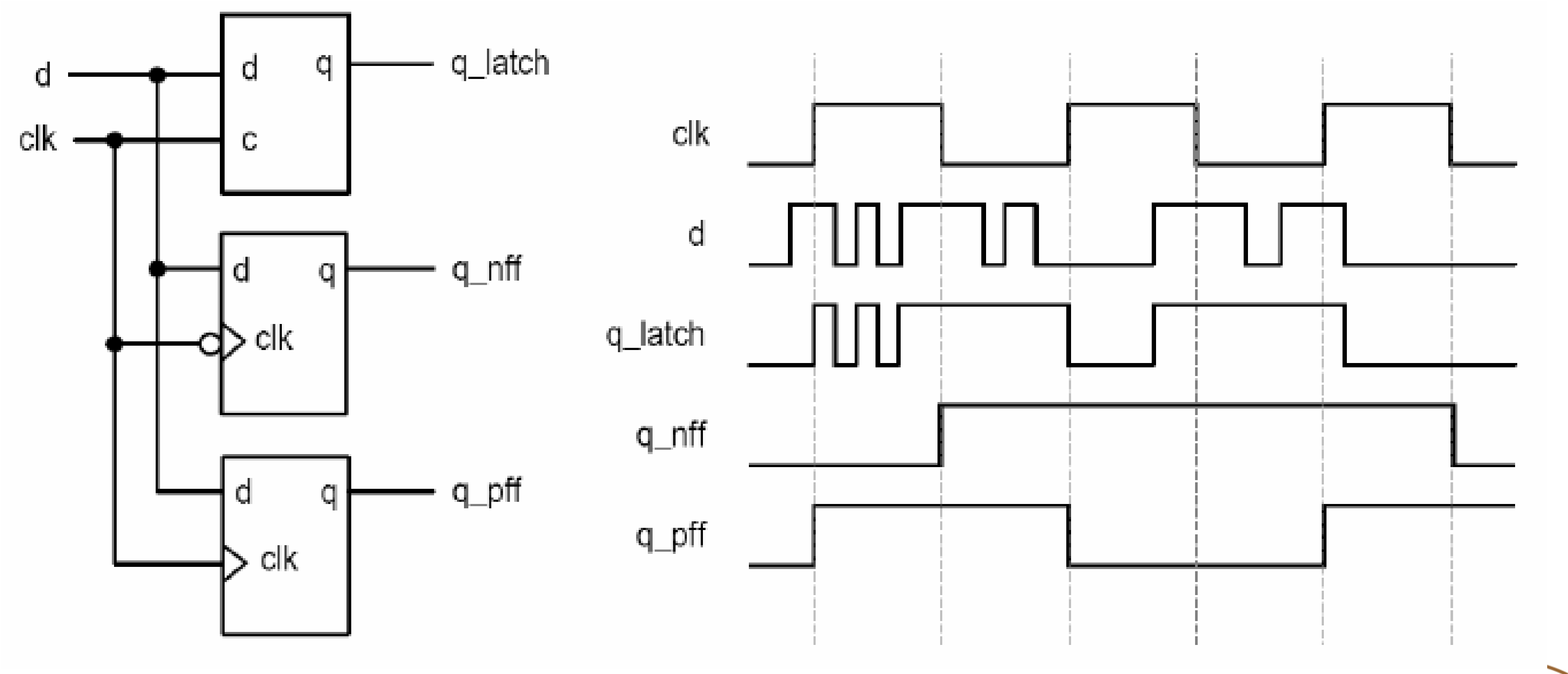


reset	clk	q*
1	-	0
0	0	q
0	1	q
0	↑	d

D-FF with reset

Basic storage element

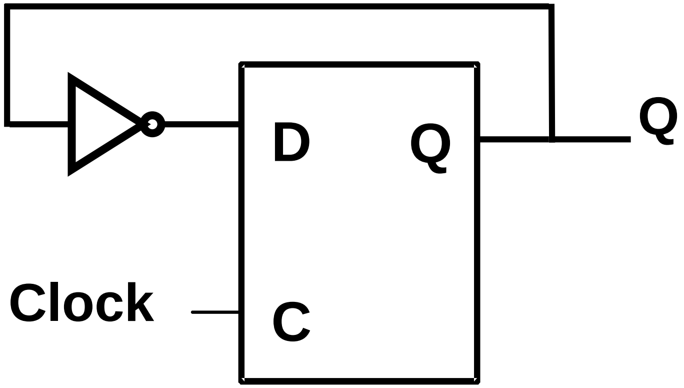
- D latch: level sensitive
- D flip-flop (D-FF): edge sensitive



Problem with Latches

- **Problem: A latch is transparent; state keep changing as long as the clock remains active**
- **Due to this uncertainty, latches can not be reliably used as storage elements.**
- **What happens if Clock=1? What will be the value of Q when Clock goes to 0?**

Latch Example



Most EDA software tools have difficulty with latches.

Outline

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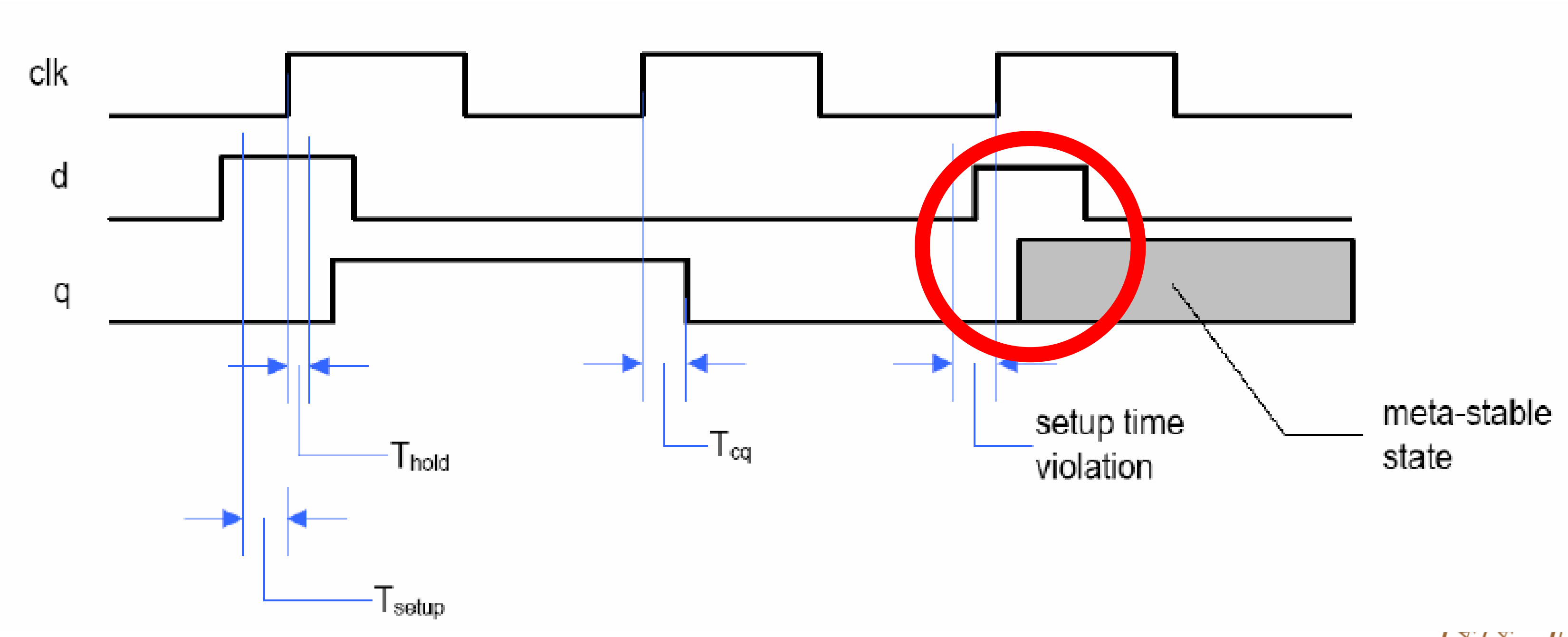
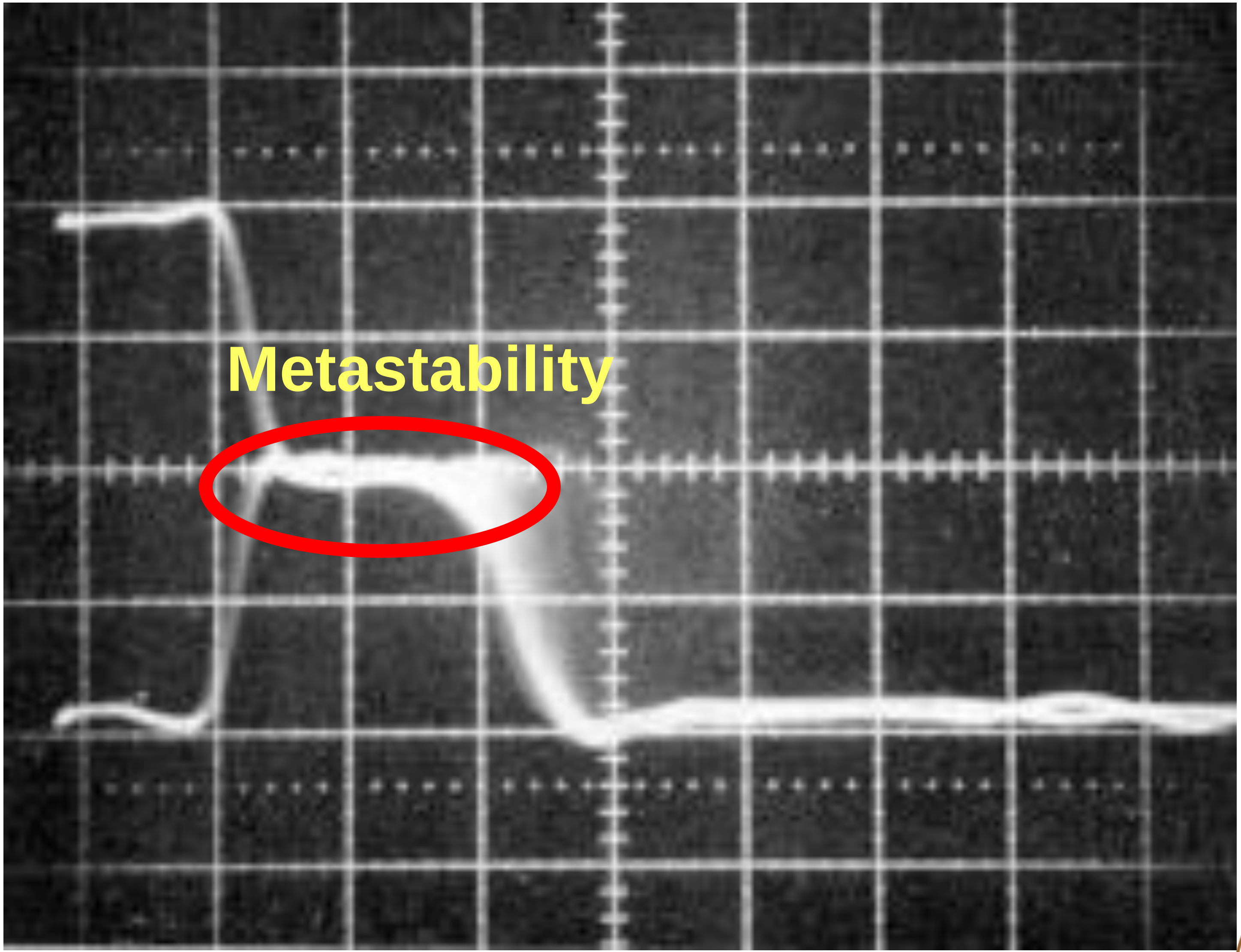
Flip Flops Timing

Very Important Timing Considerations!

- ***Setup Time*** **(Ts): The minimum time during which D input must be maintained before the clock transition occurs.**
- ***Hold Time*** **(Th): The minimum time during which D input must not be changed after the clock transition occurs.**

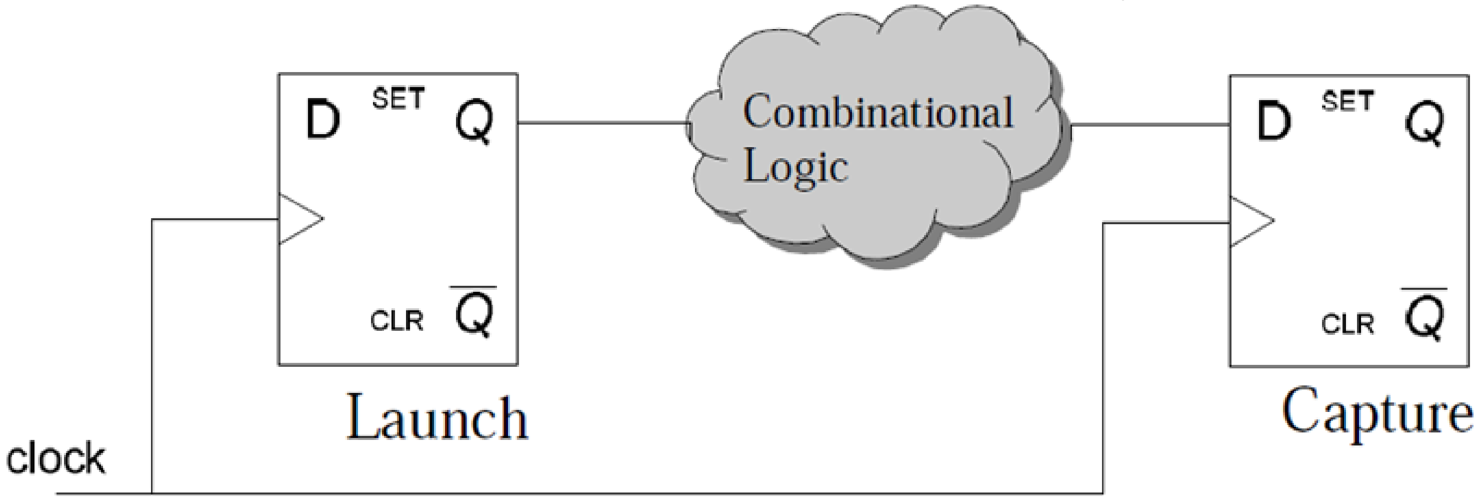


Metastability in Digital Logic



How fast can a synchronous circuit run?

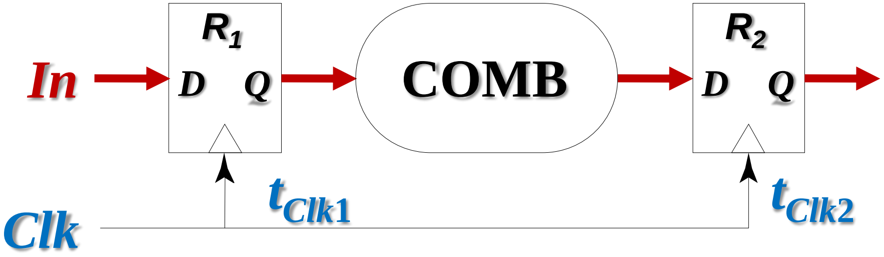
- **RTL (Register Transfer Level)**



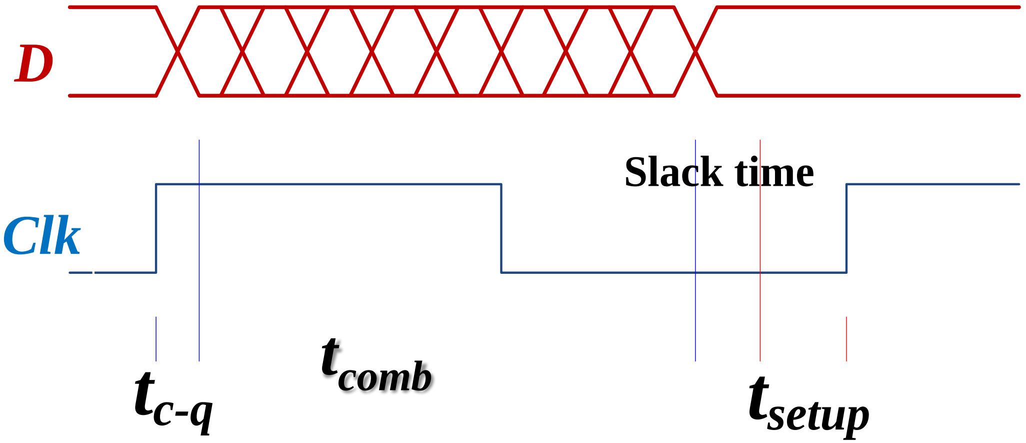
- **Timing analysis:**
  - Starting with the clock rising edge at the launch FF, end with the clock rising edge (next period or same period) of the capture FF

Setup Time

- **Setup Timing analysis:**
  - Starting with the clock rising edge at the launch FF, end with the clock rising edge (next period) of the capture FF

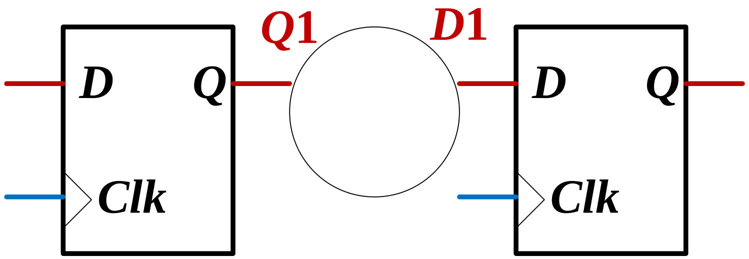


- **Data-Path (arrive time):** $T_{\text{Combinational logic}}$ + FF_{launch}(clk -> Q)
- **Clock-Path (required time):** Clock Period - FF $_{\text{tSetup}}$
- **Timing constraint :** $T_{\text{Combinational logic}}$ + FF_{launch}(clk -> Q) < Clock Period - FF $_{\text{tSetup}}$

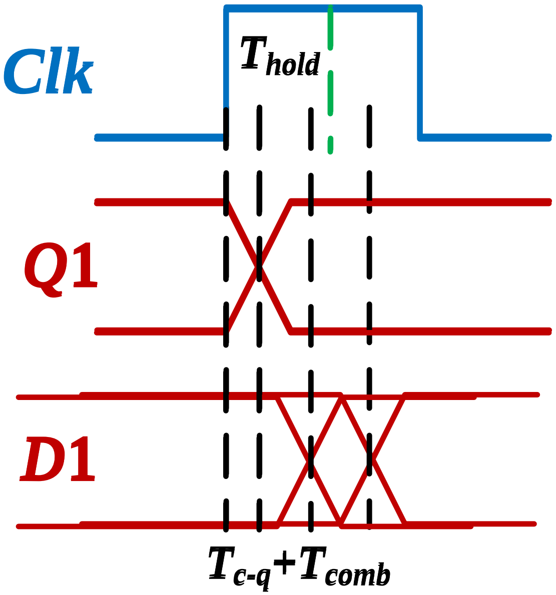


Hold Time

- **Hold Timing analysis:**
  - Starting with the clock rising edge at the launch FF, end with the clock rising edge (same period) of the capture FF

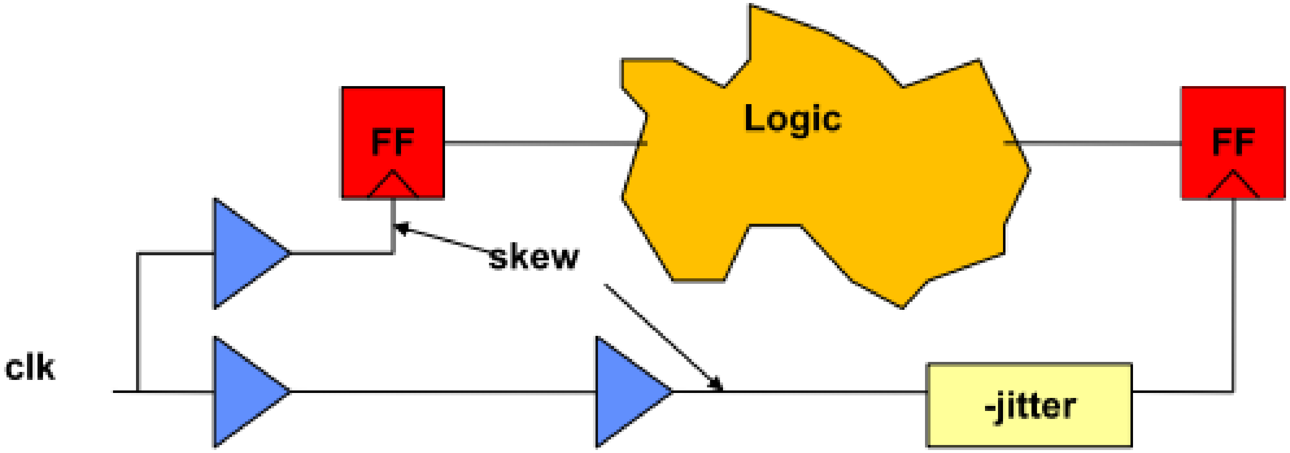


- **Data-Path (arrive time):** $T_{\text{Combinational logic}}$ + FF_{launch}(clk -> Q)
- **Clock-Path (required time):** FF_{tHold}
- **Timing constraint :** $T_{\text{Combinational logic}}$ + FF_{launch}(clk -> Q)> FF_{tHold}



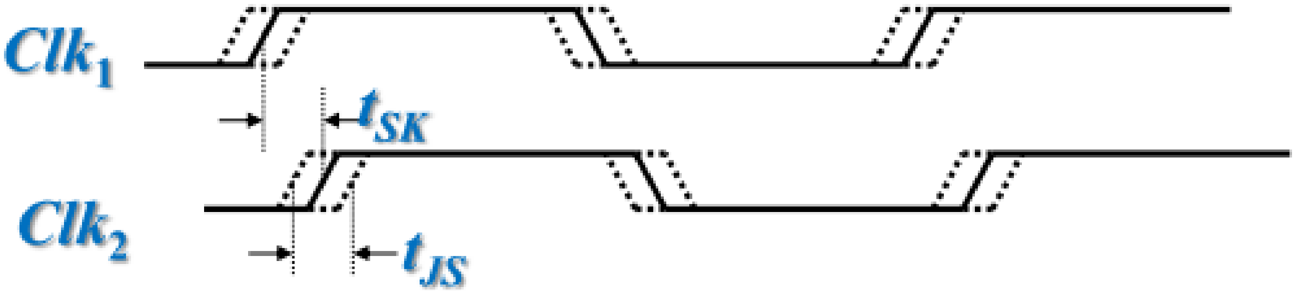
Clock uncertainty

- **Clock uncertainty = skew ±jitter**
- **Clock skew**
  - The (knowable) difference in clock arrival times at each flip-flop
  - Caused mainly by imperfect balancing of clock tree/mesh
- **Clock jitter**
  - The random (unknowable) difference in clock arrival times at each flip-flop
  - Caused by on-die process, V_{dd}, temperature variation, PLL jitter, crosstalk.
- **Clock tree to minimize clock uncertainty**

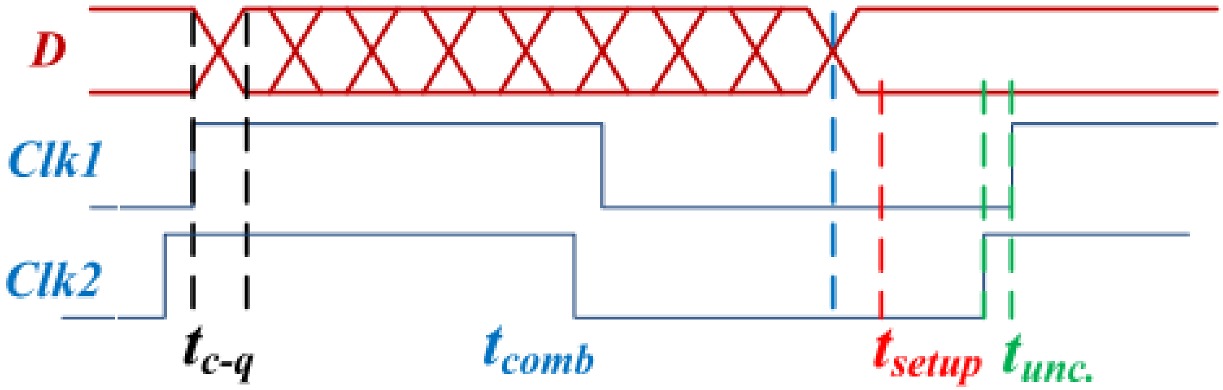


Clock uncertainty

- □Clock uncertainty = skew ±jitter

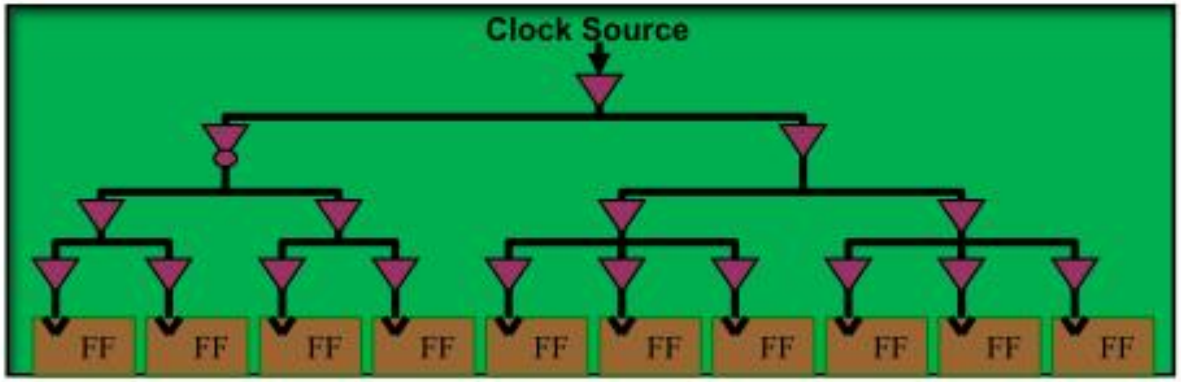


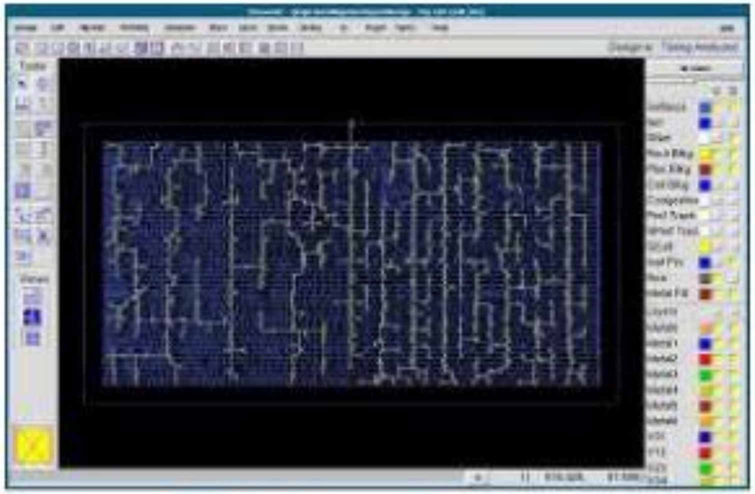
- □Timing analysis with clock uncertainty



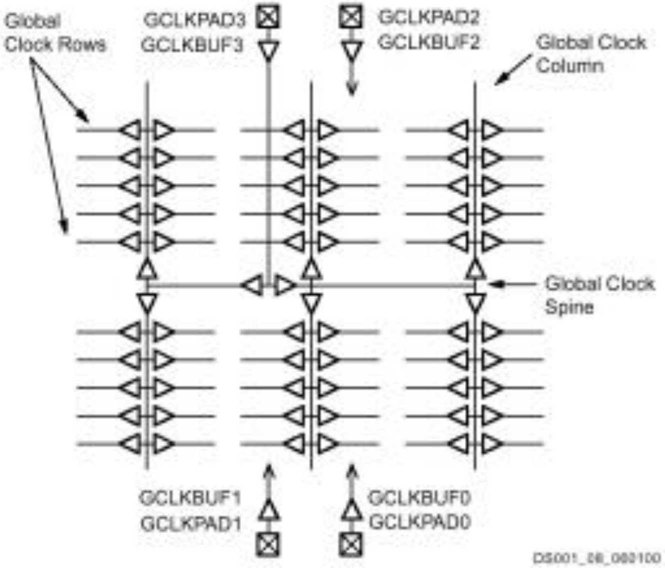
$T_{\text{Combinational logic}}$ + FF_{lauch}(clk -> Q) < Clock Period - FF $_{\text{tSetup}}$ - Clock Uncertainty

Clock tree



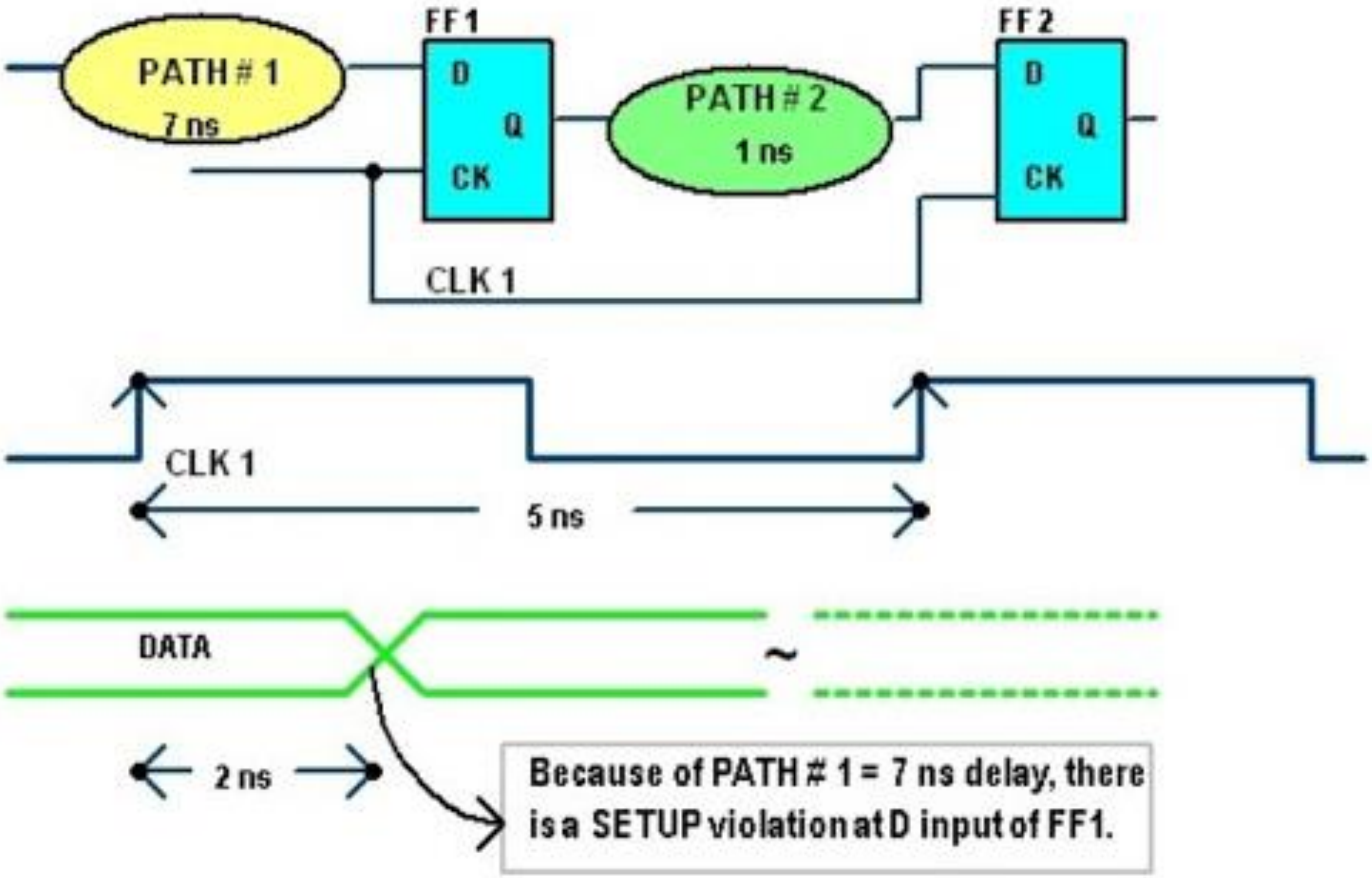


Clock-tree Asic



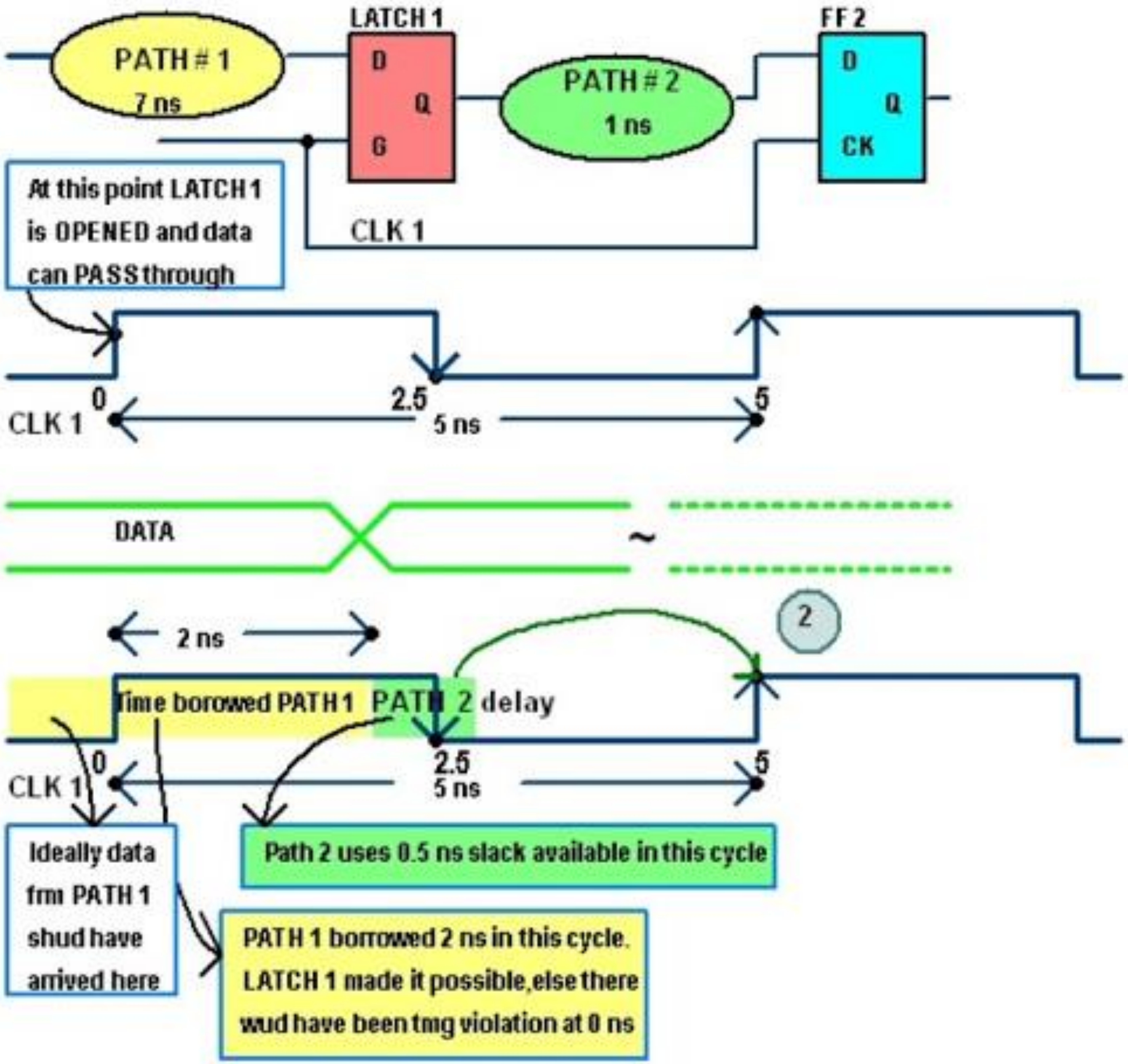
Global clock network in Xilinx FPGA

Time borrow



Because of PATH # 1 = 7 ns delay, there is a SETUP violation at D input of FF1.

Time borrow



At this point LATCH 1 is OPENED and data can PASS through

Ideally data frm PATH 1 shud have arrived here

Path 2 uses 0.5 ns slack available in this cycle

PATH 1 borrowed 2 ns in this cycle. LATCH 1 made it possible,else there wud have been tmg violation at 0 ns

Outline

- Sequential vs. Combinational
- Synchronous vs. Asynchronous
- Basic Storage Elements
- Timing
- **Folding & Pipeline**

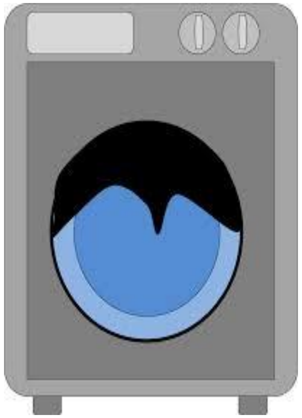
Pipeline

□Acknowledgement:

- The following slides have been provided by Prof. Ward in September 2004.
- Reformatting of PowerPoint and addition of two more slide done September 2007 by Jens Sparsø.
- Slides are used in DTU course 02154 Digital Systems Engineering (fall 2008).
- Due to Joachim Rodrigues' position at DTU, I used some of the slides in EITF35.

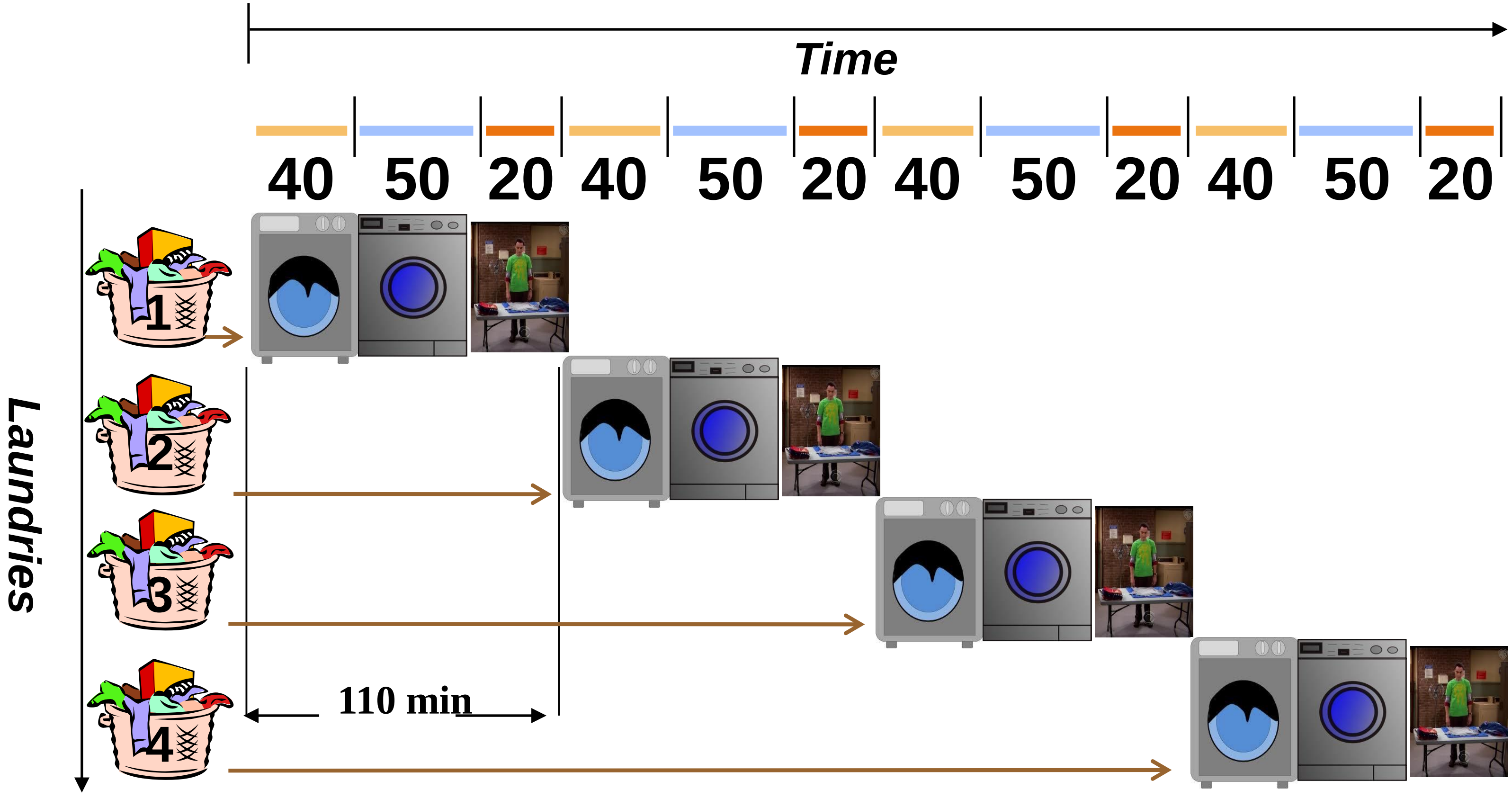
Pipelining

- Start again from laundry room



- **Small laundry has one washer, one dryer and one folder, it takes 110 minutes to finish one load:**
  - Washer takes 40 minutes
  - Dryer takes 50 minutes
  - “Folding” takes 20 minutes
- **Need to do 4 laundries**

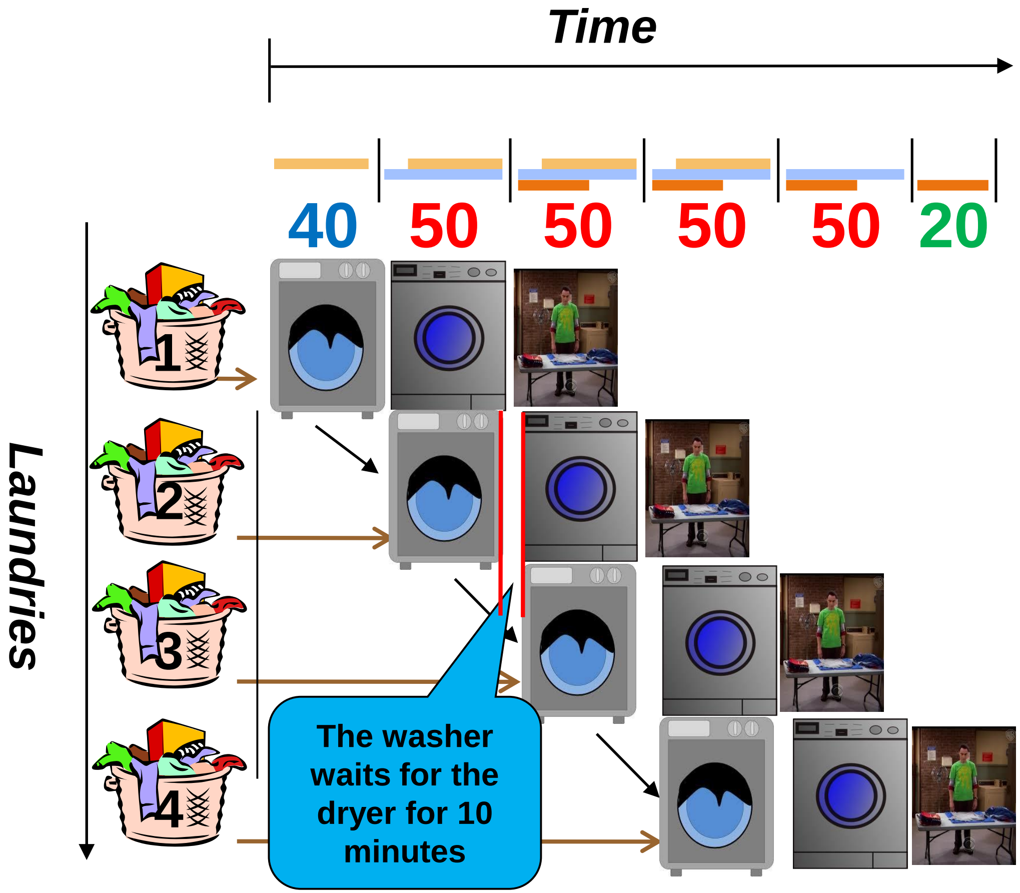
Not very smart way...



Total = N*(Washer+ Dryer+Folder)

= 440 mins

If we pipelining

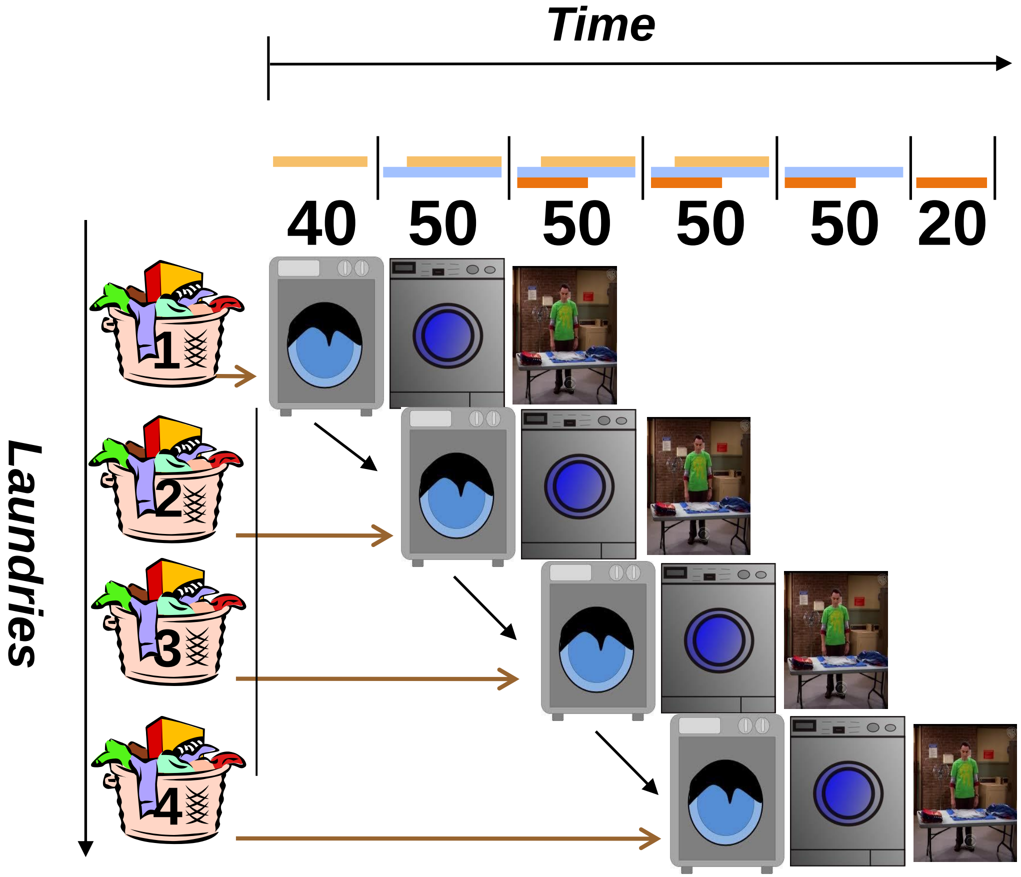


The washer waits for the dryer for 10 minutes

Total = Washer+N*Max(Washer,Dryer,Folder)+Folder

= ____260____ mins

Pipeline Facts



- □Multiple tasks operating simultaneously
- □Pipelining doesn’t help latency of single task, it helps throughput of entire workload
- □Pipeline rate limited by slowest pipeline stage
- □Unbalanced lengths of pipe stages reduces speedup
- □Potential speedup ∝ Number of pipe stages

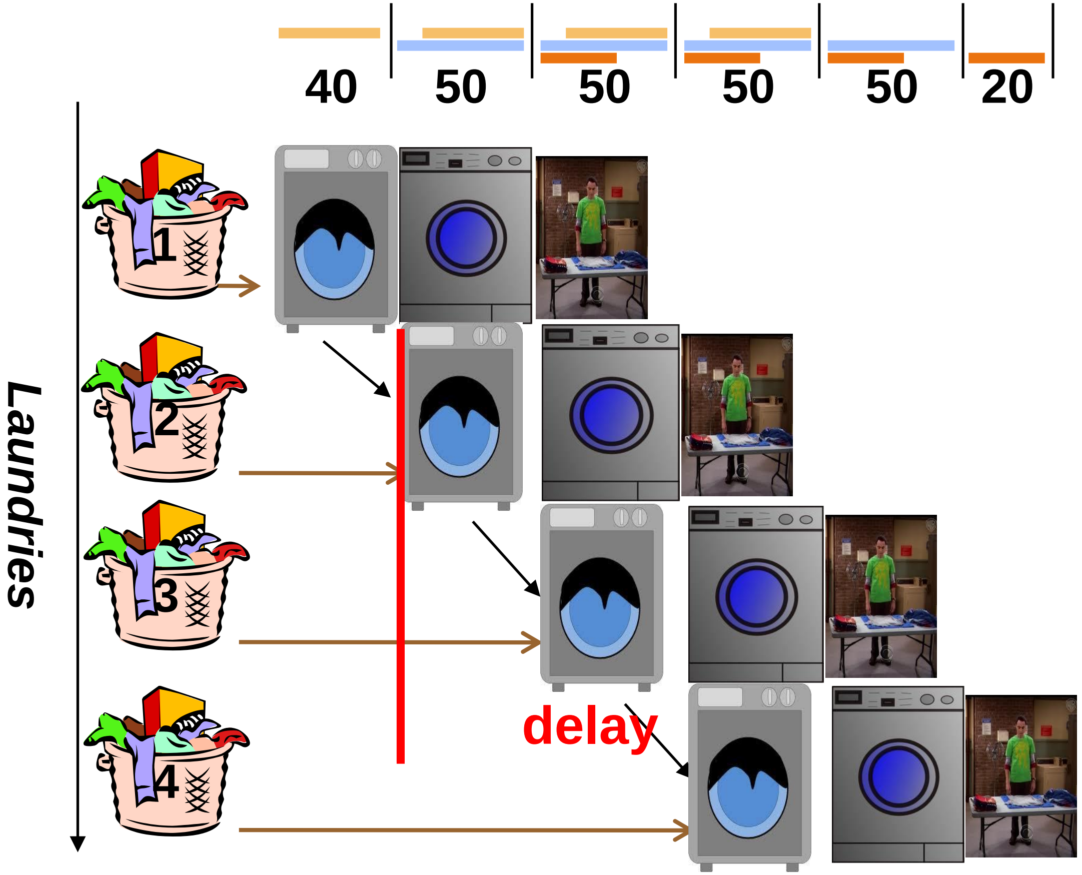
Some definitions

Very Important!

- **Latency: The delay from when an input is established until the output associated with that input becomes valid.**

(non-pipeline Laundry = ___**110**___ *mins)*

(pipeline Laundry = ___**120**___ *mins)*



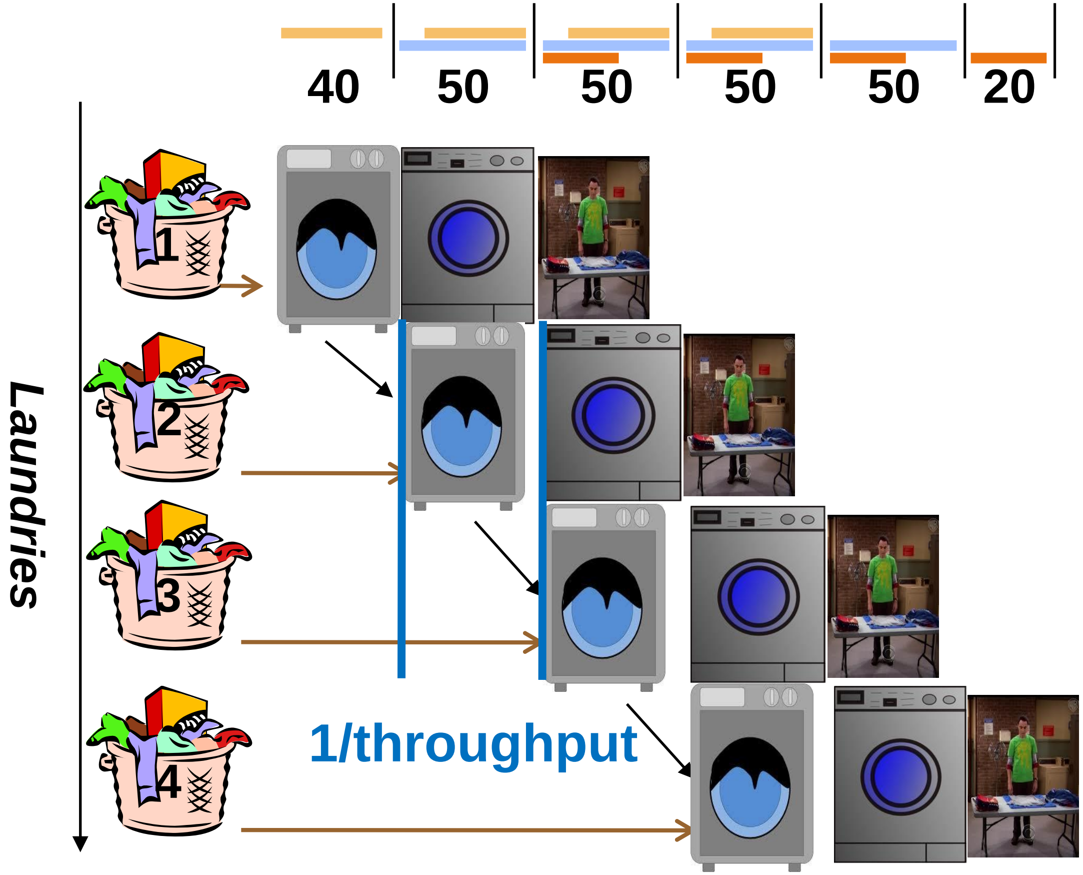
Some definitions

Very Important!

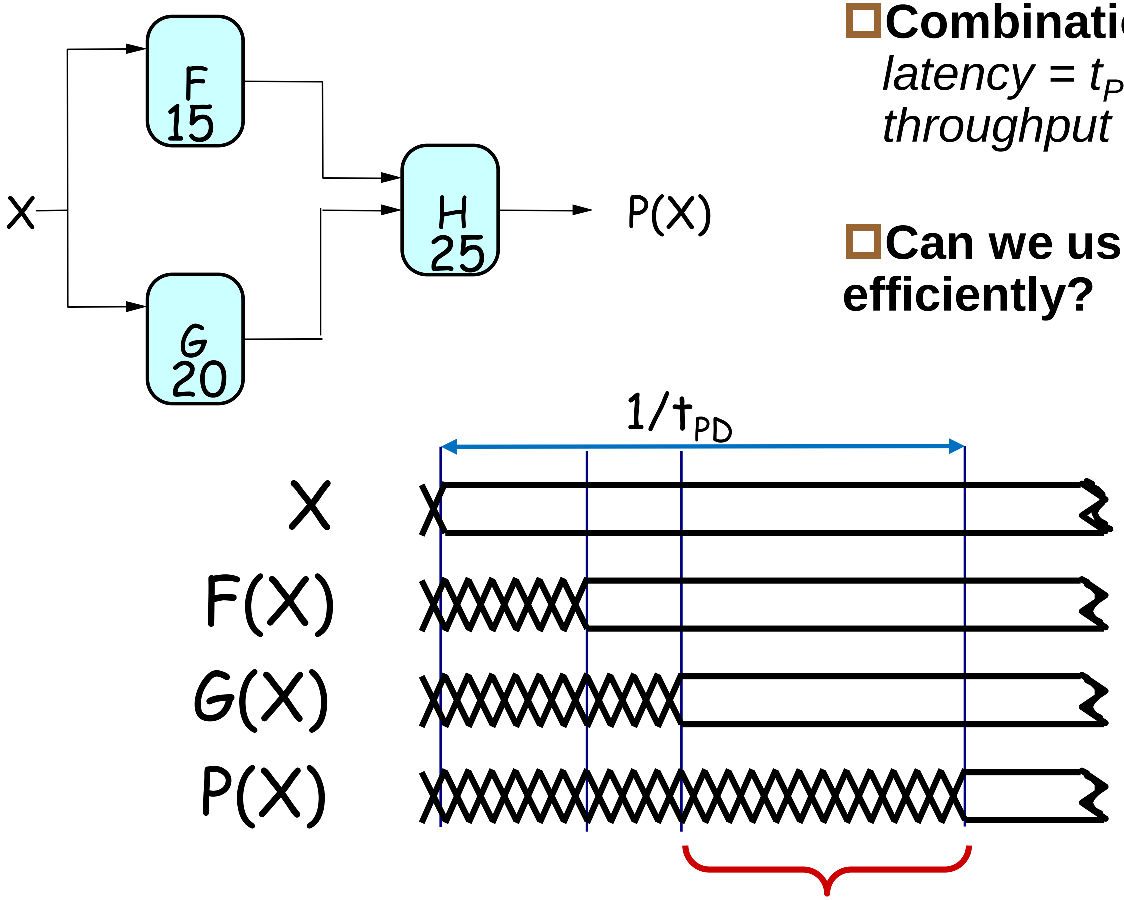
- **Throughput: The rate of which inputs or outputs are processed or how frequently a laundry can be loaded**

(non-pipeline Laundry = **1/110** *outputs/min)*

(pipeline Laundry = **1/50** *outputs/min)*



Okay, back to circuits…



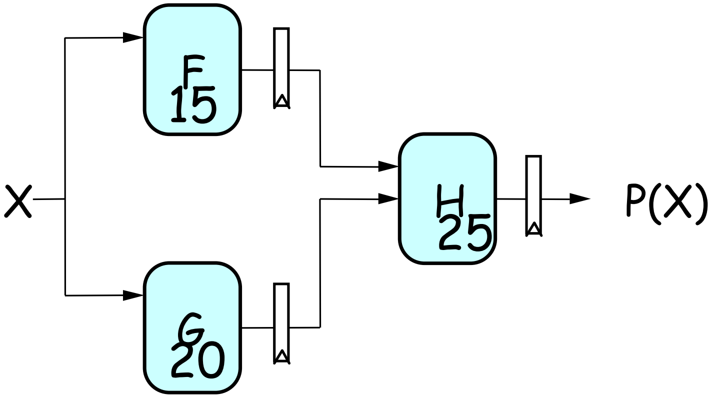
F & G are "idle", just holding their outputs stable while H performs its computation

- **Combinational logic:**
  latency = t_{PD},
  throughput = $1/t_{PD}$.

- **Can we use the hardware more efficiently?**

Pipelined Circuits

use registers to hold H's input stable!



□**Pipelined circuit:**

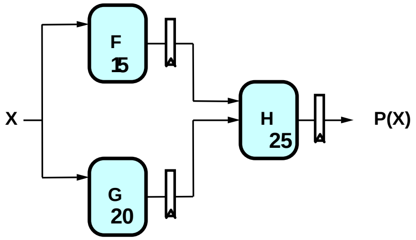
- 2-stage *pipeline*: if we have a valid input X during clock cycle j, P(X) is valid during clock j+2.
- Now F & G can be working on input X_{i+1} while H is performing its computation on X_i.

Exercise
1min

Suppose F, G, H have propagation delays of 15, 20, 25 ns and we are using ideal zero-delay registers:

	latency	throughput
un-pipelined	45	1/45
2-stage pipelined	50	1/25
	worse	better

Pipeline timing diagrams



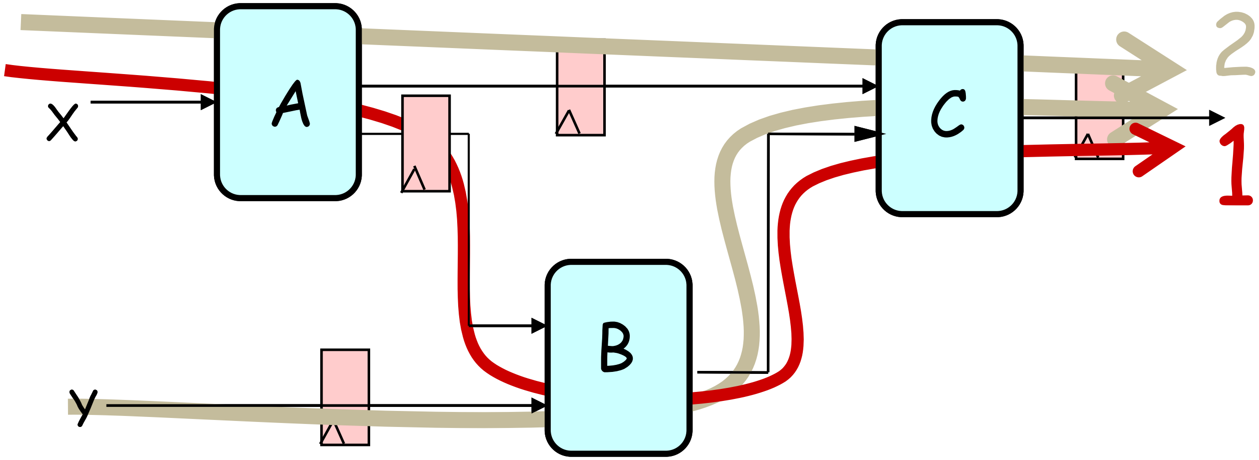
Clock cycle

Pipeline stages

	i	i+1	i+2	i+3	
Input	X_i	X_{i+1}	X_{i+2}	X_{i+3}	…
F Reg G Reg		$F(X_i)$ $G(X_i)$	$F(X_{i+1})$ $G(X_{i+1})$	$F(X_{i+2})$ $G(X_{i+2})$	…
H Reg			$H(X_i)$	$H(X_{i+1})$	$H(X_{i+2})$

Ill-formed pipelines

Consider a BAD job of pipelining:



Problem:

Some paths from inputs to outputs had 2 registers, and some had only 1!

Make sure every paths have been pipelined with same stages

Combinational, Folding and Pipelined

□ Combinational Circuits

- Advantage: low latency
- Disadvantage: low throughput, more hardware, low utilization

□ Folding

- Advantage: less hardware, high utilization
- Disadvantage: high latency, limited application

□ Pipeline

- Advantage: very high throughput
- Disadvantages: pipeline latency, more hardware

Thanks!